



NIST Special Publication
NIST SP 800-90C 3pd

Recommendation for Random Bit Generator (RBG) Constructions

Third Public Draft (3pd)

Elaine Barker
John Kelsey
Kerry McKay
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This publication is available free of charge from:
<https://doi.org/10.6028/NIST.SP.800-90C.3pd>

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Computer Security Division
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September 2022



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National Institute of Standards and Technology
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Publication History

Approved by the NIST Editorial Review Board on YYYY-MM-DD [will be added in final published version]

How to Cite this NIST Technical Series Publication:

Barker EB, Kelsey JM, McKay KA, Roginsky AL, Sönmez Turan M (2022) Recommendation for Random Bit Generator (RBG) Constructions. (National Institute of Standards and Technology, Gaithersburg, MD), NIST Special Publication (SP) 800-90C 3pd. <https://doi.org/10.6028/NIST.SP.800-90c.3pd>

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76 **Public Comment Period**
77 September 7, 2022 – December 7, 2022

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Abstract

The NIST Special Publication (SP) 800-90 series of documents supports the generation of high-quality random bits for cryptographic and non-cryptographic use. SP 800-90A specifies several deterministic random bit generator (DRBG) mechanisms based on cryptographic algorithms. SP 800-90B provides guidance for the development and validation of entropy sources. This document (SP 800-90C) specifies constructions for the implementation of random bit generators (RBGs) that include DRBG mechanisms as specified in SP 800-90A and that use entropy sources as specified in SP 800-90B. Constructions for three classes of RBGs (namely, RBG1, RBG2, and RBG3) are specified in this document.

Keywords

deterministic random bit generator (DRBG); entropy; entropy source; random bit generator (RBG); randomness source; RBG1 construction; RBG2 construction; RBG3 construction; subordinate DRBG (sub-DRBG).

Note to Reviewers

1. This draft of SP800-90C describes three RBG constructions. Note that in this draft, a non-deterministic random bit generator (NRBG) is presented as an RBG3 construction.

Question: *In a future revision of SP 800-90C, should other constructions be included?*

This version of SP 800-90C does not address the use of an RBG software implementation in which a) a cryptographic library or an application is loaded into a system and b) the software accesses entropy sources or RBGs already associated with the system for its required randomness. NIST intends to address this situation in the near future.

2. The RBG constructions provided in this draft use NIST-**approved** cryptographic primitives (such as block ciphers and hash functions) as underlying components. Note that non-vetted conditioning components may be used within SP 800-90B entropy sources.

Although NIST still allows three-key TDEA as a block-cipher algorithm, Section 4 of [[SP800-131A](#)] indicates that its use is deprecated through 2023 and will be disallowed thereafter for applying cryptographic protection. This document (i.e., SP 800-90C) **does not approve** the use of three-key TDEA in an RBG.

Although SHA-1 is still approved by NIST, NIST is planning to remove SHA-1 from a future revision of FIPS 180-4, so the SP 800-90 series will not be including the use of SHA-1.

The use of the SHA-3 hash functions are **approved** in SP 800-90C for Hash_DRBG and HMAC_DRBG but are not currently included in [[SP800-90A](#)]. SP 800-90A will be revised to exclude the use of TDEA and SHA-1 and include the use of the SHA-3 family of hash functions.

3. Since the projected date for requiring a minimum security strength of 128 bits for U.S. Government applications is 2030 (see [[SP800-57Part1](#)]), RBGs are only specified to provide 128, 192, and 256 bits of security strength (i.e., the 112-bit security strength has been removed). Note that a consuming application may still request a lower security strength, but the RBG output will be generated at the instantiated security strength.

4. Guidance is provided for accessing entropy sources and for obtaining full-entropy bits using the output of an entropy source that does not inherently provide full-entropy output (see [Section 3.3](#)).

5. SP 800-90A requires that when instantiating a CTR_DRBG without a derivation function, the randomness source needs to provide full-entropy bits (see SP 800-90A). However, this draft (SP 800-90C) relaxes this requirement in the case of an RBG1 construction, as specified in [Section 4](#). In this case, the external randomness source may be another RBG construction. An addendum to SP 800-90A has been prepared as a temporary specification in SP 800-90C, but SP 800-90A will be revised in the future to accommodate this change.

6. The DRBG used in RBG3 constructions supports a security strength of 256 bits. The RBG1 and RBG2 constructions may support any valid security strength (i.e., 128, 192 or 256 bits).

7. SP 800-90A currently allows the acquisition of a nonce (when required) for DRBG instantiation from any randomness source. However, SP 800-90C does not include an explicit requirement for the generation of a nonce when instantiating a DRBG. Instead, additional bits

beyond those needed for the security strength are acquired from the randomness source. SP 800-90A will be revised to agree with this change.

8. SP 800-90C allows the use of both physical and non-physical entropy sources. See the definitions of physical and non-physical entropy sources in [Appendix E](#). Also, multiple validated entropy sources may be used to provide entropy, and two methods are provided in [Section 2.3](#) for counting the entropy provided in a bitstring.

9. The CMVP is considering providing information on an entropy source validation certificate that indicates whether an entropy source is physical or non-physical.

10. The CMVP is developing a program to validate entropy sources against SP 800-90B with the intent of allowing the re-use of those entropy sources in different RBG implementations.

Question: *Are there any issues that still need to be addressed in SP 800-90C to allow the re-use of validated entropy sources in different RBG implementations? Note that in many cases, specific issues need to be addressed in the FIPS 140 implementation guide rather than in this document.*

Call for Patent Claims

This public review includes a call for information on essential patent claims (claims whose use would be required for compliance with the guidance or requirements in this Information Technology Laboratory (ITL) draft publication). Such guidance and/or requirements may be directly stated in this ITL Publication or by reference to another publication. This call also includes disclosure, where known, of the existence of pending U.S. or foreign patent applications relating to this ITL draft publication and of any relevant unexpired U.S. or foreign patents.

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The assurance shall also indicate that it is intended to be binding on successors-in-interest regardless of whether such provisions are included in the relevant transfer documents.

Such statements should be addressed to: rbg_comments@nist.gov

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317 **Acknowledgments**

318 The National Institute of Standards and Technology (NIST) gratefully acknowledges and
319 appreciates contributions from Chis Celi (NIST); Darryl Buller, Aaron Kaufer, and Mike Boyle
320 (National Security Agency); Werner Schindler, Matthias Peter, Johannes Mittman (Bundesamt für
321 Sicherheit in der Informationstechnik); and the members of the Cryptographic Module User Forum
322 (CMUF) for assistance in the development of this Recommendation. NIST also thanks the many
323 contributions by the public and private sectors.

1. Introduction and Purpose

Cryptography and security applications make extensive use of random bits. However, the generation of random bits is challenging in many practical applications of cryptography.

The National Institute of Standards and Technology (NIST) developed the Special Publication (SP) 800-90 series to support the generation of high-quality random bits for both cryptographic and non-cryptographic purposes. The SP 800-90 series consists of three parts:

- SP 800-90A, *Recommendation for Random Number Generation Using Deterministic Random Bit Generators*, specifies several **approved** deterministic random bit generator (DRBG) mechanisms based on **approved** cryptographic algorithms that – once provided with seed material that contains sufficient entropy – can be used to generate random bits suitable for cryptographic applications.
- SP 800-90B, *Recommendation for the Entropy Sources Used for Random Bit Generation*, provides guidance for the development and validation of entropy sources – mechanisms that generate entropy from physical or non-physical noise sources and that can be used to generate the input for the seed material needed by a DRBG or for input to an RBG.
- SP 800-90C, *Recommendation for Random Bit Generator (RBG) Constructions*, specifies constructions for random bit generators (RBGs) using entropy sources that comply with SP 800-90B and DRBGs that comply with SP 800-90A. Three classes of RBGs are specified in this document (see Sections 5, 6, and 7). SP 800-90C also provides high-level guidance for testing RBGs for conformance to this Recommendation.

The RBG constructions defined in this Recommendation consist of two main components: the *entropy sources* that generate true random variables (variables that may be biased, i.e., each possible outcome does not need to have the same chance of occurring) and the DRBGs that ensure that the outputs of the RBG are indistinguishable from the ideal distribution to a computationally bounded adversary.

Throughout this document, the phrase “this Recommendation” refers to the aggregate of SP 800-90A, SP 800-90B, and SP 800-90C, while the phrase “this document” refers only to SP 800-90C.

SP 800-90C has been developed in coordination with NIST’s Cryptographic Algorithm Validation Program (CAVP) and Cryptographic Module Validation Program (CMVP). The document uses “**shall**” and “**must**” to indicate requirements and uses “**should**” to indicate an important recommendation. The term “**shall**” is used when a requirement is testable by a testing lab during implementation validation using operational tests or a code review. The term “**must**” is used for requirements that may not be testable by the CAVP or CMVP. An example of such a requirement is one that demands certain actions and/or considerations from a system administrator. Meeting these requirements can be verified by a CMVP review of the cryptographic module’s documentation. If the requirement is determined to be testable at a later time (e.g., after SP 800-90C is published and before it is revised), the CMVP will so indicate in the [Implementation Guidance](#) for [FIPS 140](#), *Security Requirements for Cryptographic Modules*.

1.1. Audience

The intended audience for this Recommendation includes 1) developers who want to design and implement RBGs that can be validated by NIST's CMVP and CAVP, 2) testing labs that are accredited to perform the validation tests and the evaluation of the RBG constructions, and 3) users who install RBGs in systems.

1.2. Document Organization

This document is organized as follows:

- [Section 2](#) provides background and preliminary information for understanding the remainder of the document.
- [Section 3](#) provides guidance on accessing and handling entropy sources, including the external conditioning of entropy-source output.
- Sections [4](#), [5](#), and [6](#) specify the RBG constructions.
- [Section 7](#) discusses health and implementation-validation testing.
- [References](#) contains a list of papers and publications cited in this document.

The following informational appendices are also provided:

- [Appendix A](#) provides discussions on entropy versus security strength.
- [Appendix B](#) provides examples of each RBG construction.
- [Appendix C](#) is an addendum to SP 800-90A that includes two additional derivation functions that may be used with the CTR_DRBG. These functions will be moved into SP 800-90A as part of the next revision of that document.
- [Appendix D](#) provides a list of abbreviations, symbols, functions, and notations used in this document.
- [Appendix E](#) provides a glossary with definitions for terms used in this document.

2. General Information

2.1. RBG Security

Ideal randomness sources generate identically distributed and independent uniform random bits that provide full-entropy outputs (i.e., one bit of entropy per output bit). Real-world RBGs are designed with a security goal of *indistinguishability* from the output of an ideal randomness source. That is, given some limits on an adversary's data and computing power, it is expected that there is no adversary that can reliably distinguish between RBG outputs and outputs from an ideal randomness source.

Consider an adversary that can perform 2^w computations (typically, these are guesses of the RBG's internal state) and is given an output sequence from either an RBG with a security strength of s bits (where $s \geq w$) or an ideal randomness source. It is expected that an adversary has no better probability of determining which source was used for its random bits than

$$1/2 + 2^{w-s-1} + \varepsilon,$$

where ε is negligible. In this Recommendation, the size of the output is limited to 2^{64} output bits and $\varepsilon \leq 2^{-32}$.

An RBG that has been designed to support a security strength of s bits is suitable for any application with a targeted security strength that does not exceed s . An RBG that is compliant with this Recommendation can support requests for output with a security strength of 128, 192, or 256 bits, except for an RBG3 construction (as described in [Section 6](#)), which can provide full-entropy output.

A bitstring with full entropy has an amount of entropy equal to its length. Full-entropy bitstrings are important for cryptographic applications, as these bitstrings have ideal randomness properties and may be used for any cryptographic purpose. They may be truncated to any length such that the amount of entropy in the truncated bitstring is equal to its length. However, due to the difficulty of generating and testing full-entropy bitstrings, this Recommendation assumes that a bitstring has full entropy if the amount of entropy per bit is at least $1 - \varepsilon$, where ε is at most 2^{-32} . [NISTIR 8427](#)¹ provides a justification for the selection of ε .

2.2. RBG Constructions

A *construction* is a method of designing an RBG or some component of an RBG to accomplish a specific goal. Three classes of RBG constructions are defined in this document: RBG1, RBG2, and RBG3 (see [Table 1](#)). Each RBG includes a DRBG from [\[SP800-90A\]](#) and is based on the use of a randomness source that is validated for compliance with [\[SP800-90B\]](#) or SP 800-90C. Once instantiated, a DRBG can generate output at a security strength that does not exceed the DRBG's instantiated security strength.

¹ See NISTIR 8427, Discussion on the Full Entropy Assumption of SP 800-90 series.

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Table 1. RBG Capabilities

Construction	Internal Entropy Source	Prediction Resistance	Full Entropy	Type of randomness source
RBG1	No	No	No	Physical
RBG2	Yes	Yes ^a	No	Physical or Non-physical
RBG3	Yes	Yes ^a	Yes	Physical

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^a If sufficient entropy is available or can be obtained when reseeding the RBG's DRBG.

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1. An RBG1 construction (see [Section 4](#)) does not have access to a randomness source after instantiation. It is instantiated once in its lifetime over a secure channel from an external RBG with appropriate security properties. An RBG1 construction does not support reseeding and cannot provide *prediction resistance* as described in [Section 2.4.2](#) and [\[SP800-90A\]](#). The construction can be used to initialize subordinate DRBGs.

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2. An RBG2 construction (see [Section 5](#)) includes one or more entropy sources that are used to instantiate and reseed the DRBG within the construction. This construction can provide prediction resistance (see [Section 2.4.2](#) and [\[SP800-90A\]](#)) when sufficient entropy is available or can be obtained from the RBG's entropy source(s) at the time that prediction resistance is requested. The construction has two variants that depend on the type of entropy source(s) employed (i.e., physical and non-physical).

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3. An RBG3 construction is designed to provide output with a security strength equal to the requested length of its output by producing outputs that have full entropy (i.e., an RBG designed as an RBG3 construction can, in effect, support all security strengths) (see [Section 2.1](#)). This construction provides prediction resistance and has two types, namely RBG3(XOR) and RBG3(RS).

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- a. An RBG3(XOR) construction (see [Section 6.2](#)) combines the output of one or more validated entropy sources with the output of an instantiated, **approved** DRBG using an exclusive-or (XOR) operation.

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- b. An RBG3(RS) construction (see [Section 6.3](#)) uses one or more validated entropy sources to provide randomness input for the DRBG by continuously reseeding.

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This document also provides constructions for 1) subordinate DRBGs (sub-DRBGs) that are instantiated and possibly reseeded by an RBG1 construction (see [Section 4.3](#)) and 2) acquiring entropy from an entropy source and conditioning the output to provide a bitstring with full entropy (see [Section 3.3](#)). SP 800 90A provides constructions for instantiating and reseeding DRBGs and requesting the generation of pseudorandom bitstrings.

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All constructions in SP 800-90C are described in pseudocode. These pseudocode conventions are not intended to constrain real-world implementations but to provide a consistent notation to describe the constructions. By convention, unless otherwise specified, integers are unsigned 32-bit values, and when used as bitstrings, they are represented in the big-endian format.

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451 2.3. Sources of Randomness for an RBG

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The RBG constructions specified in this document are based on the use of validated entropy sources. Some RBG constructions (e.g., the RBG3 construction) access these entropy sources

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directly to obtain entropy. Other constructions (e.g., the RBG1 construction) fulfill their entropy requirements by accessing another RBG as a randomness source. In this case, the source RBG may include its own entropy source.

SP 800 90B provides guidance for the development and validation of entropy sources – mechanisms that provide entropy for an RBG. Validated entropy sources (i.e., entropy sources that have been successfully validated by the CMVP as complying with SP 800-90B) provide fixed-length outputs and have been validated as reliably providing a specified minimum amount of entropy for each output (e.g., each eight-bit output has been validated as providing at least five bits of entropy).²

An entropy source is a *physical entropy source* if the primary noise source of the entropy source is physical – that is, it uses dedicated hardware to provide entropy (e.g., from ring oscillators, thermal noise, shot noise, jitter, or metastability). Similarly, a validated entropy source is a *non-physical entropy source* if the primary noise source of the entropy source is non-physical – that is, entropy is provided by system data (e.g., the entropy present in the RAM data or system time). The entropy-source type is certified during SP 800-90B validation.

One or more validated entropy sources are used to provide entropy for instantiating and reseeding the DRBGs in RBG2 or RBG3 constructions or used by an RBG3 construction to generate output upon request by a consuming application.

An implementation could be designed to use a combination of physical and non-physical entropy sources. When requests are made to the sources, bitstring outputs are concatenated until the amount of entropy in the concatenated bitstring meets or exceeds the request. Two methods are provided for counting the entropy provided in the concatenated bitstring.

Method 1: The RBG implementation includes one or more physical entropy sources, and one or more non-physical entropy sources may also be included in the implementation. However, only the entropy in a bitstring that is provided from physical entropy sources is counted toward fulfilling the amount of entropy requested in an entropy request. Any entropy in a bitstring that is provided by a non-physical entropy source is not counted, even if bitstrings produced by the non-physical entropy source are included in the concatenated bitstring that is used by the RBG.

Method 2: The RBG implementation includes one or more non-physical entropy sources, and one or more physical entropy sources may also be included in the implementation. The entropy from both non-physical entropy sources and (if present) physical entropy sources is counted when fulfilling an entropy request.

Example: Let pes_i be the i^{th} output of a physical entropy source, and $npes_i$ be the j^{th} output of a non-physical entropy source. If an implementation consists of one physical and one non-physical entropy source, and a request has been made for 128 bits of entropy, the concatenated bitstring might be something like:

$$pes_1 \parallel pes_2 \parallel npes_1 \parallel pes_3 \parallel \dots \parallel npes_m \parallel pes_n,$$

which is the concatenated output of the physical and non-physical entropy sources.

² Note that this document also discusses the use of non-validated entropy sources. When discussing such entropy sources, “non-validated” will always precede “entropy sources.” The use of the term “validated entropy source” may be shortened to just “entropy source” to avoid repetition.

According to Method 1, only the entropy in $pes_1, pes_2, \dots, pes_n$ would be counted toward fulfilling the 128-bit request. Any entropy in $npes_1, \dots, npes_m$ is not counted.

According to Method 2, all of the entropy in $pes_1, pes_2, \dots, pes_n$ and in $npes_1, npes_2, \dots, npes_m$ is counted. Since the entropy from both non-physical and physical entropy sources is counted in Method 2, the concatenated output string is expected to be shorter compared to that credited using Method 1.

When multiple entropy sources are used, there is no requirement on the order in which the entropy sources are accessed or the number of times that each entropy source is accessed to fulfill an entropy request (e.g., if two physical entropy sources are used, it is possible that a request would be fulfilled by only one of the entropy sources because entropy is not available at the time of the request from the other entropy source). However, the Method 1 or Method 2 criteria for counting entropy still applies.

This Recommendation assumes that the entropy produced by a validated physical entropy source is generally more reliable than the entropy produced by a validated non-physical entropy source since non-physical entropy sources are typically influenced by human actions or network events, the unpredictability of which is difficult to accurately quantify. Therefore, Method 1 is considered to provide more assurance that the concatenated bitstring actually contains at least the requested amount of entropy (128 bits for the example). Note that RBG2(P) and RBG3 constructions only count the entropy using Method 1 (see Sections 5 and 6).

2.4. DRBGs

Approved DRBG designs are specified in [SP800-90A]. A DRBG includes instantiate, generate, and health-testing functions and may include reseed and uninstantiate functions. The instantiation of a DRBG involves acquiring sufficient randomness to initialize the DRBG to support a targeted security strength and establish the internal state, which includes the secret information for operating the DRBG. The generate function produces output upon request and updates the internal state. Health testing is used to determine that the DRBG continues to operate correctly. Reseeding introduces fresh entropy into the DRBG's internal state and is used to recover from a potential (or actual) compromise (see Section 2.4.2 for additional discussion). An uninstantiate function is used to terminate a DRBG instantiation and destroy the information in its internal state.

2.4.1. DRBG Instantiations

A DRBG implementation consists of software code, hardware, or both hardware and software that is used to implement a DRBG design. The same implementation can be used to create multiple “copies” of the same DRBG (e.g., for different purposes) without replicating the software code or hardware. Each “copy” is a separate instantiation of the DRBG with its own internal state that is accessed via a state handle that is unique to that instantiation (see Figure 1). Each instantiation may be considered a different DRBG, even though it uses the same software code or hardware.

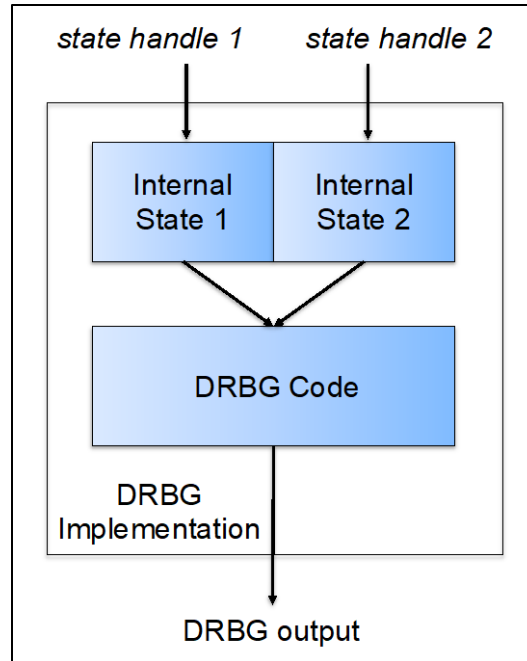


Fig. 1. DRBG Instantiations

Each DRBG instantiation is initialized with input from some randomness source that establishes the security strengths that can be supported by the DRBG. During this process, an optional but recommended personalization string may also be used to differentiate between instantiations in addition to the output of the randomness source. The personalization string could, for example, include information particular to the instantiation or contain entropy collected during system activity (e.g., from a non-validated entropy source). An implementation **should** allow the use of a personalization string. More information on personalization strings is provided in [SP800-90A].

A DRBG may be implemented to accept further input during operation from the randomness source (e.g., to reseed the DRBG) and/or additional input from inside or outside of the cryptographic module that contains the DRBG. This additional input could, for example, include information particular to a request for generation or reseeding or could contain entropy collected during system activity (e.g., from a validated or non-validated entropy source).³

2.4.2. DRBG Reseeding, Prediction Resistance, and Recovery from Compromise

Under some circumstances, the internal state of an RBG (containing the RBG's secret information) could be leaked to an adversary. This would typically happen as the result of a side-channel attack or tampering with a hardware device, and it may not be detectable by the RBG or any consuming application.

All DRBGs in [SP800-90A] are designed with *backtracking resistance* – that is, learning the DRBG's current internal state does not provide knowledge of previous outputs. Since all RBGs in SP 800-90C are based on the use of SP 800-90A DRBGs, they also inherit this property. However,

³ Entropy provided in additional input does not affect the instantiated security strength of the DRBG instantiation. However, it is good practice to include any additional entropy when available to provide more security.

once the secret information within the DRBG's internal state is compromised, all future DRBG outputs are known to the adversary unless the DRBG is reseeded – a process that returns the DRBG to a non-compromised state.

A DRBG is reseeded when at least s bits of fresh entropy are used to update the internal state (where s is the security strength of the DRBG) so that the updated internal state is unknown and extremely unlikely to be correctly guessed. A DRBG that has been reseeded has *prediction resistance* against an adversary who knows its previous internal state. Reseeding may be performed upon request from a consuming application (either an explicit request for reseeded or a request for the generation of bits with prediction resistance); on a fixed schedule based on time, number of outputs, or events; or as sufficient entropy becomes available.

Although reseeded provides fresh entropy bits that are incorporated into an already instantiated DRBG at a security strength of s bits, this Recommendation does not consider the reseed process as increasing the DRBG's security strength. For example, a reseed of a DRBG that has been instantiated to support a security strength of 128 bits does not increase the DRBG's security strength to 256 bits when reseeded with 128 bits of fresh entropy.

An RBG1 construction has no access to a randomness source after instantiation and so cannot be reseeded or recover from a compromise (see [Section 4](#)). Thus, it can never provide prediction resistance.

An RBG2 construction contains an entropy source that is used to reseed the DRBG within the construction (see [Section 5](#)) and recover from a possible compromise of the RBG's internal state. Prediction resistance may be requested by a consuming application during a request for the generation of (pseudo) random bits. If sufficient entropy can be obtained from the entropy source(s) at that time, the DRBG is reseeded before the requested bits are generated. If sufficient entropy is not available, an error indication is returned, and no bits are generated for output. Therefore, it is recommended that prediction resistance not be claimed for an RBG implementation unless sufficient entropy is reliably available upon request.

An RBG3 construction is provided with fresh entropy for every RBG output (see [Section 6](#)). As a result, every output from an RBG3 construction has prediction resistance.

For a more complete discussion of backtracking and prediction resistance, see [\[SP800-90A\]](#).

2.5. RBG Security Boundaries

An RBG exists within a *conceptual* RBG security boundary that **should** be defined with respect to one or more threat models that include an assessment of the applicability of an attack and the potential harm caused by the attack. The RBG security boundary **must** be designed to assist in the mitigation of these threats using physical or logical mechanisms or both.

The primary components of an RBG are a randomness source (i.e., an entropy source or an RBG construction), a DRBG, and health tests for the RBG. RBG input (e.g., entropy bits and a personalization string) **shall** enter an RBG only as specified in the functions described in [Section 2.8](#). The security boundary of a DRBG is discussed in [\[SP800-90A\]](#). The security boundary for an entropy source is discussed in [\[SP800-90B\]](#). Both the entropy source and the DRBG contain their own health tests within their respective security boundaries.

[Figure 2](#) shows an RBG implemented within a [FIPS 140]-validated cryptographic module. The RBG security boundary **shall** either be the same as the cryptographic module boundary or be completely contained within that boundary. The data input may be a personalization string or additional input (see [Section 2.4.1](#)). The data output is status information and possibly random bits or a state handle. Within the RBG security boundary of the figure are an entropy source and a DRBG – each with its own (conceptual) security boundary. An entropy-source security boundary includes a noise source, health tests, and (optionally) a conditioning component. A DRBG security boundary contains the chosen DRBG, memory for the internal state, and health tests. An RBG security boundary contains health tests and may also contain an (optional) external conditioning function. The RBG2 and RBG3 constructions in [Sections 5](#) and [6](#), respectively, use this model.

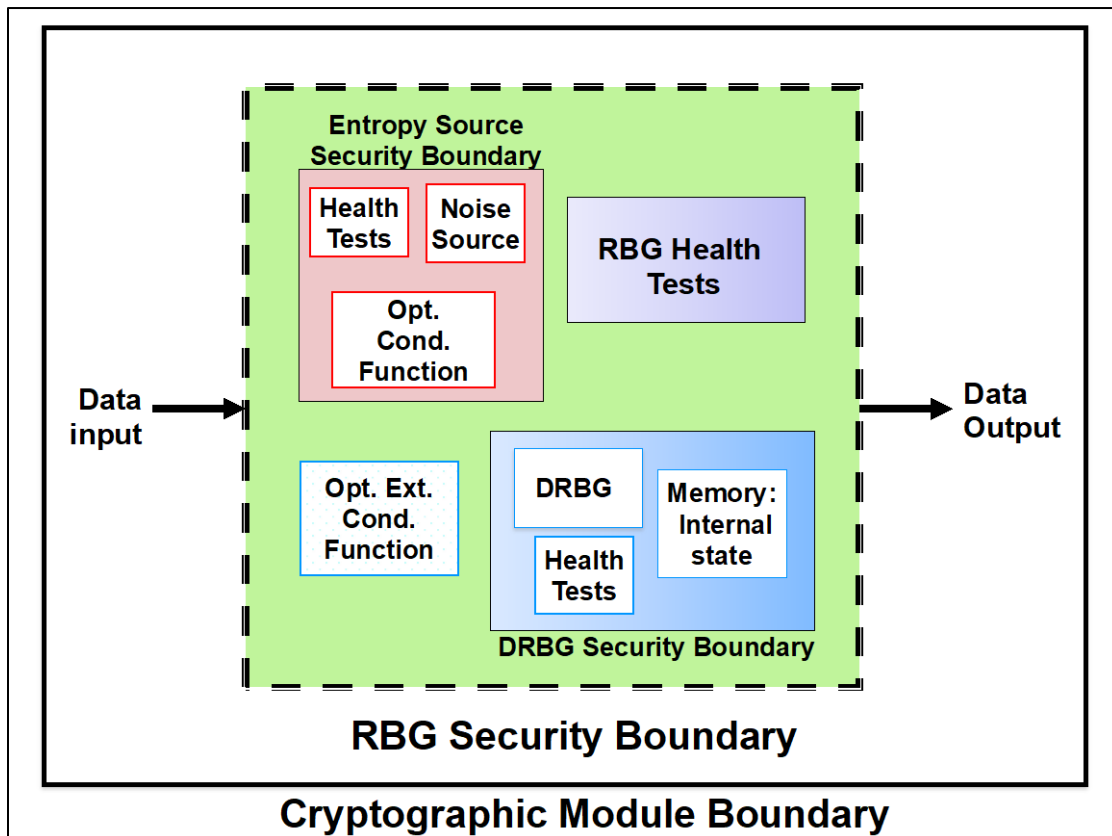


Fig. 2. Example of an RBG Security Boundary within a Cryptographic Module

Note that in the case of the RBG1 construction in [Section 4](#), the security boundary containing the DRBG does not include a randomness source (shown as an entropy source in [Figure 2](#)).

A cryptographic primitive (e.g., an **approved** hash function) used by an RBG may be used by other applications within the same cryptographic module. However, these other applications **shall not** modify or reveal the RBG's output, intermediate values, or internal state.

2.6. Assumptions and Assertions

The RBG constructions in SP 800-90C are based on the use of validated entropy sources and the following assumptions and assertions for properly functioning entropy sources:

1. An entropy source is independent of another entropy source if a) their security boundaries do not overlap (e.g., they reside in separate cryptographic modules, or one is a physical entropy source and the other is a non-physical entropy source), b) there are no common noise sources,⁴ and c) statistical tests provide evidence of the independence of the entropy sources.
2. The use of both validated and non-validated entropy sources is permitted in an implementation, but only entropy sources that have been validated for compliance with [SP800-90B] are used to provide the randomness input for seeding and reseeding a DRBG or providing entropy for an RBG3 construction.

The following assumptions and assertions pertain to the use of validated entropy sources for providing entropy bits:

3. For the purpose of analysis, it is assumed that a) the number of bits that are output by an entropy source is never more than 2^{64} , and b) the number of output bits from the RBG is never more than 2^{64} bits for a DRBG instantiation. In the case of an RBG1 construction with one or more subordinate DRBGs, the output limit applies to the total output provided by the RBG1 construction and all of its subordinate DRBGs.
4. Each entropy-source output has a fixed length, ES_len (in bits).
5. Each entropy-source output is assumed to contain a fixed amount of entropy, denoted as $ES_entropy$, that was assessed during entropy-source implementation validation. (See [SP800-90B] for entropy estimation.) $ES_entropy$ is assumed to be at least 0.1 bits per bit of output.
6. Each entropy source has been characterized as either a physical entropy source or a non-physical entropy source upon successful validation.
7. The outputs from a single entropy source can be concatenated. The entropy of the resultant bitstring is the sum of the entropy from each entropy-source output. For example, if m outputs are concatenated, then the length of the bitstring is $m \times ES_len$ bits, and the entropy for that bitstring is assumed to be $m \times ES_entropy$ bits. (This is a consequence of the model of entropy used in [SP800-90B].)
8. The output of multiple independent entropy sources can be concatenated in an RBG. The entropy in the resultant bitstring is the sum of the entropy in the output of each independent entropy-source output that is considered to be contributing to the entropy in the bitstring (see Methods 1 and 2 in Section 2.3). For example, suppose that the output from independent physical entropy sources A and B and non-physical entropy source C are concatenated. The length of the concatenated bitstring is the sum of the lengths of the component bitstrings (i.e., $ES_len_A + ES_len_B + ES_len_C$).

⁴ They may, however, use the same *type* of noise source (e.g., both entropy sources could use ring oscillators but not the same ones).

- Using Method 1 in [Section 2.3](#), the amount of entropy in the concatenated bitstring is $ES_entropy_A + ES_entropy_B$.
 - Using Method 2 in [Section 2.3](#), the amount of entropy in the concatenated bitstring is the sum of the entropies in the bitstrings (i.e., $ES_entropy_A + ES_entropy_B + ES_entropy_C$).
9. Under certain conditions, the output of one or more entropy sources can be externally conditioned to provide full-entropy output. See [Section 3.3.2](#) and [Section 6.3.1](#) for the use of this assumption and [\[NISTIR8427\]](#) for rationale.

Furthermore,

10. The amount of entropy in a subset bitstring that is “extracted” from the output block of an approved hash function or block cipher is a proportion of the entropy in that block, such that

$$entropy_{subset} = \left(\frac{subset_len}{output_len} \right) entropy_{output_block}$$

where $subset_len$ is the length of the subset bitstring, $output_len$ is the length of the output block, $entropy_{output_block}$ is the amount of entropy in the output block, and $entropy_{subset}$ is the amount of entropy in the subset bitstring.

11. Full entropy bits can be extracted from the output block of a hash function or block cipher when the amount of fresh entropy inserted into the algorithm exceeds the number of bits to be extracted by at least 64 bits. For example, if $output_len$ is the length of the output block, all bits of the output block can be assumed to have full entropy if at least $output_len + 64$ bits of entropy are inserted into the algorithm. As another example, if a DRBG is reseeded at its security strength s , $(s - 64)$ bits with full entropy can be extracted from the DRBG’s output block.
12. To instantiate a DRBG at a security strength of s bits, a bitstring of at least $3s/2$ bits long is needed from a randomness source for an RBG1 construction, and a bitstring with at least $3s/2$ bits of entropy is needed from an entropy source for an RBG2 or RBG3 construction.
13. One or more of the constructions provided herein are used in the design of an RBG.
14. All components of an RBG2 and RBG3 construction (as specified in [Sections 5](#) and [6](#)) reside within the physical boundary of a single [\[FIPS140\]](#)-validated cryptographic module.
15. The DRBGs specified in [\[SP800-90A\]](#) are assumed to meet their explicit security claims (e.g., backtracking resistance, prediction resistance, claimed security strength, etc.).

The following assumptions and assertions have been made for the subordinate DRBGs (sub-DRBGs) that are seeded (i.e., initialized) using an RBG1 construction:

16. A sub-DRBG is considered to be part of the RBG1 construction that initializes it.
17. The assumptions and assertions in items 3, 10, and 14 (above) apply to sub-DRBGs.

2.7. General Implementation and Use Requirements and Recommendations

When implementing the RBGs specified in this Recommendation, an implementation:

1. **Shall** destroy intermediate values before exiting the function or routine in which they are used,
2. **Shall** employ an “atomic” generate operation whereby a generate request is completed before using any of the requested bits,
3. **Should** consider the threats posed by quantum computers in the future, and
4. **Should** be implemented with the capability to support a security strength of 256 bits or to provide full-entropy output.

When using RBGs, the user or application requesting the generation of random or pseudorandom bits **should** request only the number of bits required for a specific immediate purpose rather than generating bits to be stored for future use. Since, in most cases, the bits are intended to be secret, the stored bits (if not properly protected) are potentially vulnerable to exposure, thus defeating the requirement for secrecy.

2.8. General Function Calls

Functions used within this document for accessing the DRBGs in [\[SP800-90A\]](#), the entropy sources in [\[SP800-90B\]](#), and the RBG3 constructions specified in SP 800-90C are provided below. Each function **shall** return a status code that **shall** be checked (e.g., a status of success or failure by the function).

If the status code indicates a success, then additional information may also be returned, such as a state handle from an instantiate function or the bits that were requested to be generated during a generate function.

If the status code indicates a failure of an RBG component, then see [Section 7.1.2](#) for error-handling guidance. Note that if the status code does not indicate a success, an invalid output (e.g., a null bitstring) **shall** be returned with the status code if information other than the status code could be returned.

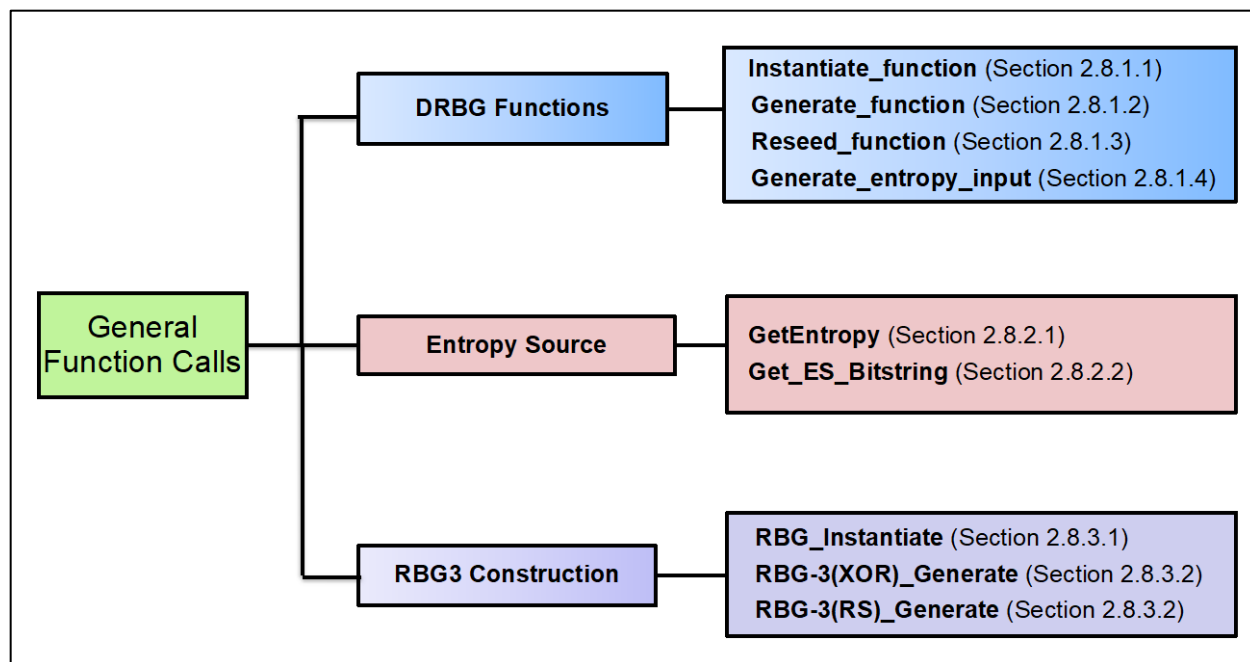


Fig. 3. General Function Calls

2.8.1. DRBG Functions

SP 800-90A specifies several functions for use within a DRBG, indicating the input and output parameters and other implementation details. Note that, in some cases, some input parameters may be omitted, and some output information may not be returned.

At least two functions are required in a DRBG:

1. An instantiate function that seeds the DRBG using the output of a randomness source and other input (see [Section 2.8.1.1](#)) and
2. A generate function that produces output for use by a consuming application (see [Section 2.8.1.2](#)).

A DRBG may also support a reseed function (see [Section 2.8.1.3](#)). A **Get_randomness-source_input** function is used in SP 800-90A to request output from a randomness source during instantiation and reseeding (see [Section 2.8.1.4](#)).

The use of the **Uninstantiate_function** specified in SP 800-90A is not explicitly discussed in SP 800-90C but may be required by an implementation.

2.8.1.1. DRBG Instantiation

A DRBG **shall** be instantiated prior to the generation of pseudorandom bits at the highest security strength to be supported by the DRBG instantiation using the following call:

(status, state_handle) = Instantiate_function(requested_instantiation_security_strength, prediction_resistance_flag, personalization_string).

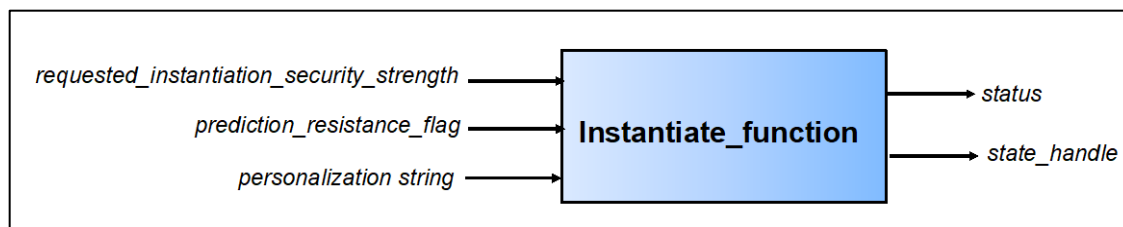


Fig. 4. Instantiate_function

The **Instantiate_function** (shown in Figure 4) is used to instantiate a DRBG at the *requested_instantiation_security_strength* using the output of a randomness source⁵ and an optional *personalization_string* to create seed material. A *prediction_resistance_flag* may be used to indicate whether subsequent **Generate_function** calls may request prediction resistance. As stated in [Section 2.4.1](#), a *personalization_string* is optional but strongly recommended. (Details about the **Instantiate_function** are provided in [\[SP800-90A\]](#).)

If the returned status code for the **Instantiate_function** indicates a success (i.e., the DRBG has been instantiated at the requested security strength), a state handle may⁶ be returned to indicate the particular DRBG instance. When provided, the state handle will be used in subsequent calls to the DRBG (e.g., during a **Generate_function** call) to identify the internal state information for the instantiation. The information in the internal state includes the security strength of the instantiation, the number of times that the instantiation has produced output, and other information that changes during DRBG execution (see [\[SP800-90A\]](#) for each DRBG design).

When the DRBG has been instantiated at the *requested_instantiation_security_strength*, the DRBG will operate at that security strength even if the *requested_security_strength* in subsequent **Generate_function** calls (see [Section 2.8.1.2](#)) is less than the instantiated security strength.

If the *status* code indicates an error and an implementation is designed to return a state handle, an invalid (e.g., *Null*) state handle **shall** be returned.

2.8.1.2. DRBG Generation Request

Pseudorandom bits are generated after DRBG instantiation using the following call:

(*status*, *returned_bits*) = **Generate_function**(*state_handle*, *requested_number_of_bits*,
requested_security_strength, *prediction_resistance_request*, *additional_input*).

⁵ The randomness source provides the randomness input required to instantiate the security strength of the DRBG.

⁶ In cases where only one instantiation of a DRBG will ever exist, a state handle need not be returned since only one internal state will be created.

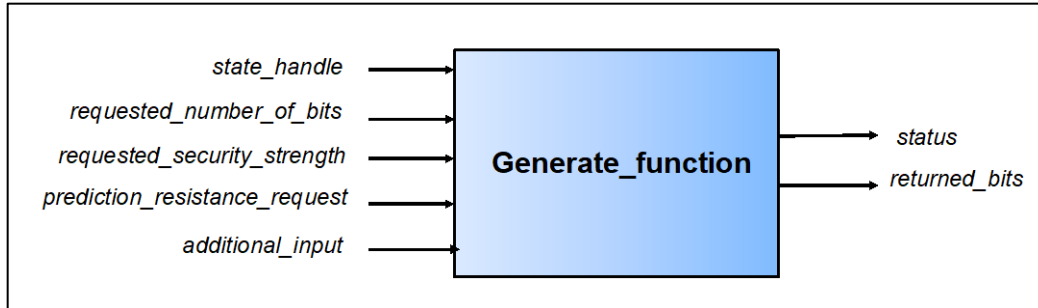


Fig. 5. Generate_function

The **Generate_function** (shown in Figure 5) requests that a DRBG generate a specified number of bits. The request may indicate the DRBG instance to be used (using the state handle returned by an **Instantiate_function** call; see [Section 2.8.1.1](#)), the number of bits to be returned, the security strength that the DRBG needs to support for generating the bitstring, and whether or not prediction resistance is to be obtained during this execution of the **Generate_function**. Optional additional input may also be incorporated into the function call. As stated in [Section 2.4.1](#), the ability to handle and use additional input is recommended.

The **Generate_function** returns status information – either an indication of success or an error. If the returned *status* code indicates a success, the requested number of bits is returned.

- If *requested_number_of_bits* is equal to or greater than the instantiated security strength, the security strength that the *returned_bits* can support (if used as a key) is:

$$ss_key = \text{the instantiated security strength,}$$

where *ss_key* is the security strength of the key.

- If the *requested_number_of_bits* is less than the instantiated security strength, and the *returned_bits* are to be used as a key, the key is capable of supporting a security strength of:

$$ss_key = \text{requested_number_of_bits.}$$

If the status code indicates an error, the *returned_bits* **shall** consist of an invalid (e.g., *Null*) bitstring that **must not** be used. Examples of conditions in which an error indication **shall** be returned include the following:

- The *requested_security_strength* exceeds the instantiated security strength for the DRBG (i.e., the security strength recorded in the DRBG’s internal state during instantiation).
- Prediction resistance has been requested but cannot be obtained at this time.

Details about the **Generate_function** are provided in Section 9.3 of [\[SP800-90A\]](#).

2.8.1.3. DRBG Reseed Request

The reseeding of a DRBG instantiation is intended to insert additional entropy into that DRBG instantiation (e.g., to recover from a possible compromise or to provide prediction resistance). This is accomplished using the following call (note that this does not increase the security strength of the DRBG):

status = **Reseed_function**(*state_handle*, *additional_input*).

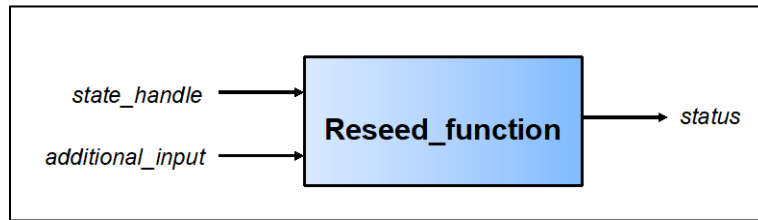


Fig. 6. Reseed_function

A **Reseed_function** (shown in Figure 6) is used to acquire at least s bits of fresh entropy for the DRBG instance indicated by the state handle (or the only instance if no state handle has been provided), where s is the security strength of the DRBG.⁷ In addition to the randomness input provided from the randomness source(s) during reseeding, optional additional input may be incorporated into the reseed process. As discussed in [Section 2.4.1](#), the capability for handling and using additional input is recommended. (Details about the **Reseed_function** are provided in [\[SP800-90A\]](#).)

An indication of the *status* is returned.

The **Reseed_function** is not permitted in an RBG1 construction (see [Section 4](#)) but is permitted in the RBG2 and RBG3 constructions (see [Sections 5](#) and [6](#), respectively).

2.8.1.4. The Get_randomness-source_input Call

A **Get_randomness-source_input** call is used in the **Instantiate_function** and **Reseed_function** in [\[SP800-90A\]](#) to indicate when a randomness source (i.e., an entropy source or RBG) needs to be accessed to obtain randomness input. Details are not provided in SP 800-90A about how the **Get_randomness-source_input** call needs to be implemented. SP 800-90C provides guidance on how the call should actually be implemented based on various situations. [Sections 4](#), [5](#), and [6](#) provide instructions for obtaining input from a randomness source when the **Get_randomness-source_input** call is encountered in SP 800-90A.⁸

2.8.2. Interfacing with Entropy Sources Using the GetEntropy and Get_ES_Bitstring Functions

2.8.2.1. The GetEntropy Call

An entropy source, as discussed in [\[SP800-90B\]](#), is a mechanism for producing bitstrings that cannot be predicted and whose unpredictability can be quantified in terms of min-entropy. SP 800-90B uses the following call for accessing an entropy source:

$(status, ES_output) = \text{GetEntropy}(bits_of_entropy),$

⁷ The value of s is available in the DRBG's internal state.

⁸ Note that, at this time, modifications to the **Instantiate_function** and **Reseed_function** specification in SP 800-90A and to the appropriate algorithms in Section 10 of that document may be required to accommodate the specific requests for entropy for each RBG construction.

where *bits_of_entropy* is the amount of entropy requested, *ES_output* is a bitstring containing the requested amount of entropy, and *status* indicates whether or not the request has been satisfied. See Figure 7.

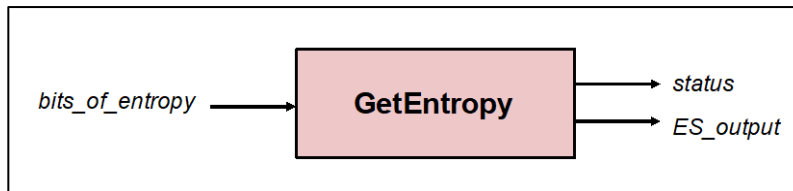


Fig. 7. GetEntropy function

If the *status* indicates a success, a bitstring of at least *bits_of_entropy* long is returned as the *ES_output*. *ES_output* **must** contain at least the requested amount of entropy indicated by the *bits_of_entropy* input parameter. If the *status* does not indicate a success, an invalid *ES_output* bitstring is returned (e.g., *ES_output* could be a null bitstring).

2.8.2.2. The Get_ES_Bitstring Function

A single **GetEntropy** call may not be sufficient to obtain the entropy required for seeding and reseeding a DRBG and for providing input for the exclusive-or operation in an RBG3(XOR) construction (see [Section 6.2](#)). Therefore, SP 800-90C uses a **Get_ES_Bitstring** function (see [Figure 8](#)) to obtain the required entropy from one or more **GetEntropy** calls. The **Get_ES_Bitstring** function is invoked as follows:

$(status, entropy_bitstring) = \text{Get_ES_Bitstring}(bits_of_entropy),$

where *bits_of_entropy* is the amount of entropy requested in the returned *entropy_bitstring*, and *status* indicates whether or not the request has been satisfied.

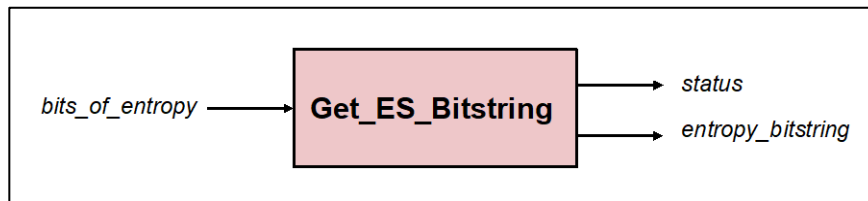


Fig. 8. Get_ES_Bitstring function

Note that if non-validated entropy sources are used (e.g., to provide entropy to be used as additional input), they **shall** be accessed using a different function than is used to access validated entropy sources (i.e., the **Get_ES_Bitstring** function).

If the returned *status* from the **Get_ES_Bitstring** function indicates a success, the requested amount of entropy (i.e., indicated by *bits_of_entropy*) **shall** be returned in the *entropy_bitstring*, whose length is equal to or greater than *bits_of_entropy*. If the *status* does not indicate a success, an invalid *entropy_bitstring* **shall** be returned (e.g., *entropy_bitstring* is a null bitstring).

The **Get_ES_Bitstring** function will be used in this document to access validated entropy sources to obtain one or more bitstrings with entropy using **GetEntropy** calls.

See [Section 3.1](#) for additional discussion about the **Get_ES_Bitstring** function.

2.8.3. Interfacing with an RBG3 Construction

An RBG3 construction requires interface functions to instantiate its DRBG (see [Section 2.8.3.1](#)) and to request the generation of full-entropy bits (see [Section 2.8.3.2](#)).

2.8.3.1. Instantiating a DRBG within an RBG3 Construction

The **RBG3_DRBG_Instantiate** function is used to instantiate the DRBG within the RBG3 construction using the following call:

```
(status, state_handle) = RBG3_DRBG_Instantiate(prediction_resistance_flag,
                                              personalization_string).
```

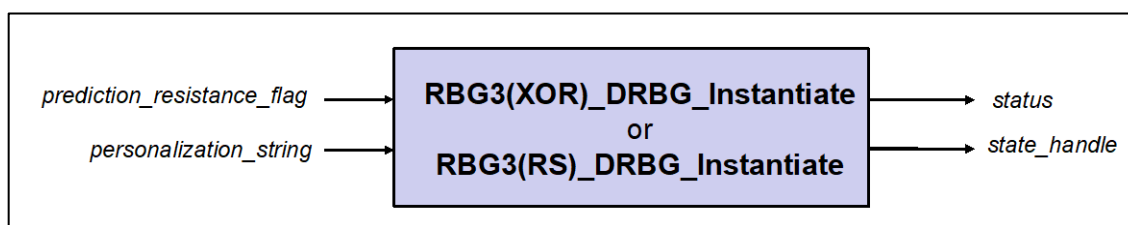


Fig. 9. RBG3 DRBG_Instantiate function

The RBG3's instantiate function (shown in Figure 9) will result in a call to the DRBG's **Instantiate_function** (provided in [Section 2.8.1.1](#)). An optional but recommended *personalization_string* (see [Section 2.4.1](#)) may be provided as an input parameter. If included, the *personalization_string* **shall** be passed to the DRBG that is instantiated in the **Instantiate_function** request. See Sections [6.2.1.1](#) and [6.3.1.1](#) for more specificity.

If the returned *status* code indicates a success, a state handle may be returned to indicate the particular DRBG instance that is to be used by the construction. Note that if multiple instances of the DRBG are used, a separate state handle **shall** be returned for each instance. When provided, the state handle **shall** be used in subsequent calls to that RBG (e.g., during a call to the generate function) when multiple instances of the DRBG have been instantiated. If the status code indicates an error (e.g., entropy is not currently available, or the entropy source has failed), an invalid (e.g., *Null*) state handle **shall** be returned.

2.8.3.2. Generation Using an RBG3 Construction

The RBG3(XOR) and RBG3(RS) generate functions are different because of the difference in their designs (see Sections [6.2.1.2](#) and [6.3.1.2](#)).

For the RBG3(XOR) construction, the generate function is invoked using the following call:

```
(status, returned_bits) = RBG3(XOR)_Generate(state_handle, requested_number_of_bits,
                                              prediction_resistance_request, additional_input).
```

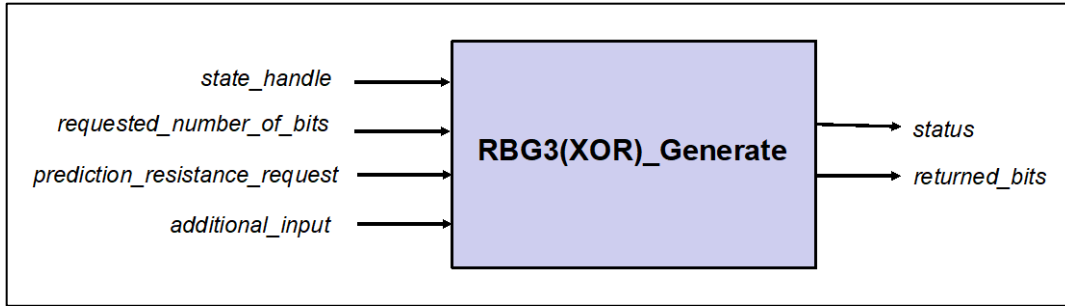



Fig. 10. RBG3(XOR)_Generate function

For the RBG3(RS) construction, the generate function is invoked using the following call:

$(status, returned_bits) = \text{RBG3(RS)_Generate}(state_handle, requested_number_of_bits, additional_input).$

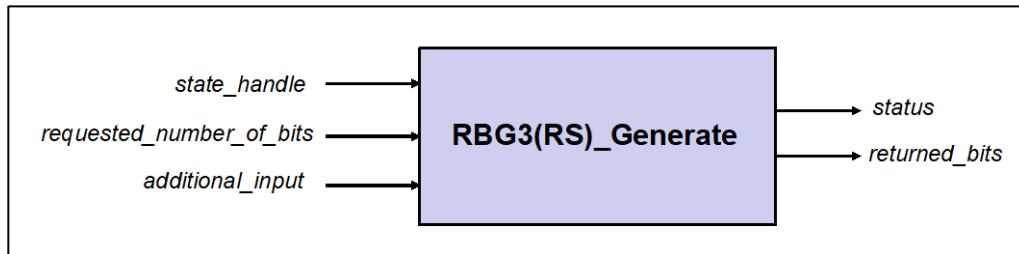


Fig. 11. RBG3(RS)_Generate function

The **RBG3(XOR)_Generate** function (shown in [Figure 10](#)) includes a *prediction_resistance_request* parameter to request a reseed of the RBG3(XOR)'s DRBG instantiation, when desired. This parameter is not included as a parameter for the **RBG3(RS)_Generate** function (shown in [Figure 11](#)) since this design always reseeds itself during execution.

The generate functions result in calls to the entropy sources and the DRBG instantiation used by the RBG3 construction. This call accesses the DRBG using the **Generate_function** call provided in [Section 2.8.1.2](#). The input parameters to the two generate functions are used when calling the DRBG instantiation used by that RBG3 construction.

If the returned status code indicates a success, a bitstring that contains the newly generated bits is returned. The RBG then uses the resulting bitstring as specified for each RBG3 construction (see [Section 6](#)).

If the status code indicates an error (e.g., the entropy source has failed), an invalid (e.g., *Null*) bitstring **shall** be returned as the *returned_bits*.

3. Accessing Entropy Source Output

The security provided by an RBG is based on the use of validated entropy sources. [Section 3.1](#) discusses the use of the **Get_ES_Bitstring** function to request entropy from one or more entropy sources. [Section 3.2](#) discusses the behavior required by an entropy source. [Section 3.3](#) discusses the conditioning of the output of one or more entropy sources to obtain a bitstring with full entropy before further use by an RBG.

3.1. The Get_ES_Bitstring Function

The **Get_ES_Bitstring** function specified in [Section 2.8.2.2](#) is used within an RBG to obtain entropy from one or more validated entropy sources using one or more **GetEntropy** calls (see [Sections 2.8.2.1](#) and [3.2](#)) in whatever manner is required (e.g., by polling the entropy sources or by extracting bits containing entropy from a pool of collected bits). The **Get_ES_Bitstring** function **shall** only be used to access validated entropy sources to obtain the entropy for seeding and reseeding a DRBG and for providing input for the exclusive-or operation of an RBG3(XOR) construction (see [Section 6.2](#)).

In many cases, the **Get_ES_Bitstring** function will need to query an entropy source (or a set of entropy sources) multiple times to obtain the amount of entropy requested. For the most part, the construction of the **Get_ES_Bitstring** function itself is not specified in this document but is left to the developer to implement appropriately for the selected entropy sources.

The behavior of the **Get_ES_Bitstring** function **shall** be as follows:

1. A **Get_ES_Bitstring** function **shall** only be used to access one or more validated entropy sources.
2. The entropy bitstrings produced from multiple entropy-source calls to a single validated entropy source or by calls to multiple validated entropy sources **shall** be concatenated into a single bitstring. The entropy in the bitstring is computed as the sum of the entropy produced by each call to a validated entropy source that is to be counted as contributing entropy to the bitstring (see [Section 2.3](#)).⁹
3. If a failure is reported during an invocation of the **Get_ES_Bitstring** function by any physical or non-physical entropy source whose entropy is counted toward fulfilling an entropy request, the failure **shall** be handled as discussed in [Section 7.1.2](#).
4. If a non-physical entropy source whose entropy is not counted reports a failure, the failure **shall** be reported to the RBG or the consuming application.
5. The **Get_ES_Bitstring** function **shall** not return an *entropy_bitstring* unless the bitstring contains sufficient entropy to fulfill the entropy request. The returned *status* **shall** indicate a success only when this condition is met.

⁹ For Method 1 in [Section 3.3](#), only entropy contributed by one or more validated physical entropy sources is counted. For Method 2, the entropy from all validated entropy sources is counted.

3.2. Entropy Source Requirements

This Recommendation requires the use of one or more validated entropy sources to provide entropy for seeding and reseeding a DRBG and for input to the XOR operation in the RBG3(XOR) construction specified in [Section 6.2](#). In addition to the assumptions and assertions concerning entropy sources in [Section 2.6](#), the following conditions **shall** be met when using these entropy sources:

1. Only validated entropy sources **shall** be used to provide the entropy bitstring for seeding and reseeding a DRBG and for providing input to the XOR operation in the RBG3(XOR) construction.

Non-validated entropy sources may be used by an RBG to provide input for personalization strings and/or the additional input in DRBG function calls (see [Section 2.4.1](#)).

2. Each validated entropy source **shall** be independent of all other validated or non-validated entropy sources used by the RBG.
3. The outputs from an entropy source **shall not** be reused (e.g., the value in the entropy source is erased after being output).
4. When queried for entropy, the validated entropy sources **must** respond as follows:
 - a. The requested output **must** be returned only if the returned status indicates a success. In this case, the *ES-output* bitstring **must** contain the requested amount of entropy. (Note that the *ES-output* bitstring may be longer than the amount of entropy requested, i.e., the bitstring may not have full entropy.)
 - b. If an indication of a failure is returned by a validated entropy source as the status, an invalid (e.g., *Null*) bitstring **shall** be returned as *ES_output*.
5. If the validated entropy-source components operate continuously regardless of whether requests are received and a failure is determined, the entropy source **shall** immediately report the failure to the RBG (see [Section 7.1.2](#)).
6. If a validated entropy source reports a failure (e.g., because of a failed health test), the entropy source **shall not** produce output (except possibly for a failure status indication) until the failure is corrected. The entropy source **shall** immediately report the failure to the **Get_ES_Bitstring** function (see [Section 3.1](#)). If multiple validated entropy sources are used, the report **shall** identify the entropy source that reported the failure.
7. A detected failure of any entropy source **shall** cause the RBG to report the failure to the consuming application and terminate the RBG operation. The RBG **must not** be returned to normal operation until the conditions that caused the failure have been corrected and tested for successful operation.

3.3. External Conditioning to Obtain Full-Entropy Bitstrings

An RBG3(XOR) construction (see [Section 6.2](#)) and a CTR_DRBG without a derivation function in an RBG2 or RBG3 construction (see Sections [5](#) and [6](#)) require bitstrings with full entropy from an entropy source. If the validated entropy source does not provide full-entropy output, a method

for conditioning the output to obtain a bitstring with full entropy is needed. Since this conditioning is performed outside an entropy source, the output is said to be *externally conditioned*.

When external conditioning is performed, the vetted conditioning function listed in [SP800-90B] **shall** be used.

3.3.1. Conditioning Function Calls

The conditioning functions operate on bitstrings obtained from one or more calls to the entropy source(s).

The following format is used in [Section 3.3.2](#) for a conditioning-function call:

conditioned_output = **Conditioning_function**(*input_parameters*),

where the *input_parameters* for the selected conditioning function are discussed in [Sections 3.3.1.2](#) and [3.3.1.3](#), and *conditioned_output* is the output returned by the conditioning function.

3.3.1.1. Keys Used in External Conditioning Functions

The **HMAC**, **CMAC**, and **CBC-MAC** vetted conditioning functions require the input of a *Key* of a specific length (*keylen*). Unlike other cryptographic applications, keys used in these external conditioning functions do not require secrecy to accomplish their purpose so may be hard-coded, fixed, or all zeros.

For the **CMAC** and **CBC-MAC** conditioning functions, the length of the key **shall** be an **approved** key length for the block cipher used (e.g., *keylen* = 128, 192, or 256 bits for AES).

For the **HMAC** conditioning function, the length of the key **shall** be equal to the length of the hash function's output block (i.e., *output_len*).

Table 2. Key Lengths for the Hash-based Conditioning Functions

Hash Function	Length of the output block (<i>output_len</i>) and key (<i>keylen</i>)
SHA-224, SHA-512/224, SHA3-224	224
SHA-256, SHA-512/256, SHA3-256	256
SHA-384, SHA3-384	384
SHA-512, SHA3-512	512

Using random keys may provide some additional security in case the input is more predictable than expected. Thus, these keys **should** be chosen randomly in some way (e.g., by drawing bits directly from the entropy source and inserting them into the key or by providing entropy-source bits to a conditioning function with a fixed key to derive the new key). Note that any entropy used to randomize the key **shall not** be used for any other purpose (e.g., as input to the conditioning function).

3.3.1.2. Hash Function-based Conditioning Functions

Conditioning functions may be based on **approved** hash functions.

One of the following calls **shall** be used for external conditioning when the conditioning function is based on a hash function:

1. Using an **approved** hash function directly:

$\text{conditioned_output} = \text{Hash}(\text{entropy_bitstring}),$

where the hash function operates on the *entropy_bitstring* provided as input.

2. Using HMAC with an **approved** hash function:

$\text{conditioned_output} = \text{HMAC}(\text{Key}, \text{entropy_bitstring}),$

where HMAC operates on the *entropy_bitstring* using a *Key* determined as specified in [Section 3.3.1.1](#).

3. Using Hash_df as specified in SP 800-90A:

$\text{conditioned_output} = \text{Hash_df}(\text{entropy_bitstring}, \text{output_len}),$

where the derivation function operates on the *entropy_bitstring* provided as input to produce a bitstring of *output_len* bits.

In all three cases, the length of the conditioned output is equal to the length of the output block of the selected hash function (i.e., *output_len*).

3.3.1.3. Block Cipher-based Conditioning Functions

Conditioning functions may be based on **approved** block ciphers.¹⁰ TDEA **shall not** be used as the block cipher (see [Section 2.6](#)).

For block cipher-based conditioning functions, one of the following calls **shall** be used for external conditioning:

1. Using CMAC (as specified in [\[SP800-38B\]](#)) with an **approved** block cipher:

$\text{conditioned_output} = \text{CMAC}(\text{Key}, \text{entropy_bitstring}),$

where CMAC operates on the *entropy_bitstring* using a *Key* determined as specified in [Section 3.3.1.1](#).

2. Using CBC-MAC (specified in Appendix F of [\[SP800-90B\]](#)) with an **approved** block cipher:

$\text{conditioned_output} = \text{CBC-MAC}(\text{Key}, \text{entropy_bitstring}),$

where CBC-MAC operates on the *entropy_bitstring* using a *Key* determined as specified in [Section 3.3.1.1](#).

¹⁰ At the time of publication, only AES-128, AES-192, and AES-256 were **approved** as block ciphers for the conditioning functions (see SP 800-90B). In all three cases, the block length is 128 bits.

CBC-MAC **shall** only be used as an external conditioning function under the following conditions:

- a. The length of the input is an integer multiple of the block size of the block cipher (e.g., a multiple of 128 bits for AES) – no padding is done by CBC-MAC itself.¹¹
- b. All inputs to CBC-MAC in the same RBG **shall** have the same length.
- c. If the CBC-MAC conditioning function is used to obtain full entropy from an entropy source for CTR_DRBG instantiation or reseeding:
 - A personalization string **shall not** be used during instantiation.
 - Additional input **shall not** be used during the reseeding of the CTR_DRBG but may be used during the generate process.

CBC-MAC is not approved for any use other than in an RBG (see [SP800-90B]).

3. Using the **Block_cipher_df** as specified in [SP800-90A] with an **approved** block cipher:

conditioned_output = **Block_cipher_df**(*entropy_bitstring*, *block_length*),

where **Block_cipher_df** operates on the *entropy_bitstring* using a key specified within the function, and the *block_length* is 128 bits for AES.

In all three cases, the length of the conditioned output is equal to the length of the output block (i.e., 128 bits for AES). If the requested amount of entropy is requested for subsequent use by an RBG,¹² then multiple iterations of the conditioning function may be required, each using a different *entropy_bitstring*.

3.3.2. Using a Vetted Conditioning Function to Obtain Full-Entropy Bitstrings

This construction will produce a bitstring with full entropy using one of the conditioning functions identified in [Section 3.3.1.1](#) for an RBG2 or RBG3 construction whenever a bitstring with full entropy is required (e.g., to seed or reseed a CTR_DRBG with no derivation function or to provide full entropy for the RBG3(XOR) construction). This process is unnecessary if the entropy source provides full-entropy output.

Let *output_len* be the length of the output block of the vetted conditioning function to be used; *output_len* is the length of the hash function's output block when a hash-based conditioning function is used (see [Section 3.3.1.2](#)); *output_len* = 128 when an AES-based conditioning function is used (see [Section 3.3.1.3](#)).

The approach used by this construction is to acquire sufficient entropy from the entropy source to produce *output_len* bits with full entropy in the conditioning function's output block, where *output_len* is the length of the output block. The amount of entropy required for each use of the conditioning function is *output_len* + 64 bits (see item 11 of [Section 2.6](#)). This process is repeated until the requested number of full-entropy bits have been produced.

¹¹ Any padding required could be done before submitting the *entropy_bitstring* to the CBC-MAC function.

¹² Since the output block of AES is only 128 bits, this will often be the case when seeding or reseeding a DRBG.

The **Get_conditioned_full_entropy_input** function below obtains entropy from one or more entropy sources using the **Get_ES_Bitstring** function discussed in [Section 3.1](#) and conditions it to provide an n -bit string with full entropy.

Get_conditioned_full_entropy_input:

Input: integer n . Comment: the requested number of full-entropy bits.

Output: integer $status$, bitstring $returned_bitstring$.

Process:

1. $temp$ = the *Null* string.
2. $ctr = 0$.
3. While $ctr < n$, do
 - 3.1 $(status, entropy_bitstring) = \text{Get_ES_Bitstring}(output_len + 64)$.
 - 3.2 If $(status \neq \text{SUCCESS})$, then return $(status, invalid_bitstring)$.
 - 3.3 $conditioned_output = \text{Conditioning_function}(input_parameters)$.
 - 3.4 $temp = temp \parallel conditioned_output$.
 - 3.5 $ctr = ctr + output_len$.
4. $returned_bitstring = \text{leftmost}(temp, n)$.
5. Return $(\text{SUCCESS}, returned_bitstring)$.

Steps 1 and 2 initialize the temporary bitstring ($temp$) for storing the full-entropy bitstring being assembled and the counter (ctr) that counts the number of full-entropy bits produced for each iteration of step 3.

Step 3 obtains and processes the entropy for each iteration.

- Step 3.1 requests $output_len + 64$ bits from the validated entropy sources. When the output of multiple entropy sources is used, the entropy counted for fulfilling the request for $outlen + 64$ bits is determined using Method 1 or Method 2 as specified in [Section 2.3](#) in the following situations:

Method 1 **shall** be used when:

Instantiating and reseeding an RBG2(P) construction containing a CTR_DRBG with no derivation function (see [Section 5.2.1](#), item 1b, and [Section 5.2.3](#)),

Instantiating and reseeding a CTR_DRBG with no derivation function that is used within an RBG3 construction (see [Section 6.1](#), requirement 1), or

Generating bits in an RBG3(XOR) construction (see [Section 6.2.1.2](#), step 1).

Method 2 **shall** be used when instantiating and reseeding an RBG2(NP) construction containing a CTR_DRBG with no derivation function (see [Section 5.2.1](#), item 1b, and [Section 5.2.3](#)).

- 1087 • Step 3.2 checks whether or not the *status* returned in step 3.1 indicated a success. If the
1088 *status* did not indicate a success, the *status* is returned along with an invalid bitstring as the
1089 *returned_bitstring* (e.g., *invalid_bitstring* is *Null*).
- 1090 • Step 3.3 invokes the conditioning function for processing the *entropy_bitstring* obtained
1091 from step 3.1. The *input_parameters* for the selected **Conditioning_function** are specified
1092 in Sections [3.3.1.2](#) or [3.3.1.3](#), depending on the conditioning function used.
- 1093 • Step 3.4 concatenates the *conditioned_output* received in step 3.3 to the temporary bitstring
1094 (*temp*), and step 3.5 increments the counter for the number of full-entropy bits that have
1095 been produced so far.
- 1096 • If at least n full-entropy bits have not been produced, repeat the process starting at step 3.1.
- 1097 • Step 4 truncates the full-entropy bitstring to n bits.
- 1098 • Step 5 returns an n -bit full-entropy bitstring as the *returned_bitstring*.

1099 **4. RBG1 Constructions Based on RBGs with Physical Entropy Sources**

1100 An RBG1 construction provides a source of cryptographic random bits from a device that has no
1101 internal randomness source. Its security depends entirely on being instantiated securely from an
1102 RBG with access to a physical entropy source that resides outside of the device.

1103 An RBG1 construction is instantiated (i.e., seeded) only once before its first use by an RBG2(P)
1104 construction (see [Section 5](#)) or an RBG3 construction (see [Section 6](#)). Since a randomness source
1105 is not available after DRBG instantiation, an RBG1 construction cannot be reseeded and, therefore,
1106 cannot provide prediction resistance.

1107 An RBG1 construction may be useful for constrained devices in which an entropy source cannot
1108 be implemented or in any device in which access to a suitable source of randomness is not available
1109 after instantiation. Since an RBG1 construction cannot be reseeded, the use of the DRBG is limited
1110 to the DRBG's seedlife (see [\[SP800-90A\]](#)).

1111 Subordinate DRBGs (sub-DRBGs) may be used within the security boundary of an RBG1
1112 construction (see [Section 4.3](#)). The use of one or more sub-DRBGs may be useful for
1113 implementations that use flash memory, such as when the number of write operations to the
1114 memory is limited (resulting in short device lifetimes) or when there is a need to use different
1115 DRBG instantiations for different purposes. The RBG1 construction is the source of the
1116 randomness that is used to (optionally) instantiate one or more sub-DRBGs. Each sub-DRBG is a
1117 DRBG specified in SP 800-90A and is intended to be used for a limited time and a limited purpose.
1118 A sub-DRBG is, in fact, a different instantiation of the DRBG design implemented within the
1119 RBG1 construction (see [Section 2.4.1](#)).

1120 **4.1. RBG1 Description**

1121 As shown in [Figure 12](#), an RBG1 construction consists of a DRBG contained within a DRBG
1122 security boundary in one cryptographic module and an RBG (serving as a randomness source)
1123 contained within a separate cryptographic module from that of the RBG1 construction. Note that
1124 the required health tests are not shown in the figure.

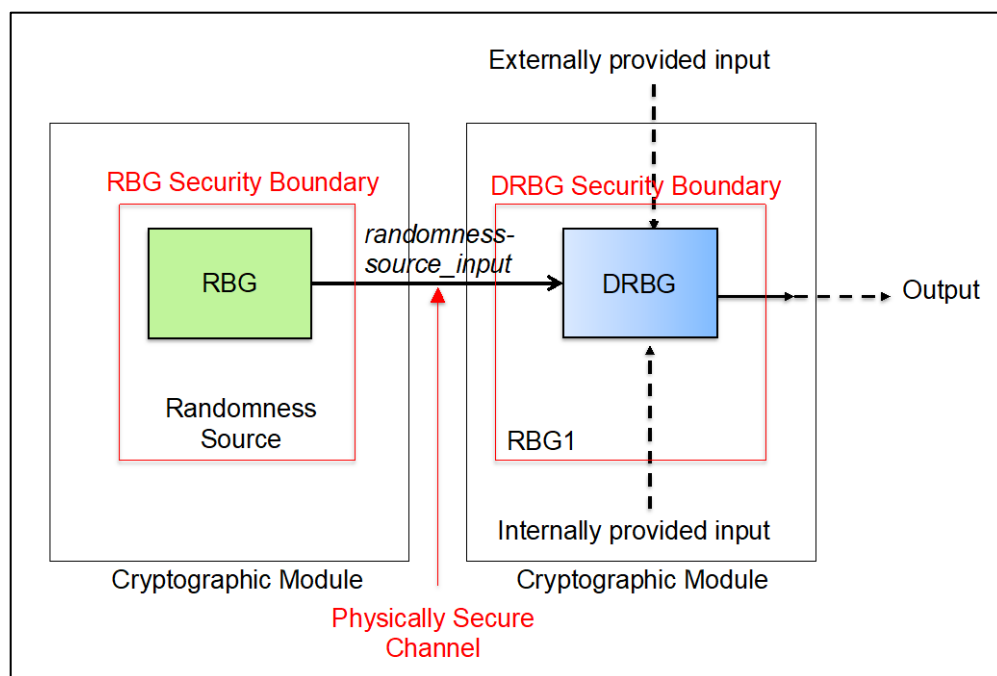


Fig. 12. RBG1 Construction

The RBG for instantiating the DRBG within the RBG1 construction **must** be either an RBG2(P) construction that has support for prediction resistance requests (see [Section 5](#)) or an RBG3 construction (see [Section 6](#)). A physically secure channel between the randomness source and the DRBG is used to securely transport the randomness input required for the instantiation of the DRBG. An optional recommended personalization string and optional additional input may be provided from within the DRBG’s cryptographic module or from outside of that module (see [Section 2.4.1](#)).

An external conditioning function is not needed for this design because the output of the RBG has already been cryptographically processed.

The output from an RBG1 construction may be used within the cryptographic module (e.g., to seed a sub-DRBG as specified in [Section 4.3](#)) or by an application outside of the RBG1 security boundary.

The security strength provided by the RBG1 construction is the minimum of the security strengths provided by the DRBG within the construction, the secure channel, and the RBG used to seed the DRBG.

Examples of RBG1 and sub-DRBG constructions are provided in Appendices [B.2](#) and [B.3](#), respectively.

4.2. Conceptual Interfaces

Interfaces to the DRBG within an RBG1 construction include function calls for instantiating the DRBG and generating pseudorandom bits upon request (see Sections [4.2.1](#) and [4.2.2](#)).

Note that reseeding is not included in this construction.

4.2.1. Instantiating the DRBG in the RBG1 Construction

The DRBG within the RBG1 construction may be instantiated at any security strength possible for the DRBG design using the **Instantiate_function** discussed in [Section 2.8.1.1](#) and [\[SP800-90A\]](#), subject to the maximum security strength that is supported by the RBG used as the randomness source.

(*status*, *RBG1_state_handle*) =
Instantiate_function (*s*, *prediction_resistance_flag* = FALSE, *personalization_string*),

where *s* is the requested security strength for the DRBG in the RBG1 construction. If used, the *prediction_resistance_flag* is set to FALSE since the DRBG cannot be reseeded to provide prediction resistance.

An external RBG (i.e., the randomness source) **shall** be used to obtain the bitstring necessary for establishing the DRBG's *s*-bit security strength.

In SP 800-90A, the **Instantiate_function** specifies the use of a **Get_randomness-source_input** call to obtain randomness input from the randomness source for instantiation (see [Section 2.8.1.4](#) in this document and in [\[SP800-90A\]](#)). For an RBG1 construction, an **approved** external RBG2(P) or RBG3 construction **must** be used as the randomness source (see [Sections 5](#) and [6](#), respectively).

If the randomness source is an RBG2(P) construction (see [Figure 13](#)), the **Get_randomness-source_input** call in the **Instantiate_function** **shall** be replaced by a **Generate_function** call to the RBG2(P) construction (in whatever manner is required) (see [Sections 2.8.1.2](#) and [5.2.2](#)). The RBG2(P) construction **must** be reseeded using its internal entropy source(s) before generating bits to be provided to the RBG1 construction. This is accomplished by setting the *prediction_resistance_request* parameter in the **Generate_function** call to TRUE (see steps 1a and 2a below).

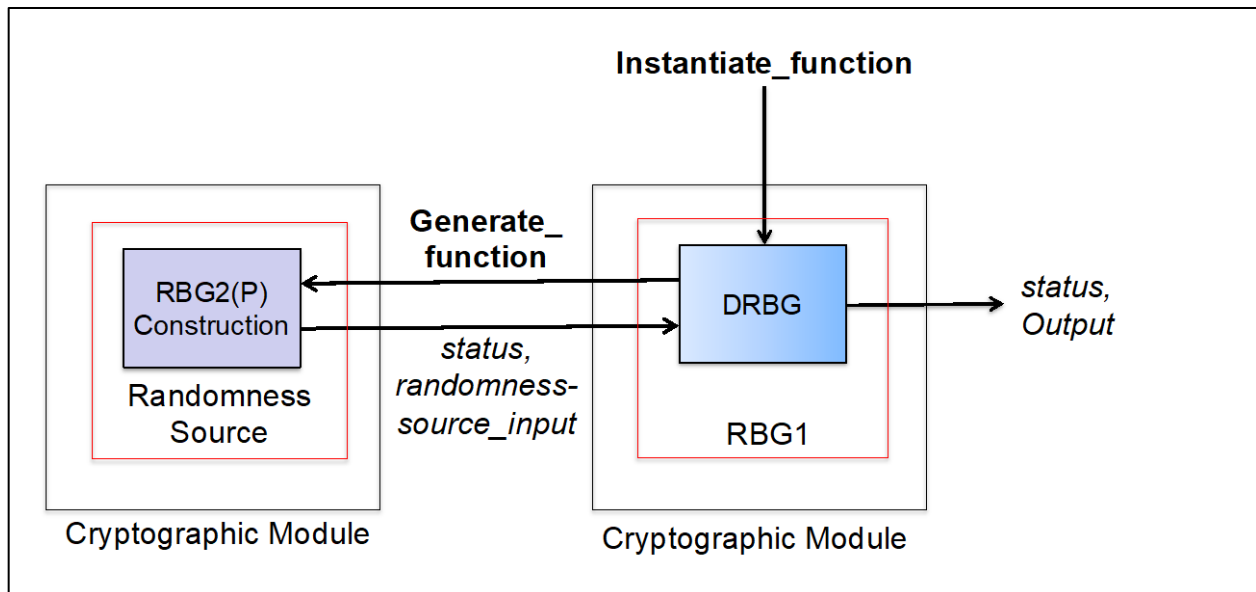


Fig. 13. Instantiation Using an RBG2(P) Construction as a Randomness Source

1173 If the randomness source is an RBG3 construction (as shown in [Figure 14](#)), the **Get_randomness-**
1174 **source_input** call **shall** be replaced by the appropriate RBG3 generate function (see Sections
1175 [2.8.3.2](#), [6.2.1.2](#), and [6.3.1.2](#) and steps 1b, 1c, 2b, and 2c below).

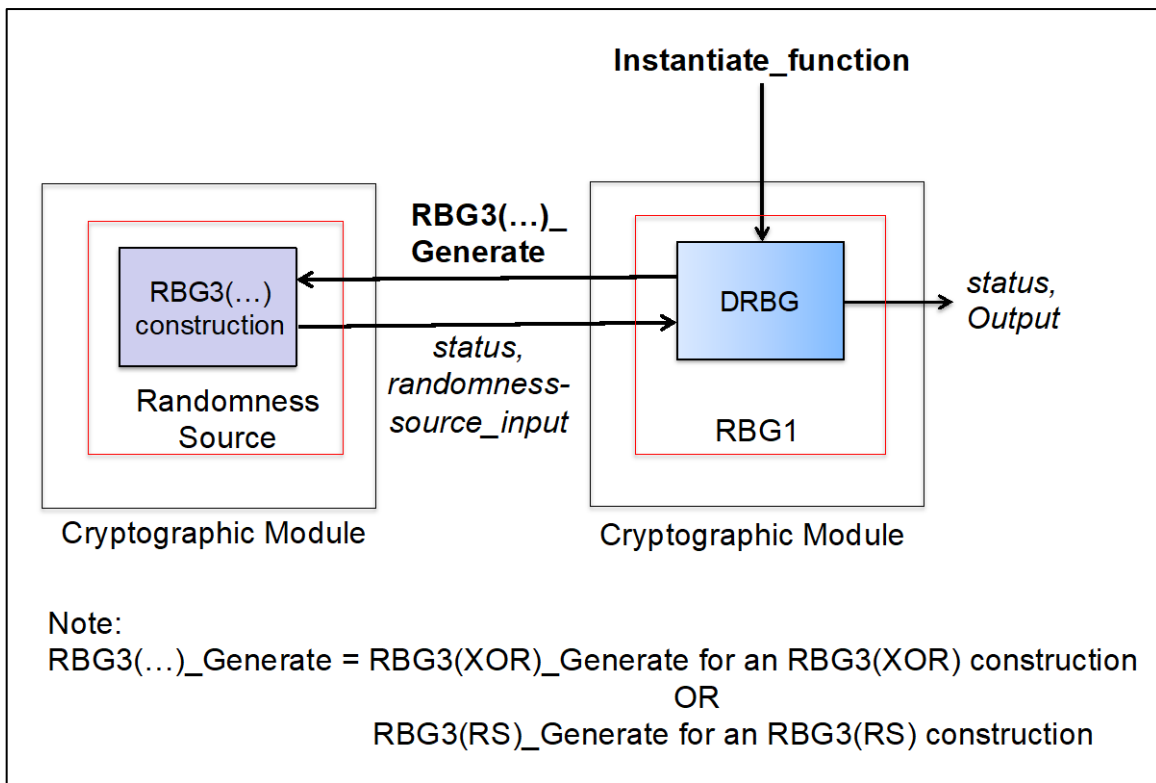


Fig. 14. Instantiation using an RBG3(XOR) or RBG3(RS) Construction as a Randomness Source

Let s be the security strength to be instantiated. The DRBG within an RBG1 construction is instantiated as follows:

1. When an RBG1 construction is instantiating a CTR_DRBG without a derivation function, $s + 128$ bits¹³ **shall** be obtained from the randomness source as follows:

If the randomness source is an RBG2(P) construction (see [Figure 13](#)), the **Get_randomness-source_input** call is replaced by:

$(status, randomness_source_input) = \text{Generate_function}(RBG2_state_handle, s + 128, s, prediction_resistance_request = \text{TRUE}, additional_input)$.

Note that the DRBG within the RBG2(P) construction **must** be reseeded before generating output.¹⁴ This may be accomplished by requesting prediction resistance (i.e., setting $prediction_resistance_request = \text{TRUE}$). See Requirement 17 in [Section 4.4.1](#).

¹³ For AES, the block length is 128 bits, and the key length is equal to the security strength s . SP 800-90A requires the randomness input from the randomness source to be key length + block length bits when a derivation function is not used.

¹⁴ See Requirement 11 in Section 5.4.1.

1190 If the randomness source is an RBG3(XOR) construction (see [Figure 14](#)), the
1191 **Get_randomness-source_input** call is replaced by:

1192 $(status, randomness_source_input) = \mathbf{RBG3(XOR)_Generate}(RBG3_state_handle, s$
1193 $+ 128, prediction_resistance_request, additional_input).$

1194 A request for prediction resistance from the DRBG used by the RBG3(XOR)
1195 construction is optional.

1196 c) If the randomness source is an RBG3(RS) construction (see [Figure 14](#)), the
1197 **Get_randomness-source_input** call is replaced by:

1198 $(status, randomness_source_input) = \mathbf{RBG3(RS)_Generate}(RBG3_state_handle,$
1199 $3s/2, additional_input).$

1200 2. When an RBG1 construction is instantiating any other DRBG (including a CTR_DRBG
1201 with a derivation function), $3s/2$ bits **shall** be obtained from a randomness source that
1202 provides a security strength of at least s bits.

1203 a) If the randomness source is an RBG2(P) construction (see [Figure 13](#)), the
1204 **Get_randomness-source_input** call is replaced by:

1205 $(status, randomness_source_input) = \mathbf{Generate_function}(RBG2_state_handle, 3s/2,$
1206 $s, prediction_resistance_request = \mathbf{TRUE}, additional_input).$

1207 Note that the DRBG within the RBG2(P) construction **must** be reseeded before
1208 generating output. This is accomplished by requesting prediction resistance (i.e., by
1209 setting $prediction_resistance_request = \mathbf{TRUE}$). See Requirement 17 in [Section 4.4](#).

1210 b) If the randomness source is an RBG3(XOR) construction (see [Figure 14](#)), the
1211 **Get_randomness-source_input** call is replaced by:

1212 $(status, randomness_source_input) = \mathbf{RBG3(XOR)_Generate}(RBG3_state_handle,$
1213 $3s/2, prediction_resistance_request, additional_input).$

1214 A request for prediction resistance from the DRBG used by the RBG3(XOR)
1215 construction is optional.

1216 c) If the randomness source is an RBG3(RS) construction (see [Figure 14](#)), the
1217 **Get_randomness-sourceinput** call is replaced by:

1218 $(status, randomness_source_input) = \mathbf{RBG3(RS)_Generate}(RBG3_state_handle,$
1219 $3s/2, additional_input).$

1220 4.2.2. Requesting Pseudorandom Bits

1221 Pseudorandom bits from the RBG1 construction **shall** be requested using the following call:

1222 $(status, returned_bits) = \mathbf{Generate_function}(RBG1_state_handle,$
1223 $requested_number_of_bits, s, prediction_resistance_request = \mathbf{FALSE}, additional_input).$

1224 The $prediction_resistance_request$ is set to **FALSE** or the parameter may be omitted since a
1225 reseeding capability is not included in an RBG1 construction.

4.3. Using an RBG1 Construction with Subordinate DRBGs (Sub-DRBGs)

[Figure 15](#) depicts an example of the use of optional subordinate DRBGs (sub-DRBGs) within the security boundary of an RBG1 construction. The RBG1 construction is used as the randomness source to provide separate outputs to instantiate each of its sub_DRBGs.

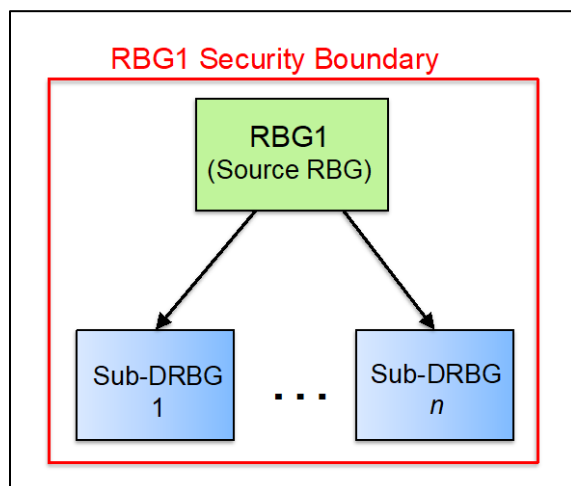


Fig. 15. RBG1 Construction with Sub-DRBGs

The RBG1 construction and each of its sub-DRBGs **shall** be implemented as separate physical or logical entities (see [Figure 15](#)).

- When implemented as separate physical entities, the DRBG algorithms used by the RBG1 construction and a sub-DRBG **shall** be the same DRBG algorithm (e.g., the RBG1 construction and all of its sub_DRBGs use HMAC_DRBG and SHA-256).
- When implemented as separate logical entities, the same software or hardware implementation of a DRBG algorithm is used but with a different internal state for each logical entity (e.g., the RBG1 construction has an internal state whose state handle is *RBG1_state_handle*, while the state handle for Sub-DRBG 1's internal state is *sub-DRBG1_state_handle*).

The sub-DRBGs have the following characteristics:

1. A sub-DRBG cannot be reseeded or provide prediction resistance.
2. Sub-DRBG outputs are considered outputs from the RBG1 construction.
3. The security strength that can be provided by a sub-DRBG is no more than the security strength of its randomness source (i.e., the RBG1 construction).
4. Each sub-DRBG has restrictions on its use (e.g., the number of outputs) as specified for its DRBG algorithm in [\[SP800-90A\]](#).
5. Sub-DRBGs cannot provide output with full entropy.
6. The number of sub-DRBGs that can be instantiated by a RBG1 construction is limited only by practical considerations associated with the implementation or application.

4.3.1. Instantiating a Sub-DRBG

Instantiation of the sub-DRBG is requested (e.g., by a consuming application) using the **Instantiate_function** discussed in [Section 2.8.1.1](#) and [\[SP800-90A\]](#).

(status, sub-DRBG_state_handle) =
Instantiate_function(*s, prediction_resistance_flag = FALSE, personalization_string*),

where *s* is the requested security strength for the (target) sub-DRBG (note that *s* **must** be no greater than the security strength of the RBG1 construction).¹⁵

The (target) sub-DRBG is instantiated as follows:

1. When the sub-DRBG uses CTR_DRBG without a derivation function, *s* + 128 bits¹⁶ **shall** be obtained from the RBG1 construction as follows:

(status, randomness-source_input) = Generate_function(*RBG1_state_handle, s +*
128, s, prediction_resistance_request = FALSE, additional_input).

2. When the sub-DRBG uses any other DRBG (including a CTR_DRBG with a derivation function), *3s/2* bits **shall** be obtained from the RBG1 construction as follows:

(status, randomness-source_input) = Generate_function(*RBG1_state_handle, 3s/2,*
s, prediction_resistance_request = FALSE, additional_input).

4.3.2. Requesting Random Bits

Pseudorandom bits may be requested from a sub-DRBG using the following call (see [Section 2.8.1.2](#)):

(status, returned_bits) = Generate_function(*sub_DRBG_state_handle,*
requested_number_of_bits, requested_security_strength, prediction_resistance_request =
FALSE, additional_input),

where *sub_DRBG_state_handle* (if used) was returned by the **Instantiate_function** (see Sections [2.8.1.1](#) and [4.3.1](#)).

4.4. Requirements

4.4.1. RBG1 Requirements

An RBG1 construction being instantiated has the following testable requirements (i.e., testable by the validation labs):

1. An **approved** DRBG from [\[SP800-90A\]](#) whose components are capable of providing the targeted security strength for the RBG1 construction **shall** be employed.

¹⁵ The implementation is required to check the requested security strength (for the sub-DRBG) against the security strength recorded in the internal state of the RBG1's DRBG (see SP 800-90A).

¹⁶ For AES, the block length is 128 bits, and the key length is equal to the security strength *s*. SP 800-90A requires the randomness input from the randomness source to be (key length + block length) bits when a derivation function is not used.

2. The RBG1 components **shall** be successfully validated for compliance with [\[SP800-90A\]](#), SP 800-90C, [\[FIPS140\]](#), and the specification of any other **approved** algorithm used within the RBG1 construction, as applicable.
 3. The RBG1 construction **shall not** produce any output until it is instantiated.
 4. The RBG1 construction **shall not** include a reseed capability.
 5. The RBG1 construction **shall not** permit itself to be instantiated more than once.¹⁷
 6. For a Hash_DRBG, HMAC_DRBG or CTR_DRBG (with a derivation function), $3s/2$ bits **shall** be obtained from a randomness source (see Requirements 13 - 17), where s is the targeted security strength for the DRBG used in the RBG1 construction.
 7. For a CTR_DRBG (without a derivation function), $s + 128$ bits¹⁸ **shall** be obtained from the randomness source (see Requirements 13 - 17), where s is the targeted security strength for the DRBG used in the RBG1 construction.
 8. The internal state of the RBG1 construction **shall** be maintained¹⁹ and updated to produce output on demand.
 9. The RBG1 construction **shall not** provide output for generating requests that specify a security strength greater than the instantiated security strength of its DRBG.
 10. If the RBG1 construction is used to instantiate a sub-DRBG, the RBG1 construction **may** directly produce output in addition to instantiating the sub-DRBG.
 11. If the seedlife of the DRBG within the RBG1 construction is ever exceeded or a health test of the DRBG fails, the use of the RBG1 construction **shall** be terminated.
 12. If a health test on the RBG1 construction fails, the RBG1 construction and all of its sub-DRBGs **shall** be terminated.
- The non-testable requirements for the RBG1 construction are listed below. If these requirements are not met, no assurance can be obtained about the security of the implementation.
13. An **approved** RBG2(P) construction with support for prediction resistance requests or an RBG3 construction **must** be used as the randomness source for the DRBG in the RBG1 construction.
 14. The randomness source **must** fulfill the requirements in [Section 5](#) (for an RBG(P) construction) or [Section 6](#) (for an RBG3 construction), as appropriate.
 15. The randomness source **must** provide the requested number of bits at a security strength of s bits or higher, where s is the targeted security strength for the RBG1 construction.
 16. The specific output of the randomness source (or portion thereof) that is used for the instantiation of an RBG1 construction **must not** be used for any other purpose, including for seeding a different instantiation.

¹⁷ While technically possible to reseed the DRBG, doing so outside of very controlled conditions (e.g., “in the field”) might result in seeds with less than the required amount of randomness.

¹⁸ Note that $s + 128 = keylen + blocklen = seedlen$, as specified in SP 800-90A.

¹⁹ This means ever-changing but maintained regardless of access to power for its entire lifetime.

17. If an RBG2(P) construction is used as the randomness source for the RBG1 construction, the RBG2(P) construction **must** be reseeded (i.e., prediction resistance must be obtained within the RBG2(P) construction) before generating bits for each RBG1 instantiation.
18. A physically secure channel **must** be used to insert the randomness input from the randomness source into the DRBG of the RBG1 construction.
19. An RBG1 construction **must not** be used for applications that require a higher security strength than has been instantiated.

4.4.2. Sub-DRBG Requirements

A sub-DRBG has the following testable requirements (i.e., testable by the validation labs).

1. The randomness source for a sub-DRBG **shall** be an RBG1 construction; a sub-DRBG **shall not** serve as a randomness source for another sub-DRBG.
2. A sub-DRBG **shall** employ the same DRBG components as its randomness source.
3. A sub-DRBG **shall** reside in the same security boundary as the RBG1 construction that initializes it.
4. The RBG1 construction **shall** fulfill the appropriate requirements of [Section 4.4.1](#).
5. A sub-DRBG **shall** exist only for a limited time and purpose, as determined by the application or developer.
6. The output from the RBG1 construction that is used for sub-DRBG instantiation **shall not** be output from the security boundary of the construction and **shall not** be used for any other purpose, including for seeding a different sub-DRBG.
7. A sub-DRBG **shall not** permit itself to be instantiated more than once.
8. A sub-DRBG **shall not** provide output for use by the RBG1 construction (e.g., as additional input) or another sub-DRBG in the security boundary.
9. The security strength s requested for a target sub-DRBG instantiation **shall not** exceed the security strength that is supported by the RBG1 construction.
10. For a Hash_DRBG, HMAC_DRBG or CTR_DRBG (with a derivation function), $3s/2$ bits **shall** be obtained from the RBG1 construction for instantiation, where s is the requested security strength for the target sub-DRBG.
11. For a CTR_DRBG (without a derivation function), $s + 128$ bits **shall** be obtained from the RBG1 construction for instantiation, where s is the requested security strength for the target sub-DRBG.
12. A sub-DRBG **shall not** produce output until it is instantiated.
13. A sub-DRBG **shall not** provide output for generating requests that specify a security strength greater than the instantiated security strength of the sub-DRBG.
14. A sub-DRBG **shall** not include a reseed capability.

1351 15. If the seedlife of a sub-DRBG is ever exceeded or a health test of the sub-DRBG fails, the
1352 use of the sub-DRBG **shall** be terminated.

1353 A non-testable requirement for a sub-DRBG (not testable by the validation labs) is:

1354 16. The output of a sub-DRBG **must not** be used as input to seed other DRBGs (e.g., the
1355 DRBGs in other RBGs).

5. RBG2 Constructions Based on Physical and/or Non-Physical Entropy Sources

An RBG2 construction is a cryptographically secure RBG with continuous access to one or more validated entropy sources within its RBG security boundary. The RBG is instantiated before use, generates outputs on demand, and can be used in an RBG3 construction (see [Section 6](#)). An RBG2 construction **may** support reseeding and may provide prediction resistance during generation requests (i.e., by performing a reseed of the DRBG prior to generating output). Both reseeding and providing prediction resistance are optional for this construction.

If full-entropy output is required by a consuming application, an RBG3 construction from [Section 6](#) needs to be used rather than an RBG2 construction.

An RBG2 construction may be useful for all devices in which an entropy source can be implemented.

5.1. RBG2 Description

The DRBG for an RBG2 construction is contained within the same RBG security boundary and cryptographic module as its validated entropy source(s) (see [Figure 16](#)). The entropy source is used to provide the entropy bits for both DRBG instantiation and the reseeding of the DRBG used by the construction (e.g., to provide prediction resistance). An optional recommended personalization string and optional additional input may be provided from within the cryptographic module or from outside of that module.

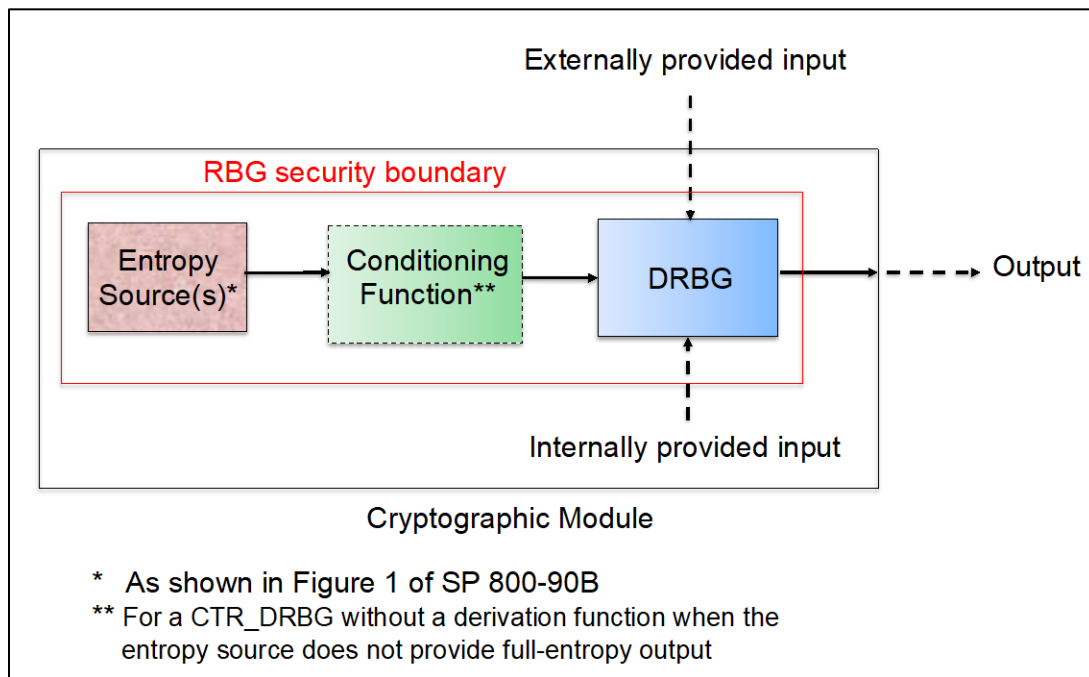


Fig. 16. RBG2 Construction

The output from the RBG may be used within the cryptographic module or by an application outside of the module.

An example of an RBG2 construction is provided in [Appendix B.4](#).

An RBG2 construction may be implemented to use one or more validated physical and/or non-physical entropy sources for instantiation and reseeding. Two variants of the RBG2 construction may be implemented.

1. An RBG2(P) construction uses the output of one or more validated physical entropy sources and (optionally) one or more validated non-physical entropy sources as discussed in Method 1 of [Section 2.3](#) (i.e., only the entropy produced by validated physical entropy sources is counted toward the entropy required for instantiating or reseeding the RBG). Any amount of entropy may be obtained from a non-physical entropy source as long as sufficient entropy has been obtained from the physical entropy sources to fulfill an entropy request.
2. An RBG2(NP) construction uses the output of any validated non-physical or physical entropy sources as discussed in Method 2 of [Section 2.3](#) (i.e., the entropy produced by both validated physical and non-physical entropy sources is counted toward the entropy required for instantiating or reseeding the RBG).

These variants affect the implementation of a **Get_ES_Bitstring** function (as specified in [Section 2.8.2.2](#) and discussed in [Section 3.1](#)), either accessing the entropy source directly or via the **Get_conditioned_full_entropy_input** function during instantiation and reseeding (see [Sections 5.2.1](#) and [5.2.3](#)). That is, when instantiating and reseeding an RBG2(P) construction (including a DRBG within an RBG3 construction as discussed in [Section 6](#)), Method 1 in [Section 2.3](#) is used to combine the entropy from the entropy sources, and Method 2 is used when instantiating and reseeding an RBG2(NP) construction.

5.2. Conceptual Interfaces

The RBG2 construction interfaces to the DRBG include function calls for instantiating the DRBG (see [Section 5.2.1](#)), generating pseudorandom bits on request (see [Section 5.2.2](#)), and (optionally) reseeding the DRBG at the end of the DRBG's seedlife and providing prediction resistance upon request (see [Section 5.2.3](#)).

Once instantiated, an RBG2 construction with a reseed capability may be reseeded on demand or whenever sufficient entropy is available.

5.2.1. RBG2 Instantiation

An RBG2 construction may be instantiated at any valid²⁰ security strength possible for the DRBG and its components using the following call:

(status, RBG2_state_handle) = Instantiate_function(s, prediction_resistance_flag, personalization_string),

²⁰ A security strength of either 128, 192, or 256 bits.

where s is the requested instantiation security strength for the DRBG. The *prediction_resistance_flag* (if used) is set to TRUE if prediction resistance is to be supported and FALSE otherwise.

An RBG2 construction obtains entropy for its DRBG from one or more validated entropy sources, either directly or using a conditioning function to process the output of the entropy source to obtain a full-entropy bitstring for instantiation (e.g., when employing a CTR_DRBG without a derivation function using entropy sources that do not provide full-entropy output).

SP 800-90A uses a **Get_randomness-source_input** call to obtain the entropy needed for instantiation (see SP 800-90A).

1. When the DRBG is a CTR_DRBG without a derivation function, full-entropy bits **shall** be obtained as follows:

- a) If the entropy source provides full-entropy output, the **Get_randomness-source_input** call is replaced by:^{21, 22}

$(status, entropy_bitstring) = \text{Get_ES_Bitstring}(s + 128).$ ²³

For an RBG2(P) construction, only validated physical entropy sources **shall** be used. The output of the entropy sources **shall** be concatenated to obtain the $s + 128$ full-entropy bits to be returned as *entropy_bitstring*.

(This recommendation assumes that non-physical entropy sources cannot provide full-entropy output. Therefore, the **Get_ES_bitstring** function **shall not** be used with non-physical entropy sources in this case.)

- b) If the entropy sources does not provide full-entropy output, the **Get_randomness-source_input** call is replaced by:^{24, 25}

$(status, Full_entropy_bitstring) =$
Get_conditioned_full_entropy_input($s + 128$).

Validated physical and/or non-physical entropy sources **shall** be used to provide the requested entropy. For an RBG2(P) construction, the requested $s + 128$ bits of entropy **shall** be counted as specified in Method 1 of [Section 2.3](#). For an RBG2(NP) construction, the requested $s + 128$ bits of entropy **shall** be counted as specified in Method 2 of [Section 2.3](#).

2. For the Hash_DRBG, HMAC_DRBG and CTR_DRBG (with a derivation function), the entropy source **shall** provide $3s/2$ bits of entropy to establish the security strength.

- a) If the consuming application requires full entropy in the returned bitstring, the **Get_randomness-source_input** call is replaced by:

$(status, Full_entropy_bitstring) =$
Get_conditioned_full_entropy_input($3s/2$).

²¹ Appropriate changes may be required for the **Instantiate_function** in [\[SP800-90A\]](#) and the algorithms in Section 10 of that document.

²² See Section 3.8.2.2 for a specification of the **Get_ES_Bitstring** function.

²³ For a CTR_DRBG using AES, $s + 128$ = the length of the key + the length of the AES block = *seedlen* (see Table 2 in SP 800-90A).

²⁴ Appropriate changes may be required for the **Instantiate_function** in [\[SP800-90A\]](#) and the algorithms in Section 10.2 of that document.

²⁵ See Section 4.3.2 for a specification of the **Get_conditioned_full_entropy_input** function.

- b) If the consuming application does not require full entropy in the returned bitstring, the **Get_randomness-source_input** call is replaced by:

$(status, entropy_bitstring) = \text{Get_ES_Bitstring}(3s/2).$

Validated physical and/or non-physical entropy sources **shall** be used to provide the requested entropy. For an RBG2(P) construction, the requested $3s/2$ bits of entropy **shall** be counted as specified in Method 1 of [Section 2.3](#). For an RBG2(NP) construction, the requested $3s/2$ bits of entropy **shall** be counted as specified in Method 2 of Section 3.3.

5.2.2. Requesting Pseudorandom Bits from an RBG2 Construction

Pseudorandom bits may be requested using the following call (see [Section 2.8.1.2](#)):

$(status, returned_bits) = \text{Generate_function}(RBG2_state_handle, requested_number_of_bits, requested_security_strength, prediction_resistance_request, additional_input),$

where *state_handle* (if used) was returned by the **Instantiate_function** (see Sections [2.8.1.1](#) and [5.2.1](#)).

Support for prediction resistance is optional. If prediction resistance is supported, its use is optional. This RBG may be designed to always provide prediction resistance, to only provide prediction resistance upon request, or to be unable to provide prediction resistance (i.e., to not support prediction-resistance requests during generation).

Note that when prediction resistance is requested, the **Generate_function** will invoke the **Reseed_function**. If sufficient entropy is not available for reseeding, an error indication **shall** be returned, and the requested bits **shall not** be generated.

5.2.3. Reseeding an RBG2 Construction

As discussed in [Section 2.4.2](#), when the RBG2 construction includes a reseed capability, the reseeding of the DRBG may be performed 1) upon request from a consuming application (either an explicit request for reseeding or a request for the generation of bits with prediction resistance); 2) on a fixed schedule based on time, number of outputs, or events; or 3) as sufficient entropy becomes available.

An RBG2 construction is reseeded using the following call:

$status = \text{Reseed_function}(RBG2_state_handle, additional_input),$

where the *RBG2_state_handle* (when used) was obtained during the instantiation of the RBG (see Sections [2.8.1.1](#) and [5.2.1](#)).

SP 800-90A uses a **Get_randomness-source_input** call to obtain the entropy needed for reseeding the DRBG (see [Section 2.8.1.3](#) herein and in [\[SP800-90A\]](#)). The DRBG is reseeded at the instantiated security strength recorded in the DRBG's internal state. The **Get_randomness-source_input** call in SP 800-90A **shall** be replaced with the following:

1. For the CTR_DRBG without a derivation function, use the appropriate replacement as specified in step 1 of [Section 5.2.1](#).

2. For the Hash_DRBG, HMAC_DRBG and CTR_DRBG (with a derivation function), replace the **Get_randomness-sourceinput** call in the **Reseed_function** with the following:²⁶

- a) If the consuming application requires full entropy in the returned bitstring, the **Get_randomness-source_input** call is replaced by:

(status, Full_entropy_bitstring) = Get_conditioned_full_entropy_input(s).

- b) If the consuming application does not require full entropy in the returned bitstring, the **Get_randomness-source_input** call is replaced by:

(status, entropy_bitstring) = Get_ES_Bitstring(s).

Validated physical and/or non-physical entropy sources **shall** be used to provide the requested entropy. For an RBG2(P) construction, the requested s bits of entropy **shall** be counted as specified in Method 1²⁷ of [Section 2.3](#). For an RBG2(NP) construction, the requested s bits of entropy **shall** be counted as specified in Method 2²⁸ of [Section 2.3](#).

5.3. RBG2 Requirements

An RBG2 construction has the following requirements in addition to those specified in [\[SP800-90A\]](#):

1. The RBG **shall** employ an **approved** and validated DRBG from [\[SP800-90A\]](#) whose components are capable of providing the targeted security strength for the RBG.
2. The RBG and its components **shall** be successfully validated for compliance with [\[SP800-90A\]](#), [\[SP800-90B\]](#), SP 800-90C, [\[FIPS140\]](#), and the specification of any other **approved** algorithm used within the RBG, as appropriate.
3. The RBG **may** include a reseed capability. If implemented, the reseeding of the DRBG **shall** be performed either a) upon request from a consuming application (either an explicit request for reseeding or a request for the generation of bits with prediction resistance); b) on a fixed schedule based on time, number of outputs, or events; and/or c) as sufficient entropy becomes available.
4. Validated entropy sources **shall** be used to instantiate and reseed the DRBG. A non-validated entropy sources **shall not** be used for this purpose.
5. The entropy sources used for the instantiation and reseeding of an RBG(P) construction **shall** include one or more validated physical entropy sources; the inclusion of one or more validated non-physical entropy sources is optional. A bitstring that contains entropy **shall** be assembled and the entropy in that bitstring determined as specified in Method 1 of [Section 2.3](#) (i.e., only the entropy provided by validated physical entropy sources **shall** be counted toward fulfilling the amount of entropy in an entropy request).

²⁶ See Sections 2.8.2.2 and 3.1 for discussions of the Get_ES_bitstring function.

²⁷ Method 1 only counts the entropy provided by validated physical sources.

²⁸ Method 2 counts the entropy provided by both physical and non-physical entropy sources.

6. The entropy sources used for the instantiation and reseeding of an RBG2(NP) construction **shall** include one or more validated non-physical entropy sources; the inclusion of one or more validated physical entropy sources is optional. A bitstring containing entropy **shall** be assembled and the entropy in that bitstring determined as specified in Method 2 of [Section 2.3](#) (i.e., the entropy provided by both validated non-physical entropy sources and any validated physical entropy sources included in the implementation **shall** be counted toward fulfilling the requested amount of entropy).
7. The DRBG **shall** be capable of being instantiated and reseeded at the maximum security strength (s) for the DRBG design (see [[SP800-90A](#)]).
8. A specific entropy-source output (or portion thereof) **shall not** be reused (e.g., it is destroyed after use).
9. When instantiating and reseeding a CTR_DRBG without a derivation function, $(s + 128)$ bits with full entropy **shall** be obtained either directly from the entropy source or from the entropy source via an external vetted conditioning function (see [Section 3.3](#)).
10. For a Hash_DRBG, HMAC_DRBG or CTR_DRBG (with a derivation function), a bitstring with at least $3s/2$ bits of entropy **shall** be obtained from the entropy source to instantiate the DRBG at a security strength of s bits. When reseeding is performed, a bitstring with at least s bits of entropy **shall** be obtained from the entropy source.
11. The DRBG **shall** be instantiated before first use (i.e., before providing output for use by a consuming application) and reseeded using the validated entropy sources used for instantiation.
12. When health tests detect the failure of a validated entropy source, the failure **shall** be handled as discussed in [Section 7.1.2.1](#).

A non-testable requirement for the RBG (not testable by the validation labs) is:

13. The RBG **must not** be used by applications that require a higher security strength than has been instantiated in the DRBG.

6. RBG3 Constructions Based on Physical Entropy Sources

An RBG3 construction is designed to provide full entropy (i.e., an RBG3 construction can support all security strengths). The RBG3 constructions specified in this Recommendation include one or more entropy sources and an **approved** DRBG from SP 800-90A that can and will be instantiated at a security strength of 256 bits. If an entropy source fails in an undetected manner, the RBG continues to operate as an RBG2(P) construction, providing outputs at the security strength of its DRBG (256 bits) (see [Section 5](#) and [Appendix A](#)). If a failure is detected, the RBG operation **shall** be terminated.

Two RBG3 constructions are specified:

1. RBG3(XOR) – This construction is based on combining the output of one or more validated entropy sources with the output of an instantiated, **approved** DRBG using an exclusive-or operation (see [Section 6.2](#)).
2. RBG3(RS) – This construction is based on using one or more validated entropy sources to continuously reseed the DRBG (see [Section 6.3](#)).

An RBG3 construction continually accesses its entropy sources, and its DRBG may be reseeded whenever requested (e.g., to provide prediction resistance for the DRBG's output). Upon receipt of a request for random bits from a consuming application, the entropy source is accessed to obtain sufficient bits for the request. See Sections [3.1](#) and [3.2](#) for further discussion about accessing the entropy source(s).

An implementation may be designed so that the DRBG implementation used within an RBG3 construction can be directly accessed by a consuming application (i.e., the directly accessible DRBG uses the same internal state as the RBG3 construction).

An RBG3 construction is useful when bits with full entropy are required or a higher security strength than RBG1 and RBG2 constructions can support is needed.

6.1. General Requirements

RBG3 constructions have the following general security requirements. See Sections [6.2.2](#) and [6.3.2](#) for additional requirements for the RBG3(XOR) and RBG3(RS) constructions, respectively.

1. An RBG3 construction **shall** be designed to provide outputs with full entropy using one or more validated independent physical entropy sources as specified for Method 1 in [Section 3.3](#) (i.e., only the entropy provided by validated physical entropy sources **shall** be counted toward fulfilling entropy requests, although entropy provided by any validated non-physical entropy source may be used but not counted).
2. An RBG3 construction and its components **shall** be successfully validated for compliance with the corresponding requirements in [\[SP800-90A\]](#), [\[SP800-90B\]](#), SP 800-90C, [\[FIPS 140\]](#) and the specification of any other **approved** algorithm used within the RBG, as appropriate.
3. The DRBG within the RBG3 construction **shall** be capable of supporting a security strength of 256 bits (i.e., a CTR_DRBG based on AES-256 or either Hash_DRBG or HMAC_DRBG using a hash function with an output length of at least 256 bits).

4. The DRBG **shall** be instantiated at a security strength of 256 bits before the first use of the RBG3 construction or direct access of the DRBG.
5. The DRBG **shall** include a reseed function to support reseed requests.
6. A specific entropy-source output (or portion thereof) **shall not** be reused (e.g., the same entropy-source outputs **shall not** be used for an RBG3 request and a request to a separate instantiation of a DRBG).
7. If the DRBG is directly accessible, the requirements in [Section 5.3](#) for RBG2(P) constructions **shall** apply to the direct access of the DRBG.
8. When health tests detect the failure of a validated physical entropy source, the failure **shall** be handled as discussed in [Section 7.1.2.1](#). If a failure is detected in a non-physical entropy source, the consuming application **shall** be notified.

6.2. RBG3(XOR) Construction

An RBG3(XOR) construction contains one or more validated entropy sources and a DRBG whose outputs are XORed to produce full-entropy output (see [Figure 17](#)). In order to provide the required full-entropy output, the input to the XOR (shown as “ \oplus ” in the figure) from the entropy-source side of the figure **shall** consist of bits with full entropy (see [Section 2.1](#)).²⁹ If the entropy sources cannot provide full-entropy output, then an external conditioning function **shall** be used to condition the output of the entropy sources to a full-entropy bitstring before XORing with the output of the DRBG (see [Section 3.3](#)).

²⁹ Note that the DRBGs themselves are not designed to inherently provide full-entropy output.

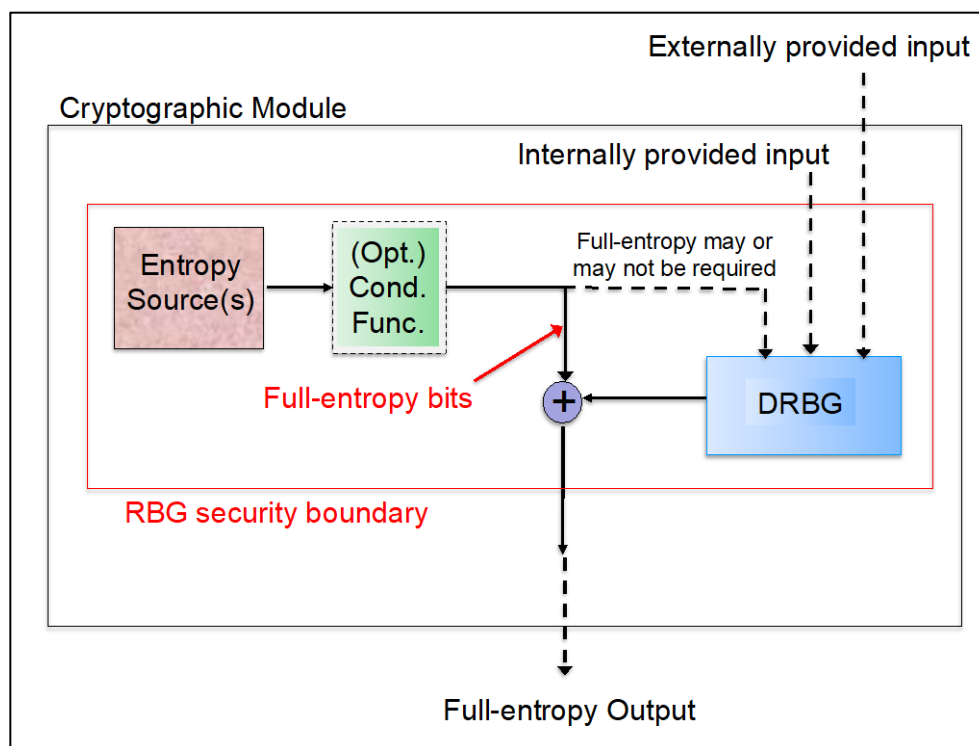


Fig. 17. RBG3(XOR) Construction

When n bits of output are requested from an RBG3(XOR) construction, n bits of output from the DRBG are XORed with n full-entropy bits obtained either directly from the entropy source or from the entropy source after cryptographic processing by an external vetted conditioning function (see [Section 3.3](#)). When the entropy source is working properly,³⁰ an n -bit output from the RBG3(XOR) construction is said to provide n bits of entropy or to support a security strength of n bits. The DRBG used in the RBG3(XOR) construction is always required to support a 256-bit security strength. If the entropy source fails without being detected and the DRBG has been successfully instantiated with at least 256 bits of entropy, the DRBG continues to produce output at a security strength of 256 bits.

An example of an RBG3(XOR) design is provided in [Appendix B.5](#).

6.2.1. Conceptual Interfaces

The RBG interfaces include function calls for instantiating the DRBG (see [Section 6.2.1.1](#)), generating random bits on request (see [Section 6.2.1.2](#)), and reseeding the DRBG instantiation(s) (see [Section 6.2.1.3](#)).

6.2.1.1. Instantiation of the DRBG

The DRBG for the RBG3(XOR) construction is instantiated as follows:

³⁰ The entropy source provides at least the amount of entropy determined during the entropy-source validation process.

1619 **RBG3(XOR)_DRBG_Instantiate:**

1620 **Input:** integer (*prediction_resistance_flag*), string *personalization_string*.

1621 **Output:** integer *status*, integer *state_handle*.

1622 **Process:**

1623 1. (*status*, *RBG3(XOR)_state_handle*) = **Instantiate_function**(256,
1624 *prediction_resistance_flag*, *personalization_string*).

1625 2. Return (*status*, *RBG3(XOR)_state_handle*).

1626 In step 1, the DRBG is instantiated at a security strength of 256 bits. The
1627 *prediction_resistance_flag* and *personalization_string* (when provided as input to the
1628 **RBG3(XOR)_DRBG_Instantiate** function) **shall** be used in step 1.

1629 In step 2, the *status* and *RBG3(XOR)_state_handle* that were obtained in step 1 are returned. Note
1630 that if the *status* does not indicate a successful instantiate process (i.e., a failure is indicated), the
1631 returned state handle **shall** be invalid (e.g., a *Null* value). The handling of status codes is discussed
1632 in [Section 2.8.3](#).

1633 **6.2.1.2. Random and Pseudorandom Bit Generation**

1634 Let *n* be the requested number of bits to be generated, and let the *RBG3(XOR)_state_handle* be
1635 the value returned by the instantiation function for RBG3's DRBG instantiation (see [Section](#)
1636 [6.2.1.1](#)). Random bits with full entropy **shall** be generated by the RBG3(XOR) construction using
1637 the following generate function:

1638 **RBG3(XOR)_Generate:**

1639 **Input:** integer (*RBG3(XOR)_state_handle*, *n*, *prediction_resistance_request*), string
1640 *additional_input*.

1641 **Output:** integer *status*, string *returned_bits*.

1642 **Process:**

1643 1. (*status*, *ES_bits*) = **Request_entropy**(*n*).

1644 2. If (*status* ≠ SUCCESS), then return (*status*, *invalid_string*).

1645 3. (*status*, *DRBG_bits*) = **Generate_function**(*RBG3(XOR)_state_handle*, *n*, 256,
1646 *prediction_resistance_request*, *additional_input*).

1647 4. If (*status* ≠ SUCCESS), then return (*status*, *invalid_string*).

1648 5. *returned_bits* = *ES_bits* ⊕ *DRBG_bits*.

1649 6. Return (SUCCESS, *returned_bits*).

1650 Step 1 requests that the entropy sources generate bits. Since full-entropy bits are required, the
1651 (place holder) **Request_entropy** call **shall** be replaced by one of the following:

- If full-entropy output is provided by all validated physical entropy sources used by the RBG3(XOR) implementation, and non-physical entropy sources are not used,³¹ step 1 becomes:

$(status, ES_bits) = \text{Get_ES_Bitstring}(n).$

The **Get_ES_Bitstring** function³² **shall** use Method 1 in [Section 2.3](#) to obtain the n full-entropy bits that were requested in order to produce the ES_bits bitstring.

- If full-entropy output is not provided by all physical entropy sources, or the output of both physical and non-physical entropy sources is also used by the implementation, step 1 becomes:

$(status, ES_bits) = \text{Get_conditioned_full_entropy_input}(n).$

The **Get_conditioned_full_entropy_input** construction is specified in [Section 3.3.2](#). It requests entropy from the entropy sources in step 3.1 of that construction with a **Get_ES_Bitstring** call. The **Get_ES_Bitstring** call **shall** use Method 1 (as specified in [Section 3.3](#)) when collecting the output of the entropy sources (i.e., only the entropy provided by physical entropy sources is counted).

In step 2, if the request in step 1 is not successful, abort the **RBG3(XOR)_Generate** function, returning the *status* received in step 1 and an invalid bitstring as the *returned_bits* (e.g., a *Null* bitstring). If *status* indicates a success, ES_bits is the full-entropy bitstring to be used in step 5.

In step 3, the RBG3(XOR)'s DRBG instantiation is requested to generate n bits at a security strength of 256 bits. The DRBG instantiation is indicated by the $RBG3(XOR)_state_handle$, which was obtained during instantiation (see [Section 6.2.1.1](#)). If a prediction-resistance request and/or additional input are provided in the **RBG3(XOR)_Generate** call, they **shall** be included in the **Generate_function** call.

Note that it is possible that the DRBG would require reseeding during the **Generate_function** call in step 3 (e.g., because of a prediction-resistance request, or the end of the seedlife of the DRBG has been reached). If a reseed of the DRBG is required during **Generate_function** execution, the DRBG **shall** be reseeded as specified in [Section 6.2.1.3](#) with bits not otherwise used by the RBG.

In step 4, if the **Generate_function** request is not successful, the **RBG3(XOR)_Generate** function is aborted, and the *status* received in step 3 and an invalid bitstring (e.g., a *Null* bitstring) are returned to the consuming application. If *status* indicates a success, $DRBG_bits$ is the pseudorandom bitstring to be used in step 5.

Step 5 combines the bitstrings returned from the entropy sources (from step 1) and the DRBG (from step 3) using an XOR operation. The resulting bitstring is returned to the consuming application in step 6.

³¹ Since non-physical entropy sources are assumed to be incapable of providing full-entropy output, they cannot contribute to the bitstring provided by the **Get_ES_Bitstring** function.

³² See Section 3.10.2.2.

6.2.1.3. Pseudorandom Bit Generation Using a Directly Accessible DRBG

Pseudorandom bit generation by a direct access of the DRBG is accomplished as specified in [Section 5.2.2](#) using the state handle obtained during instantiation (see [Section 6.2.1.1](#)).

When directly accessing the DRBG instantiation that is also used by the RBG3(XOR) construction, the following function is used:

*(status, returned_bits) = **Generate_function**(RBG3(XOR)_state_handle,
requested_number_of_bits, requested_security_strength, prediction_resistance_request,
additional_input),*

where:

- *RBG3(XOR)_state_handle* indicates the DRBG instantiation to be used.
- *requested_security_strength* ≤ 256.
- *prediction-resistance-request* is either TRUE or FALSE; requesting prediction resistance during the **Generate function** is optional.
- The use of additional input is optional.

Note that when prediction resistance is requested, the **Generate_function** will invoke the **Reseed_function** (see [Section 6.2.1.3](#)). If sufficient entropy is not available for reseeding, an error indication **shall** be returned, and the requested bits **shall not** be generated.

6.2.1.4. Reseeding the DRBG Instantiations

Reseeding is performed using the entropy sources in the same manner as an RBG2 construction using the appropriate state handle (e.g., *RBG3(XOR)_state_handle*, as specified in [Section 6.2.1.1](#)).

6.2.2. RBG3(XOR) Requirements

An RBG3(XOR) construction has the following requirements in addition to those provided in [Section 6.2](#):

1. Bitstrings with full entropy **shall** be provided to the XOR operation either directly from the concatenated output of one or more validated physical entropy sources or by an external conditioning function using the output of one or more validated entropy sources as specified in Method 1 of [Section 2.3](#). In the latter case, the output of validated non-physical entropy sources may be used without counting any entropy that they might provide.
2. The same entropy-source outputs used by the DRBG for instantiation or reseeding **shall not** be used as input into the RBG's XOR operation.
3. The DRBG instantiations **shall** be reseeded occasionally (e.g., after a predetermined period of time or number of generation requests).

6.3. RBG3(RS) Construction

The second RBG3 construction specified in this document is the RBG3(RS) construction shown in [Figure 18](#), and an example of this construction is provided in [Appendix B.6](#).

Note that external conditioning of the outputs from the entropy sources during instantiation and reseeding is required when the DRBG is a CTR_DRBG without a derivation function and the entropy sources do not provide a bitstring with full entropy.

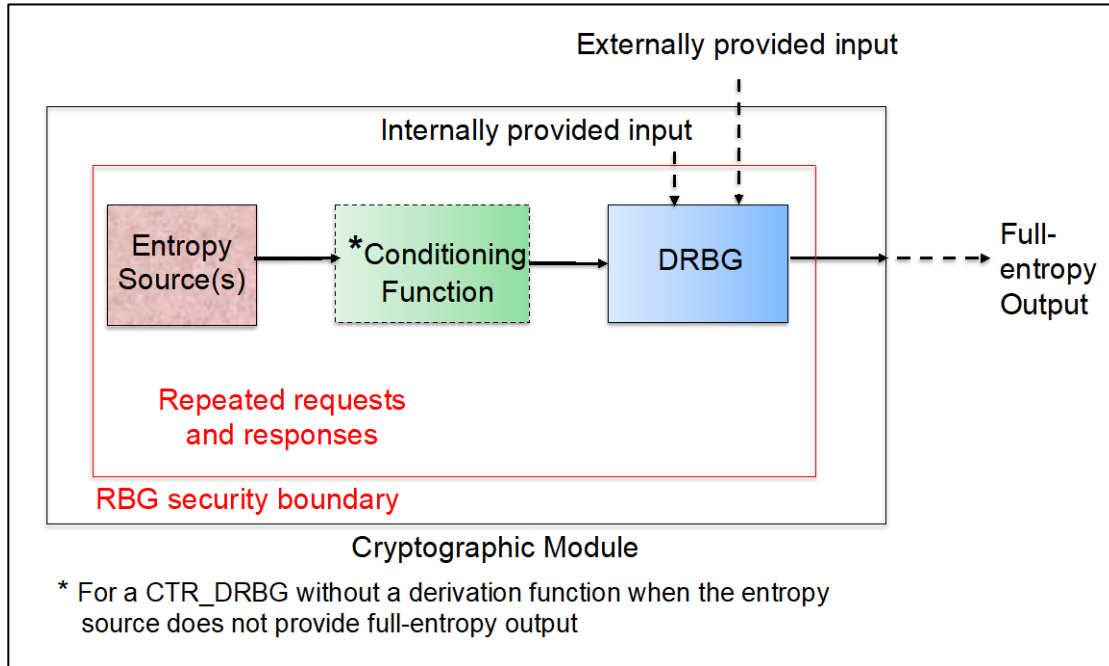


Fig. 18. RBG3(RS) Construction

6.3.1. Conceptual Interfaces

The RBG interfaces include function calls for instantiating the DRBG (see [Section 6.3.1.1](#)), generating random bits on request (see [Section 6.3.1.2](#)), and reseeding the DRBG instantiation (see [Section 6.3.1.3](#)).

6.3.1.1. Instantiation of the DRBG Within an RBG3(RS) Construction

DRBG instantiation is performed as follows:

RBG3(RS)_DRBG_Instantiate:

Input: integer (*prediction_resistance_flag*), string *personalization_string*.

Output: integer *status*, integer *state_handle*.

Process:

1. (*status*, *RBG3(RS)_state_handle*) = **Instantiate_function**(256, *prediction_resistance_flag* = TRUE, *personalization_string*).

2. Return (*status*, *RBG3(RS)_state_handle*).

In step 1, the DRBG is instantiated at a security strength of 256 bits. The *prediction_resistance_flag* is set to TRUE, and *personalization_string* (when provided as input to the **RBG3(RS)_DRBG_Instantiate** function) **shall** be used in step 1.

In step 2, the *status* and the *RBG3(RS)_state_handle* are returned. Note that if the *status* does not indicate a successful instantiate process (i.e., a failure is indicated), the returned state handle **shall** be invalid (e.g., a *Null* value). The handling of status codes is discussed in [Section 2.8.3](#).

6.3.1.2. Random and Pseudorandom Bit Generation

6.3.1.2.1 Generation Using the RBG3(RS) Construction

When an RBG3(RS) construction receives a request for *n* random bits, the DRBG instantiation used by the construction needs to be reseeded with sufficient entropy so that bits with full entropy can be extracted from the DRBG's output block.

Table 3 provides information for generating full-entropy output from the DRBGs in SP 800-90A that use the cryptographic primitives listed in the table. Each primitive in the table can support a security strength of 256 bits – the highest security strength recognized by this Recommendation. To use the table, select the row that identifies the cryptographic primitive used by the implemented DRBG.

- Column 1 lists the DRBGs.
- Column 2 identifies the cryptographic primitives that can be used by the DRBG(s) in column 1 to support a security strength of 256 bits.
- Column 3 indicates the length of the output block (*output_len*) for the cryptographic primitives in column 2.
- Column 4 indicates the amount of fresh entropy that is obtained by a **Reseed_function** when the **Generate_function** is invoked with prediction resistance requested.

Table 3. Values for generating full-entropy bits by an RBG3(RS) Construction

DRBG	DRBG Primitives	Output Block Length (<i>output_len</i>) in bits	Entropy obtained during a normal reseed operation
CTR_DRBG (with no derivation function)	AES-256	128	384
CTR_DRBG (using a derivation function)	AES-256	128	256
Hash_DRBG or HMAC_DRBG	SHA-256 SHA3-256	256	256
	SHA-384 SHA3-384	384	256
	SHA-512 SHA3-512	512	256

1763 The strategy used for obtaining full-entropy output from the RBG3(RS) construction requires
1764 obtaining sufficient fresh entropy and subsequently extracting full entropy bits from the output
1765 block in accordance with item 11 of [Section 2.6](#).

1766 For the **RBG3(RS)_Generate** function:

- 1767 • Let n be the requested number of full-entropy bits to be generated by an RBG3(RS)
1768 construction.
- 1769 • Let *RBG3(RS)_state_handle* be a state handle returned from the instantiate function (see
1770 [Section 6.3.1.1](#)).

1771 Random bits with full entropy **shall** be generated as follows:

1772 **RBG3(RS)_Generate:**

1773 **Input:** integer (*RBG3(RS)_state_handle*, n), string *additional_input*.

1774 **Output:** integer *status*, bitstring *returned_bits*.

1775 **Process:**

- 1776 1. *full-entropy_bits* = *Null*.
- 1777 2. *sum* = 0.
- 1778 3. While (*sum* < n),
 - 1779 3.1 Obtain *generated_bits* from the entropy source.
 - 1780 3.2 If (*status* ≠ SUCCESS), then return (*status*, *invalid_bitstring*).
 - 1781 3.3 *full-entropy_bits* = *full_entropy_bits* || *generated_bits*.
 - 1782 3.4 *sum* = *sum* + **len**(*generated_bits*).
- 1783 4. Return (SUCCESS, **leftmost**(*full-entropy_bits*, n)).

1784 In steps 1 and 2, the bitstring intended to collect the generated bits for returning to the calling
1785 application (i.e., *full-entropy_bits*) is initialized to the *Null* bitstring, and the counter for the number
1786 of bits obtained for fulfilling the request is initialized to zero.

1787 Step 3 is iterated until n bits have been generated.

1788 In step 3.1, the DRBG is requested to obtain sufficient entropy so that a bitstring with full
1789 entropy can be extracted from the output block. The form of the request depends on the DRBG
1790 algorithm used in the RBG3(RS) construction and the method for obtaining a full-entropy
1791 bitstring (see [Section 2.6](#), item 11). Note that extracting fewer full-entropy bits from the
1792 DRBG's output block is permitted.

1793 For a CTR_DRBG (with or without a derivation function), a maximum of 128 bits with
1794 full entropy can be provided from the AES output block for each iteration of the DRBG as
1795 follows:

1796 (*status*, *generated_bits*) = **Generate_function**(*RBG3(RS)_state_handle*, 128,
1797 256, *prediction_resistance_request* = TRUE, *additional_input*).

1798 The **Generate_function** generates 128 (full entropy) bits after reseeding the
1799 CTR_DRBG with either 256 or 384 bits of entropy (by setting
1800 *prediction_resistance_request* = TRUE).³³

1801 For a hash-based DRBG (i.e., Hash_DRBG and HMAC_DRBG), a maximum of 256 full-
1802 entropy bits can be produced from each iteration of the DRBG as follows:

1803 3.1.1 (*status*, *additional_entropy*) = **Get_ES_Bitstring** (64).

1804 3.1.2 If (*status* ≠ SUCCESS), then return (*status*, *invalid_bitstring*).

1805 3.1.3 (*status*, *generated_bits*) = **Generate_function**(*RBG3(RS)_state_handle*,
1806 256, 256, *prediction_resistance_request* = TRUE, *additional_input* ||
1807 *additional_entropy*).

1808 At least 64 bits of entropy beyond the amount obtained during reseeding are required.
1809 As shown in [Table 3](#), the reseeding process will acquire 256 bits of entropy. The (256
1810 + 64 = 384) bits of entropy are inserted into the DRBG by 1) obtaining a bitstring with
1811 at least 64 bits of entropy directly from the entropy sources (step 3.1.1), 2)
1812 concatenating the additional entropy bits with any *additional_input* provided in the
1813 **RBG3(RS)_Generate** call, and 3) requesting the generation of 256 bits with prediction
1814 resistance and including the concatenated bitstring. This results in both the reseed of
1815 the DRBG with 256 bits of entropy and the insertion of the additional 64 bits of entropy)
1816 (step 3.1.3).

1817 For a hash-based DRBG (i.e., Hash_DRBG and HMAC_DRBG), a maximum of 192 full-
1818 entropy bits can be produced from each iteration of the DRBG as follows:

1819 (*status*, *generated_bits*) = **Generate_function**(*RBG3(RS)_state_handle*, 192,
1820 256, *prediction_resistance_request* = TRUE, *additional_input*).

1821 The DRBG is reseeded with 256 bits of entropy by requesting generation with prediction
1822 resistance and extracting only (256 – 64 = 192) bits from the DRBG’s output block as
1823 full-entropy bits.

1824 In step 3.2, if the **Generate_function** request invoked in step 3.1 is not successful, the
1825 **RBG3(RS)_Generate** function is aborted, and the *status* received in step 3.1 and an invalid
1826 bitstring (e.g., a *Null* bitstring) are returned to the consuming application.

1827 Step 3.3 combines the full-entropy bitstrings obtained in step 3.1 with previously generated
1828 full-entropy bits using a concatenation operation.

1829 Step 3.4 adds the number of full-entropy bits produced in step 3.1 to those generated in
1830 previous iterations of step 3.

1831 If *sum* is less than the requested number of bits (*n*), repeat step 3 starting at step 3.1.

1832 In step 4, the leftmost *n* bits are selected from the collected bitstring (i.e., *full-entropy_bits*) and
1833 returned to the consuming application.

1834 **6.3.1.2.2 Generation Using a Directly Accessible DRBG**

³³ The use of the *prediction_resistance_request* will handle the differences between the two versions of the CTR_DRBG (i.e., with or without a derivation function).

Direct access of the DRBG is accomplished as specified in [Section 5.2.2](#) using the state handle associated with the instantiation and internal state that was returned for the DRBG (see [Section 6.3.1.1](#)).

(status, returned_bits) = Generate_function(RBG3(RS)_state_handle,
requested_number_of_bits, requested_security_strength, prediction_resistance_request,
additional_input),

where *state_handle* (if used) was returned by the **Instantiate_function** (see [Section 6.3.1.1](#)).

When the previous generate request was made to the RBG3(RS) construction rather than directly to the DRBG, the *prediction_resistance_request* parameter **shall** be set to TRUE. Otherwise, requesting prediction resistance during the **Generate_function** is optional.

6.3.1.3. Reseeding

Reseeding is performed during a **Generate_function** request to a directly accessible DRBG (see [Section 6.3.1.2.2](#)) when prediction resistance is requested or the end of the DRBG's seedlife is reached. The **Generate_function** invokes the **Reseed_function** specified in [SP800-90A].

Reseeding may also be performed on demand as specified in [Section 4.2.3](#) using the *RBG3(RS)_state_handle* if provided during instantiation.

6.3.2. Requirements for a RBG3(RS) Construction

An RBG3(RS) construction has the following requirements in addition to those provided in [Section 6.1](#):

1. Fresh entropy **shall** be acquired either directly from all independent validated entropy sources (see [Section 3.2](#)) or (in the case of a CTR_DRBG used as the DRBG when the entropy sources do not provide full-entropy output) from an external conditioning function that processes the output of the validated entropy sources as specified in [Section 3.3.2](#). Method 1 in [Section 2.3](#) **shall** be used when collecting the required entropy (i.e., only the entropy provided by validated physical entropy sources **shall** be counted toward fulfilling the amount of entropy requested).
2. If the DRBG is directly accessible, a reseed of the DRBG instantiation **shall** be performed before generating output in response to a request for output from the directly accessible DRBG when the previous use of the DRBG was by the RBG3(RS) construction. This could require an additional internal state value to record the last use of the DRBG for generation (e.g., used by an **RBG3(RS)_Generate** function as specified in [Section 6.3.1.2.1](#) or directly accessed by a (DRBG) **Generate_function** as discussed in [Section 6.3.1.2.2](#)).

7. Testing

Two types of testing are specified in this Recommendation: health testing and implementation-validation testing. Health testing **shall** be performed on all RBGs that claim compliance with this Recommendation (see [Section 7.1](#)). [Section 7.2](#) provides requirements for implementation validation.

7.1. Health Testing

Health testing is the testing of an implementation prior to and during normal operations to determine that the implementation continues to perform as expected and as validated. Health testing is performed by the RBG itself (i.e., the tests are designed into the RBG implementation).

An RBG **shall** support the health tests specified in [\[SP800-90A\]](#) and [\[SP800-90B\]](#) as well as perform health tests on the components of SP 800-90C (see [Section 7.1.1](#)). [\[FIPS 140\]](#) specifies the testing to be performed within a cryptographic module.

7.1.1. Testing RBG Components

Whenever an RBG receives a request to start up or perform health testing, a request for health testing **shall** be issued to the RBG components (e.g., the DRBG and any entropy source).

7.1.2. Handling Failures

Failures may occur during the use of entropy sources and during the operation of other components of an RBG.

Note that [\[SP800-90A\]](#) and [\[SP800-90B\]](#) discuss the error handling for DRBGs and entropy sources, respectively.

7.1.2.1. Entropy-Source Failures

A failure of a validated entropy source may be reported to the **Get_ES_Bitstring** function (see item 3 of [Section 3.1](#) and item 4 of [Section 3.2](#)) during entropy requests to the entropy sources or to the RBG when the entropy sources continue to function when entropy is not requested (see item 5 of [Section 3.2](#)).

7.1.2.2. Failures by Non-Entropy-Source Components

Failures by non-entropy-source components may be caused by either hardware or software failures. Some of these may be detected using the health testing within the RBG using known-answer tests. Failures could also be detected by the system in or on which the RBG resides.

When such failures are detected that affect the RBG, RBG operation **shall** be terminated. The RBG **must not** be resumed until the reasons for the failure have been determined and the failures have been repaired and successfully tested for proper operation.

7.2. Implementation Validation

Implementation validation is the process of verifying that an RBG and its components fulfill the requirements of this Recommendation. Validation is accomplished by:

- Validating the components from [\[SP800-90A\]](#) and [\[SP800-90B\]](#).
- Validating the use of the constructions in SP 800-90C via code inspection, known-answer tests, or both, as appropriate.
- Validating that the appropriate documentation as specified in SP 800-90C has been provided (see below).

Documentation **shall** be developed that will provide assurance to testers that an RBG that claims compliance with this Recommendation has been implemented correctly. This documentation **shall** include the following as a minimum:

- An identification of the constructions and components used by the RBG, including a diagram of the interaction between the constructions and components.
- If an external conditioning function is used, an indication of the type of conditioning function and the method for obtaining any keys that are required by that function.
- Appropriate documentation, as specified in [\[SP800-90A\]](#) and [\[SP800-90B\]](#). The DRBG and the entropy sources **shall** be validated for compliance with SP 800-90A or SP 800-90B, respectively, and the validations successfully finalized before the completion of RBG implementation validation.
- For an RBG1 or RBG2 construction, the maximum security-strength that can be supported by the DRBG.
- A description of all validated and non-validated entropy sources used by the RBG, including identifying whether the entropy source is a physical or non-physical entropy source.
- Documentation justifying the independence of all validated entropy sources from all other validated and non-validated entropy sources.
- An identification of the features supported by the RBG (e.g., access to the underlying DRBG of an RBG3 construction).
- A description of the health tests performed, including an identification of the periodic intervals for performing the tests.
- A description of any support functions other than health testing.
- A description of the RBG components within the RBG security boundary (see [Section 2.5](#)).
- For an RBG1 construction, a statement indicating that the randomness source **must** be a validated RBG2(P) or RBG3 construction (e.g., this could be provided in user documentation and/or a security policy).
- If sub-DRBGs can be used in an RBG1 construction, the maximum number of sub-DRBGs and the security strengths to be supported by the sub-DRBGs.

- 1936 • For an RBG2 construction (including a directly accessible DRBG within an RBG3
1937 construction), a statement indicating whether prediction resistance is always provided
1938 when a request is made by a consuming application, only provided when requested, or
1939 never provided.
- 1940 • For an RBG3 construction, a statement indicating whether the DRBG can be accessed
1941 directly.
- 1942 • Documentation specifying the guidance to users about fulfilling the non-testable
1943 requirements for RBG1 constructions, RBG2 constructions, and sub-DRBGs, as
1944 appropriate (see Sections [5.4](#) and [6.3](#), respectively).

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2012 **Appendix A. Entropy vs. Security Strength (Informative)**

2013 This section of the appendix compares and contrasts entropy and security strength.

2014 **A.1. Entropy**

2015 Suppose that an entropy source produces n -bit strings with m bits of entropy in each bitstring. This
2016 means that when an n -bit string is obtained from that entropy source, the best possible guess of the
2017 value of the string has a probability of no more than 2^{-m} of being correct.

2018 Entropy can be thought of as a property of a probability distribution, like the mean or variance.
2019 Entropy measures the unpredictability or randomness of the *probability distribution on bitstrings*
2020 *produced by the entropy source*, not a property of any particular bitstring. However, the
2021 terminology is sometimes slightly abused by referring to a bitstring as having m bits of entropy.
2022 This simply means that the bitstring came from a source that ensures m bits of entropy in its output
2023 bitstrings.

2024 Because of the inherent variability in the process, predicting future entropy-source outputs does
2025 not depend on an adversary's amount of computing power.

2026 **A.2. Security Strength**

2027 A deterministic cryptographic mechanism (such as one of the DRBGs defined in [\[SP800-90A\]](#))
2028 has a security strength – a measure of how much computing power an adversary expects to need
2029 to defeat the security of the mechanism. If a DRBG has an s -bit security strength, an adversary
2030 who can make 2^w computations of the underlying block cipher or hash function, where $w < s$,
2031 expects to have about a 2^{w-s} probability of defeating the DRBG's security. For example, an
2032 adversary who can perform 2^{96} AES encryptions can expect to defeat the security of the CTR-
2033 DRBG that uses AES-128 with a probability of about 2^{-32} (i.e., 2^{96-128}).

2034 **A.3. A Side-by-Side Comparison**

2035 Informally, one way of thinking of the difference between security strength and entropy is the
2036 following: suppose that an adversary somehow obtains the internal state of an entropy source (e.g.,
2037 the state of all of the ring oscillators and any internal buffer). This might allow the adversary to
2038 predict the next few bits from the entropy source (assuming that there is some buffering of bits
2039 within the entropy source), but the entropy source outputs will once more become unpredictable
2040 to the adversary very quickly. For example, knowing what faces of the dice are showing on the
2041 craps table does not allow a player to successfully predict the next roll of the dice.

2042 In contrast, suppose that an adversary somehow obtains the internal state of a DRBG. Because the
2043 DRBG is deterministic, the adversary can then predict all future outputs from the DRBG until the
2044 next reseeding of the DRBG with a sufficient amount of entropy.

2045 An entropy source provides bitstrings that are hard for an adversary to guess correctly but usually
2046 have some detectable statistical flaws (e.g., they may have slightly biased bits, or successive bits
2047 may be correlated). However, a well-designed DRBG provides bitstrings that exhibit none of these

properties. Rather, they have independent and identically distributed bits, with each bit taking on a value with a probability of exactly 0.5. These bitstrings are only unpredictable to an adversary who does not know the DRBG's internal state.

A.4. Entropy and Security Strength in this Recommendation

In the RBG1 construction specified in [Section 4](#), the DRBG is instantiated from either an RBG2(P) or an RBG3 construction. In order to instantiate the RBG1 construction at a security strength of s bits, this Recommendation requires the source RBG to support a security strength of at least s bits and provide a bitstring that is $3s/2$ bits long for most of the DRBGs. However, for a CTR_DRBG without a derivation function, a bitstring that is $s + 128$ bits long is required. (Note that an RBG3 construction supports any desired security strength.)

In the RBG2 and RBG3 constructions specified in [Sections 5](#) and [6](#), respectively, the DRBG within the construction is instantiated using a bitstring with a certain amount of entropy obtained from a validated entropy source.³⁴ In order to instantiate the DRBG to support an s -bit security strength, a bitstring with at least $3s/2$ bits of entropy is required for the instantiation of most of the DRBGs. Reseeding requires a bitstring with at least s bits of entropy. However, for a CTR_DRBG without a derivation function, a bitstring with exactly $s + 128$ full-entropy bits is required for instantiation and reseeding, either obtained directly from an entropy source that provides full-entropy output or from an entropy source via an **approved** (vetted) conditioning function (see [Section 3.3](#)).

The RBG3 constructions specified in [Section 6](#) are designed to provide full-entropy outputs but with a DRBG included in the design in case the entropy source fails undetectably. Entropy bits are possibly obtained from an entropy source via an **approved** (vetted) conditioning function. When the entropy source is working properly, an n -bit output from the RBG3 construction is said to provide n bits of entropy. The DRBG in an RBG3 construction is always required to support a 256-bit security strength. If an entropy-source fails and the failure is undetected, the RBG3 construction outputs are generated at a security strength of 256 bits. In this case, the security strength of a bitstring produced by the RBG is the minimum of 256 and its length (i.e., $security_strength = \min(256, length)$).

In conclusion, entropy sources and properly functioning RBG3 constructions provide output with entropy. RBG1 and RBG2 constructions provide output with a security strength that depends on the security strength of the RBG instantiation and the length of the output. Likewise, if the entropy source used by an RBG3 construction fails undetectably, the output is then dependent on the DRBG within the construction (an RBG(P) construction) to produce output at a security strength of 256 bits.

Because of the difference between the use of “entropy” to describe the output of an entropy source and the use of “security strength” to describe the output of a DRBG, the term “randomness” is used as a general term to mean either “entropy” or “security strength,” as appropriate. A “randomness source” is the general term for an entropy source or RBG that provides the randomness used by an RBG.

³⁴ However, note that the entropy-source output may be cryptographically processed by an **approved** conditioning function before being used.

Appendix B. RBG Examples (Informative)

[Appendix B.1](#) discusses and provides an example of the direct access to a DRBG used by an RBG3 construction.

Appendices [B.2](#) – [B.6](#) provide examples of each RBG construction. Not shown in the figures: if an error that indicates an RBG failure (e.g., a noise source in the entropy source has failed) is reported, RBG operation is terminated (see [Section 7.1.2](#)). For these examples, all entropy sources are considered to be physical entropy sources.

B.1. Direct DRBG Access in an RBG3 Construction

An implementation may be designed so that the DRBG implementation used within an RBG3 construction can be directly accessed by a consuming application³⁵ using the same or separate instantiations from the instantiation used by the RBG3 construction (see the examples in [Figure 19](#)).

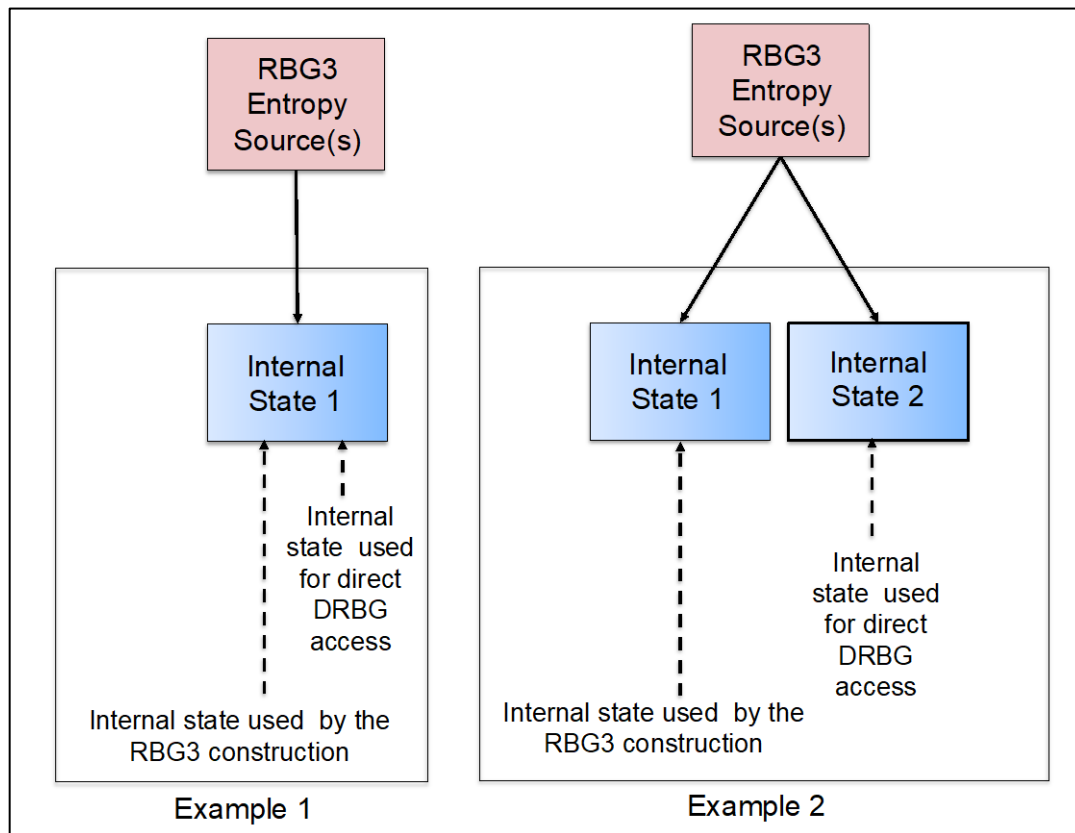


Fig. 19. DRBG Instantiations

In the leftmost example in Figure 19, the same internal state is used by the RBG3 construction and a directly accessible DRBG. The DRBG implementation is instantiated only once, and only a single state handle is obtained during instantiation (e.g., *RBG3_state handle*).³⁶ Generation and

³⁵ Without using other components or functionality used by the RBG3 construction (see Sections 6.2 and 6.3).

³⁶ Because only a single instantiation has been implemented, a state handle is not required.

2104 reseeding for RBG3 operations use RBG3 function calls (see Sections [6.2](#) and [6.3](#)), while
2105 generation and reseeding for direct DRBG access use RBG2 function calls (see [Section 5.2](#)) with
2106 the *RBG3_state_handle*. Using the same instantiation for both RBG3 operation and direct access
2107 to the DRBG requires additional reseeding processes in the case of an RBG3(RS) construction
2108 (see [Section 6.3.2](#)).

2109 In the rightmost example in [Figure 19](#), different internal states are used by the RBG3 construction
2110 and a directly accessible DRBG. The DRBG implementation is instantiated twice – once for RBG3
2111 operations and a second time for direct access to the DRBG. A different state handle needs to be
2112 obtained for each instantiation (e.g., *RBG3_state_handle* and *DRBG_state_handle*). Generation
2113 and reseeding for RBG3 operations use RBG3 function calls and *RBG3_state_handle* (see Sections
2114 [6.2](#) and [6.3](#)), while generation and reseeding for direct DRBG access use RBG2 function calls and
2115 *DRBG_state_handle* (see [Section 5.2](#)).

2116 Multiple directly accessible DRBGs may also be incorporated into an implementation by creating
2117 multiple instantiations. However, no more than one directly accessible DRBG should share the
2118 same internal state with the RBG3 construction (i.e., if n directly accessible DRBGs are required,
2119 either n or $n-1$ separate instantiations are required).

2120 The directly accessed DRBG instantiations are in the same security boundary as the RBG3
2121 construction. When accessed directly (rather than operating as part of the RBG3 construction), the
2122 DRBG instantiations are considered to be operating as RBG2(P) constructions as discussed in
2123 [Section 5](#).

2124 **B.2. Example of an RBG1 Construction**

2125 An RBG1 construction has access to a randomness source only during instantiation when it is
2126 seeded (see [Section 4](#)). For this example (see [Figure 20](#)), the DRBG used by the RBG1 construction
2127 and the randomness source reside in two different cryptographic modules with a secure channel
2128 connecting them during the instantiation process. Following DRBG instantiation, the secure
2129 channel is not available. For this example, the randomness source is an RBG2(P) construction (see
2130 [Section 5](#)) with a state handle of *RBG2_state_handle*.

2131 The targeted security strength for the RBG1 construction is 256 bits, so a DRBG from [[SP800-90A](#)]
2132 that is able to support this security strength must be used (HMAC_DRBG using SHA-256 is
2133 used in this example). A *personalization_string* is provided during instantiation, as recommended
2134 in [Section 2.4.1](#).

2135 As discussed in [Section 4](#), the randomness source (i.e., the RBG2(P) construction for this example)
2136 is not available during normal operation, so reseeding and prediction resistance cannot be
2137 provided.

2138 This example provides an RBG that is instantiated at a security strength of 256 bits.

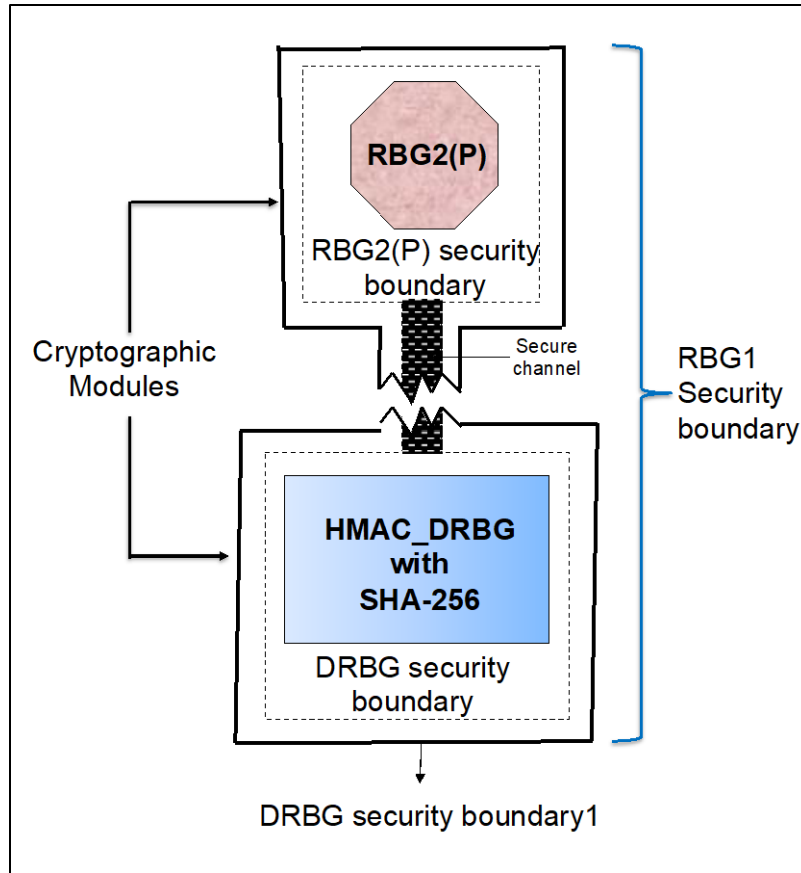


Fig. 20. RBG1 Construction Example

B.2.1. Instantiation of the RBG1 Construction

A physically secure channel is required to transport the entropy bits from the randomness source (the RBG2(P) construction) to the HMAC_DRBG during instantiation; an example of an RBG2(P) construction is provided in [Appendix B.4](#). Thereafter, the randomness source and the secure channel are no longer available.

The HMAC_DRBG is instantiated using the **Instantiate_function**, as specified in [Section 2.8.1.1](#), with the following call:

```
(status, RBG1_state_handle) = Instantiate_function (256, prediction_resistance_flag =
FALSE, "Device 7056").
```

A security strength of 256 bits is requested for the HMAC_DRBG used in the RBG1 construction.

Since an RBG1 construction does not provide prediction resistance (see [Section 4](#)), the *prediction_resistance_flag* is set to FALSE.

The *personalization string* to be used for this example is "Device 7056."

The **Get_randomness-source_input** call in the **Instantiate_function** results in a single request being sent to the randomness source to generate bits to establish the security strength (see [Section 4.2.1](#), item 2.a).

The HMAC_DRBG requests $3s/2 = 384$ bits from the randomness source, where s = the 256-bit targeted security strength for the DRBG:

$(status, randomness_bitstring) = \text{Generate_function}(RBG2_state_handle, 384, 256, prediction_resistance_request = \text{TRUE}).$

This call requests the randomness source (indicated by *RBG2_state_handle*) to generate 384 bits at a security strength of 256 bits for the randomness input required for seeding the DRBG in the RBG1 construction. Prediction resistance is requested so that the randomness source (i.e., the RBG2(P) construction) is reseeded before generating the requested 384 bits (see Requirement 17 in [Section 4.4.1](#)). Note that optional *additional_input* is not provided for this example.

2. The RBG2(P) construction checks that the request can be handled (e.g., whether a security strength of 256 bits is supported). If the request is valid, 384 bits are generated after reseeding the RBG2(P) construction, the internal state of the RBG2(P) construction is updated, and *status* = SUCCESS is returned to the RBG1 construction along with the newly generated *randomness_bitstring*.

If the request is determined to be invalid, *status* = FAILURE is returned along with a *Null* bitstring as the *randomness_bitstring*. The FAILURE *status* is subsequently returned from the **Instantiate_function** along with a *Null* value as the *RBG1_state_handle*, and the instantiation process is terminated.

If a valid *randomness_bitstring* is returned from the RBG2(P) construction, the *randomness_bitstring* is used along with the *personalization_string* to create the seed to instantiate the DRBG (see [\[SP800-90A\]](#)).³⁷ If the instantiation is successful, the internal state is established, a *status* of SUCCESS is returned from the **Instantiate_function** with a state handle of *RBG1_state_handle*, and the RBG can be used to generate pseudorandom bits.

B.2.2. Generation by the RBG1 Construction

Assuming that the HMAC_DRBG in the RBG1 construction has been instantiated (see [Appendix B.2.1](#)), pseudorandom bits are requested from the RBG by a consuming application using the **Generate_function** call as specified in [Section 2.8.1.2](#):

$(status, returned_bits) = \text{Generate_function}(RBG1_state_handle, requested_number_of_bits, requested_security_strength, prediction_resistance_request = \text{FALSE}, additional_input).$

RBG1_state_handle was returned as the state handle during instantiation (see [Appendix B.2.1](#)).

³⁷ The first 256 bits of the *randomness_bitstring* are used as the randomness input, and the remaining 128 bits are used as the nonce in SP 800-90A, Revision 1. A future update of SP 800-90A will revise this process by using the entire 384-bit string as the randomness input.

The *requested_security_strength* may be any value that is less than or equal to 256 (the instantiated security strength recorded in the DRBG's internal state).

Since prediction resistance cannot be provided in an RBG1 construction, *prediction_resistance_request* is set to FALSE. (Note that the *prediction_resistance_request* input parameter could be omitted from the **Generate_function** call for this example).

Any *additional_input* is optional.

The **Generate_function** returns an indication of the *status*. If *status* = SUCCESS, the *requested_number_of_bits* are provided as the *returned_bits* to the consuming application. If *status* = FAILURE, *returned_bits* is an empty (i.e., null) bitstring.

B.3. Example Using Sub-DRBGs Based on an RBG1 Construction

This example uses an RBG1 construction to instantiate two sub-DRBGs: sub-DRBG1 and sub-DRBG2 (see [Figure 21](#)).

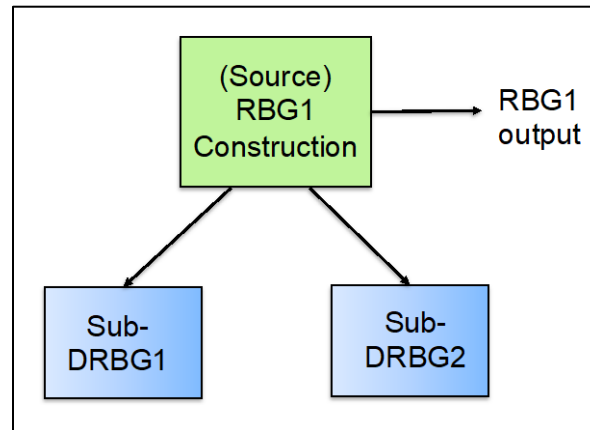


Fig. 21. Sub-DRBGs Based on an RBG1 Construction

The instantiation of the RBG1 construction is discussed in Appendix B.2. The RBG1 construction that is used as the source RBG includes an HMAC_DRBG and has been instantiated to provide a security strength of 256 bits. The state handle for the construction is *RBG1_state_handle*.

For this example, Sub-DRBG1 will be instantiated to provide a security strength of 128 bits, and Sub-DRBG2 will be instantiated to provide a security strength of 256 bits. Both sub-DRBGs use the same DRBG algorithm as the RBG1 construction.

Neither the RBG1 construction nor the sub-DRBGs can be reseeded or provide prediction resistance.

This example provides the following capabilities:

- Access to the RBG1 construction to provide output generated at a security strength of 256 bits (see Appendix B.2 for the RBG1 example)
- Access to one sub-DRBG (Sub-DRBG1) that provides output for an application that requires a security strength of no more than 128 bits

- Access to a second sub-DRBG (Sub-DRBG2) that provides output for a second application that requires a security strength of 256 bits

2221 B.3.1. Instantiation of the Sub-DRBGs

2222 Each sub-DRBG is instantiated using output from an RBG1 construction that is discussed in
2223 Appendix 62B.2.

2224 B.3.1.1. Instantiating Sub-DRBG1

2225 Sub-DRBG1 is instantiated using the following **Instantiate_function** call (see [Section 2.8.1.1](#)):

2226 $(status, sub-DRBG1_state_handle) = \text{Instantiate_function}(128, prediction_resistance_flag$
2227 $= \text{FALSE}, \text{"Sub-DRBG App 1"}).$

- A security strength of 128 bits is requested from the DRBG indicated by the *RBG1_state_handle*.
- Setting "*prediction_resistance_flag* = FALSE" indicates that a consuming application will not be allowed to request prediction resistance. Optionally, the parameter can be omitted.
- The *personalization_string* to be used for sub-DRBG1 is "Sub-DRBG App 1."
- The returned state handle for sub-DRBG1 will be *sub-DRBG1_state_handle*.

2234 The randomness input for establishing the 128-bit security strength of sub-DRBG1 is requested
2235 using the following **Generate_function** call to the RBG1 construction):

2236 $(status, randomness-source_input) = \text{Generate_function}(RBG1_state_handle, 192, 128,$
2237 $prediction_resistance_request = \text{FALSE}, additional_input).$

- 192 bits are requested from the source RBG (indicated by *RBG1_state_handle*) at a security strength of 128 bits ($192 = 128 + 64 = 3s/2$).
- Setting "*prediction_resistance_flag* = FALSE" indicates that the source RBG (the RBG1 construction) will not need to reseed itself before generating the requested output. Alternatively, the parameter can be omitted.
- Additional input is optional.

2244 If *status* = SUCCESS is returned from the **Generate_function**, the HMAC_DRBG in sub-DRBG1
2245 is seeded using the *randomness-source_input* obtained from the RBG1 construction and the
2246 *personalization_string* provided in the **Instantiate_function** call (i.e., "Sub-DRBG App 1"). The
2247 internal state is recorded for Sub-DRBG1 (including the 128-bit security strength), and *status* =
2248 SUCCESS is returned from the **Instantiate_function** along with a state handle of *sub-*
2249 *DRBG1_state_handle*.

2250 If *status* = FAILURE is returned from the **Generate_function** call, then the internal state is not
2251 created, *status* = FAILURE and a Null state handle are returned from the **Instantiate_function**,
2252 and the sub-DRBG1 cannot be used to generate bits.

2253 **B.3.1.2. Instantiating Sub-DRBG2**

2254 Sub-DRBG2 is instantiated using the following **Instantiate_function** call (see [Section 2.8.1.1](#)):

2255 $(status, sub\text{-}DRBG2_state_handle) = \text{Instantiate_function}(256, prediction_resistance_flag =$
2256 $FALSE, \text{“Sub-DRBG App 2”})$.

- 2257 • A security strength of 256 bits is requested from the randomness source (the DRBG
2258 construction indicated by *RBG1_state_handle*).
- 2259 • Setting “*prediction_resistance_flag* = FALSE” indicates that a consuming application will
2260 not be allowed to request prediction resistance. Optionally, the parameter can be omitted.
- 2261 • The *personalization_string* to be used for sub-DRBG2 is “Sub-DRBG App 2.”
- 2262 • The returned state handle will be *sub-DRBG2_state_handle*.

2263 The randomness input for establishing the 256-bit security strength of sub-DRBG2 is requested
2264 using the following **Generate_function** call to the RBG1 construction):

2265 $(status, randomness\text{-}source_input) = \text{Generate_function}(RBG1_state_handle, 384, 256,$
2266 $prediction_resistance_request = FALSE, additional_input)$.

- 2267 • 384 bits are requested from the source RBG (indicated by *RBG1_state_handle*) at a security
2268 strength of 256 bits ($384 = 256 + 128 = 3s/2$).
- 2269 • Setting “*prediction_resistance_flag* = FALSE” indicates that the source RBG (the RBG1
2270 construction) will not need to reseed itself before generating the requested output.
2271 Alternatively, the parameter can be omitted.
- 2272 • Additional input is optional.

2273 If *status* = SUCCESS is returned from the **Generate_function**, the HMAC_DRBG in sub-DRBG2
2274 is seeded using the *randomness-source_input* obtained from the RBG1 construction and the
2275 *personalization_string* provided in the **Instantiate_function** call (i.e., “Sub-DRBG App 2”). The
2276 internal state is recorded for Sub-DRBG2 (including the 256-bit security strength), and *status* =
2277 SUCCESS is returned from the **Instantiate_function** along with a state handle of *sub-*
2278 *DRBG2_state_handle*.

2279 If *status* = FAILURE is returned from the **Generate_function** call, then the internal state is not
2280 created, *status* = FAILURE and a Null state handle are returned from the **Instantiate_function**,
2281 and the sub-DRBG2 cannot be used to generate bits.

2282 **B.3.2. Pseudorandom Bit Generation by Sub-DRBGs**

2283 Assuming that the sub-DRBG has been successfully instantiated (see [Appendix B.3.1](#)),
2284 pseudorandom bits are requested from the sub-DRBG by a consuming application using the
2285 **Generate_function** call as specified in [Section 2.8.1.2](#):

2286 $(status, returned_bits) = \text{Generate_function}(state_handle, requested_number_of_bits,$
2287 $security_strength, prediction_resistance_request, additional_input)$,

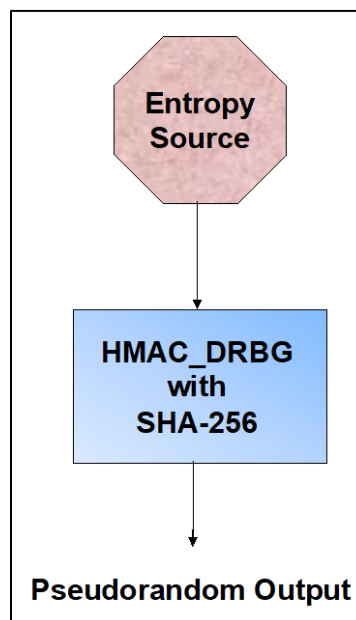
2288 where:

- 2289 • For sub_DRBG1, *state_handle* = *sub-DRBG1_state_handle*;

- 2290 For sub-DRBG2, *state_handle* = *sub-DRBG2_state_handle*;
- 2291 • *requested_number_of_bits* must be $\leq 2^{19}$ (see SP 800-90A for HMAC_DRBG);
- 2292 • For *sub_DRBG1*, *security_strength* must be ≤ 128 ;
- 2293 • For *sub_DRBG2*, *security_strength* must be ≤ 256 ;
- 2294 • *prediction_resistance_request* = FALSE (or is omitted); and
- 2295 • *additional_input* is optional.

2296 B.4. Example of an RBG2(P) or RBG2(NP) Construction

2297 For this example of an RBG2 construction, no conditioning function is used, and only a single
2298 DRBG instantiation will be used (see [Figure 22](#)), so a state handle is not needed. Full-entropy
2299 output is not provided by the entropy source, which may be either a physical or non-physical
2300 entropy source.



2302 Fig. 22. RBG2 Example

2303 The targeted security strength is 256 bits, so a DRBG from [\[SP800-90A\]](#) that can support this
2304 security strength must be used; HMAC_DRBG using SHA-256 is used in this example. A
2305 *personalization_string* may be provided, as recommended in [Section 2.4.1](#). Reseeding and
2306 prediction resistance are supported and will be available on demand.

2307 This example provides the following capabilities:

- 2308 • An RBG instantiated at a security strength of 256 bits, and
- 2309 • Access to an entropy source to provide prediction resistance.

B.4.1. Instantiation of an RBG2 Construction

The DRBG in the RBG2 construction is instantiated using an **Instantiate_function** call (see [Section 2.8.1.1](#)):

(status) = **Instantiate_function** (256, *prediction_resistance_flag* = TRUE, “RBG2 42”).

- Since there is only a single instantiation, a *state_handle* is not used for this example.
- Using “*prediction_resistance_flag* = TRUE”, the RBG is notified that prediction resistance may be requested in subsequent **Generate_function** calls.
- The *personalization_string* to be used for this example is “RBG2 42.”

The entropy for establishing the security strength (*s*) of the DRBG (i.e., *s* = 256 bits) is requested using the following **Get_ES_Bitstring** call to the entropy source (see [Section 2.8.2.2](#) and item 2 in [Section 5.2.1](#)):

(status, entropy_bitstring) = **Get_ES_Bitstring**(384),

where $3s/2 = 384$ bits of entropy are requested from the entropy source.

If *status* = SUCCESS is returned from the **Get_ES_Bitstring** call, the HMAC_DRBG is seeded using *entropy_bitstring*, and the *personalization_string* is “RBG2 42.” The internal state is recorded (including the security strength of the instantiation), and *status* = SUCCESS is returned to the consuming application by the **Instantiate_function**.

If *status* = FAILURE is returned from the **Get_ES_Bitstring** call, then the internal state is not created, *status* = FAILURE and a Null state handle are returned by the **Instantiate_function** to the consuming application, and the RBG cannot be used to generate bits.

B.4.2. Generation in an RBG2 Construction

Assuming that the RBG has been successfully instantiated (see [Appendix B.4.1](#)), pseudorandom bits are requested from the RBG by a consuming application using the **Generate_function** call as specified in [Section 2.8.1.2](#):

(status, returned_bits) = **Generate_function**(*requested_number_of_bits*, *security_strength*, *prediction_resistance_request*, *additional_input*).

- Since there is only a single instantiation of the HMAC_DRBG, a *state_handle* was not returned from the **Instantiate_function** (see [Appendix B.4.1](#)) and is not used during the **Generate_function** call.
- The *requested_security_strength* may be any value that is less than or equal to 256 (the instantiated security strength recorded in the HMAC_DRBG’s internal state).
- *prediction_resistance_request* = TRUE if prediction resistance is requested and FALSE otherwise.
- Additional input is optional.

If prediction resistance is requested, a reseed of the HMAC_DRBG is requested by the **Generate_function** before the requested bits are generated (see [Appendix B.4](#)). If *status* =

FAILURE is returned from the **Reseed_function**, *status* = FAILURE is also returned to the consuming application by the **Generate_function**, along with a Null value as the *returned_bits*.

Whether or not prediction resistance is requested, a *status* indication is returned from the **Generate_function** call. If *status* = SUCCESS, a bitstring of at least *requested_number_of_bits* is provided as the *returned_bits* to the consuming application. If *status* = FAILURE, *returned_bits* is an empty bitstring.

B.4.3. Reseeding an RBG2 Construction

The HMAC_DRBG will be reseeded 1) if explicitly requested by the consuming application, 2) whenever generation with prediction resistance is requested by the **Generate_function**, or 3) automatically during a **Generate_function** call at the end of the DRBG's designed *seedlife* (see the **Generate_function** specification in [SP800-90A]).

The **Reseed_function** call, as specified in [Section 2.8.1.3](#), is:

status = **Reseed_function**(*additional_input*).

- Since there is only a single instantiation of the HMAC_DRBG, a *state_handle* was not returned from the **Instantiate_function** (see [Appendix B.4.1](#)) and is not used during the **Reseed_function** call.
- The *additional_input* is optional.

Since entropy is obtained directly from the entropy source (case 2 in [Section 5.2.3](#)), the implementation has replaced the **Get_randomness-source_input** call used by the **Reseed_function** in [SP800-90A] with a **Get_ES_Bitstring** call.

The HMAC_DRBG is reseeded with a security strength of 256 bits as follows:

(*status*, *entropy_bitstring*) = **Get_ES_Bitstring**(256).

If *status* = SUCCESS is returned by **Get_ES_Bitstring**, the *entropy_bitstring* contains at least 256 bits of entropy and is at least 256 bits long. *Status* = SUCCESS is returned to the calling application (e.g., the **Generate_function**) by the **Reseed_function**.

If *status* = FAILURE, *entropy_bitstring* is an empty (e.g., null) bitstring. The HMAC_DRBG is not reseeded, and *status* = FAILURE is returned from **Reseed_function** to the calling application.

B.5. Example of an RBG3(XOR) Construction

This construction is specified in [Section 6.2](#) and requires a DRBG and a source of full-entropy bits. For this example, the entropy source itself does not provide full-entropy output, so the vetted Hash conditioning function listed in [SP800-90B] using SHA-256 is used as an external conditioning function.

The Hash_DRBG specified in [SP800-90A] will be used as the DRBG, with SHA-256 used as the underlying hash function for the DRBG (note the use of SHA-256 for both the Hash_DRBG and the vetted conditioning function). The DRBG will obtain input directly from the RBG's entropy source without conditioning (as shown in [Figure 23](#)), since bits with full entropy are not required

2382 for input to the DRBG, even though full-entropy bits are required for input to the XOR operation
2383 (shown as “ \oplus ” in the figure) from the entropy source via the conditioning function.

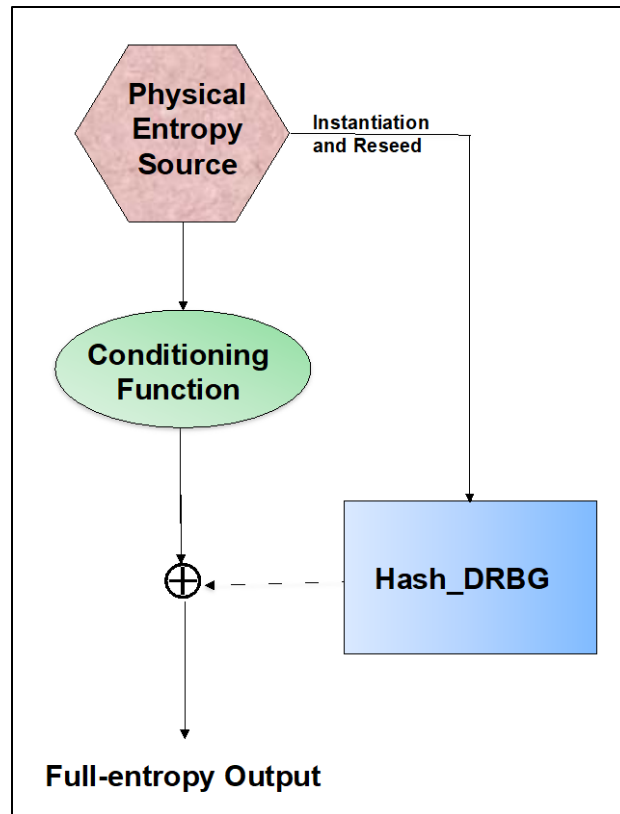


Fig. 23. RBG3(XOR) Construction Example

2386 As specified in [Section 6.2](#), the DRBG must be instantiated (and reseeded) at 256 bits, which is
2387 possible for SHA-256.

2388 In this example, only a single instantiation is used, and a personalization string is provided during
2389 instantiation. The DRBG is not directly accessible.

2390 Calls are made to the RBG using the RBG3(XOR) calls specified in [Section 6.2](#).

2391 The Hash_DRBG itself is not directly accessible.

2392 This example provides the following capabilities:

- 2393 • Full-entropy output by the RBG,
- 2394 • Fallback to the security strength provided by the Hash_DRBG (256 bits) if the entropy
2395 source has an undetected failure, and
- 2396 • Access to an entropy source to instantiate and reseed the Hash_DRBG.

2397 B.5.1. Instantiation of an RBG3(XOR) Construction

2398 The Hash_DRBG is instantiated using:

2399 `status = RBG3(XOR)_DRBG_Instantiate("RBG3(XOR)",`

- The personalization string for the DRBG is “RBG3(XOR).”

```

2405 (status) = Instantiate_function(256, prediction_resistance_flag = FALSE,
2406                                personalization_string).

```

2409 • The DRBG is instantiated at a security strength of 256 bits.

2414 • The personalization string for the DRBG is “RBG3(XOR).” It was provided in the
2415 **RBG3(XOR) DRBG Instantiate** call.

2417 The entropy for establishing the security strength (s) of the Hash_DRBG (i.e., where $s = 256$ bits)
2418 is requested using the following **Get ES Bitstring** call:

2420 where $3s/2 = 384$ bits of entropy are requested from the entropy source.

2426 If *status* = FAILURE is returned from the **Get_ES_Bitstring** call, *status* = FAILURE and a Null
2427 state handle are returned to the consuming application from the **Instantiate_function**. 'The
2428 Hash DRBG's internal state is not established, and the RBG cannot be used to generate bits.

2430 Assuming that the Hash_DRBG has been instantiated (see [Appendix B.4.1](#)), the RBG can be called
2431 by a consuming application to generate output with full entropy.

2432 **B.5.2.1. Generation**

2433 Let n indicate the requested number of bits to generate. The construction in [Section 6.3.1.2](#) is used
2434 as follows:

2435 **RBG3(XOR)_Generate:**

2436 **Input:** integer n , string *additional_input*.

2437 **Output:** integer *status*, bitstring *returned_bits*.

2438 **Process:**

- 2439 1. $(status, ES_bits) = \text{Get_conditioned_full-entropy_input}(n)$.
- 2440 2. If $(status \neq \text{SUCCESS})$, then return($status, \text{Null}$).
- 2441 3. $(status, DRBG_bits) = \text{Generate_function}(n, 256, prediction_resistance_request =$
2442 $\text{FALSE}, additional_input)$.
- 2443 4. If $(status \neq \text{SUCCESS})$, then return($status, \text{Null}$).
- 2444 5. $returned_bits = ES_bits \oplus DRBG_bits$.
- 2445 6. Return $\text{SUCCESS}, returned_bits$.

2446 Note that the *state_handle* parameter is not used in the **RBG3(XOR)_Generate** call or the
2447 **Generate_function** call (in step 3) for this example since a *state_handle* was not returned from
2448 the **RBG3(XOR)_DRBG_Instantiate** function (see [Appendix B.5.1](#)).

2449 In step 1, the entropy source is accessed via the conditioning function using the
2450 **Get_conditioned_full-entropy_input** routine (see [Appendix B.5.2.2](#)) to obtain n bits with full
2451 entropy.

2452 Step 2 checks that the **Get_conditioned_full-entropy_input** call in step 1 was successful. If it
2453 was not successful, the **RBG3(XOR)_Generate** function is aborted, returning $status \neq \text{SUCCESS}$
2454 to the consuming application along with a *Null* bitstring as the *returned_bits*.

2455 Step 3 calls the Hash_DRBG to generate n bits to be XORed with the n -bit output of the entropy
2456 source (*ES_Bits*; see step 1) in order to produce the RBG output. Note that a request for prediction
2457 resistance is not made in the **Generate_function** call (i.e., *prediction_resistance_request* =
2458 *FALSE*). Optionally, this parameter could be omitted since prediction resistance is never
2459 requested.

2460 Step 4 checks that the **Generate_function** invoked in step 3 was successful. If it was not
2461 successful, the **RBG3(XOR)_Generate** function is aborted, returning $status \neq \text{SUCCESS}$ to the
2462 consuming application along with a *Null* bitstring as the *returned_bits*.

2463 If step 3 returns an indication of success, the *ES_bits* returned in step 1 and the *DRBG_bits* obtained
2464 in step 3 are XORed together in step 5. The result is returned to the consuming application in step
2465 6.

2466 **B.5.2.2. Get_conditioned_full-entropy_input Function**

2467 The **Get_conditioned_full-entropy_input** construction is specified in [Section 3.3.2](#). For this
2468 example, the routine becomes the following:

2469 **Get_conditioned_full_entropy_input:**

2470 **Input:** integer n .

2471 **Output:** integer $status$, bitstring $Full\text{-}entropy_bitstring$.

2472 **Process:**

- 2473 1. $temp$ = the *Null* string.
- 2474 2. $ctr = 0$.
- 2475 3. While $ctr < n$, do
 - 2476 3.1 $(status, entropy_bitstring) = \text{Get_ES_Bitstring}(320)$.
 - 2477 3.2 If $(status \neq \text{SUCCESS})$, then return $(status, invalid_string)$.
 - 2478 3.3 $conditioned_output = \text{Hash}_{\text{SHA-256}}(entropy_bitstring)$.
 - 2479 3.4 $temp = temp \parallel conditioned_output$.
 - 2480 3.5 $ctr = ctr + 256$.
- 2481 4. $Full\text{-}entropy_bitstring = \text{leftmost}(temp, n)$.
- 2482 5. Return $(\text{SUCCESS}, Full\text{-}entropy_bitstring)$.

2483 Steps 1 and 2 initialize the temporary bitstring ($temp$) for holding the full-entropy bitstring being
2484 assembled, and the counter (ctr) that counts the number of full-entropy bits produced so far.

2485 Step 3 obtains and processes the entropy for each iteration.

- 2486 • Step 3.1 requests 320 bits from the entropy source(s) (i.e., $output_len + 64$ bits, where
2487 $output_len = 256$ for SHA-256).
- 2488 • Step 3.2 checks whether or not the $status$ returned in step 3.1 indicated a success. If the
2489 $status$ did not indicate a success, the $status$ is returned along with an invalid (e.g., *Null*)
2490 bitstring as the $Full\text{-}entropy_bitstring$.
- 2491 • Step 3.3 invokes the Hash conditioning function (see [Section 3.3.1.2](#)) using SHA-256 for
2492 processing the $entropy_bitstring$ obtained from step 3.1.
- 2493 • Step 3.4 concatenates the $conditioned_output$ received in step 3.3 to the temporary bitstring
2494 ($temp$), and step 3.5 increments the counter for the number of full-entropy bits that have
2495 been produced so far.

2496 After at least n bits have been produced in step 3, step 4 selects the leftmost n bits of the temporary
2497 string ($temp$) to be returned as the bitstring with full entropy.

2498 Step 5 returns the result from step 4 ($Full\text{-}entropy_bitstring$).

2499 **B.5.3. Reseeding an RBG3(XOR) Construction**

2500 The Hash_DRBG must be reseeded at the end of its designed seedlife and may be reseeded on
2501 demand (e.g., by the consuming application). Reseeding will be automatic whenever the end of
2502 the DRBG's seedlife is reached during a **Generate_function** call (see [SP800-90A]). For this
2503 example, whether reseeded is done automatically during a **Generate_function** call or is
2504 specifically requested by a consuming application, the **Reseed_function** call is:

2505 *status* = **Reseed_function**(*additional_input*).

- 2506 • The *state_handle* parameter is not used in the **Reseed_function** call since a *state_handle*
2507 was not returned from the **RBG3(XOR)_DRBG_Instantiate** function (see [Appendix](#)
2508 [B.5.1](#)).
- 2509 • The security strength for reseeded the Hash_DRBG is recorded in the internal state as 256
2510 bits.
- 2511 • Additional input is optional.

2512 [Section 6.3.1.3](#) refers to [Section 5.2.3](#) for reseeded the Hash_DRBG. Since entropy is obtained
2513 directly from the entropy source and no conditioning function is used (case 2 in [Section 6.3.2](#)), the
2514 implementation has replaced the **Get_randomness-source_input** call used by the
2515 **Reseed_function** in [SP800-90A] with a **Get_ES_Bitstring** call.

2516 The Hash_DRBG is reseeded with a security strength of 256 bits as follows:

2517 (*status*, *entropy_bitstring*) = **Get_ES_Bitstring**(256).

2518 If *status* = SUCCESS is returned by the **Get_ES_Bitstring** call, *entropy_bitstring* consists of at
2519 least 256 bits that contain at least 256 bits of entropy. These bits are used to reseed the
2520 Hash_DRBG. *Status* = SUCCESS is then returned to the calling application by the
2521 **Reseed_function**.

2522 If *status* = FAILURE, *entropy_bitstring* is an empty (e.g., null) bitstring. The Hash_DRBG is not
2523 reseeded, and *status* ≠ SUCCESS is returned from the **Reseed_function** to the calling application
2524 (e.g., the **Generate_function**).

2525 **B.6. Example of an RBG3(RS) Construction**

2526 This construction is specified in [Section 6.3](#) and requires an entropy source and a DRBG (see the
2527 left half of [Figure 24](#) outlined in green). The DRBG is directly accessible using the same
2528 instantiation that is used by the RBG3(RS) construction (i.e., they share the same internal state).
2529 When accessed directly, the DRBG behaves as an RBG2(P) construction (see the right half of
2530 [Figure 24](#) outlined in blue).

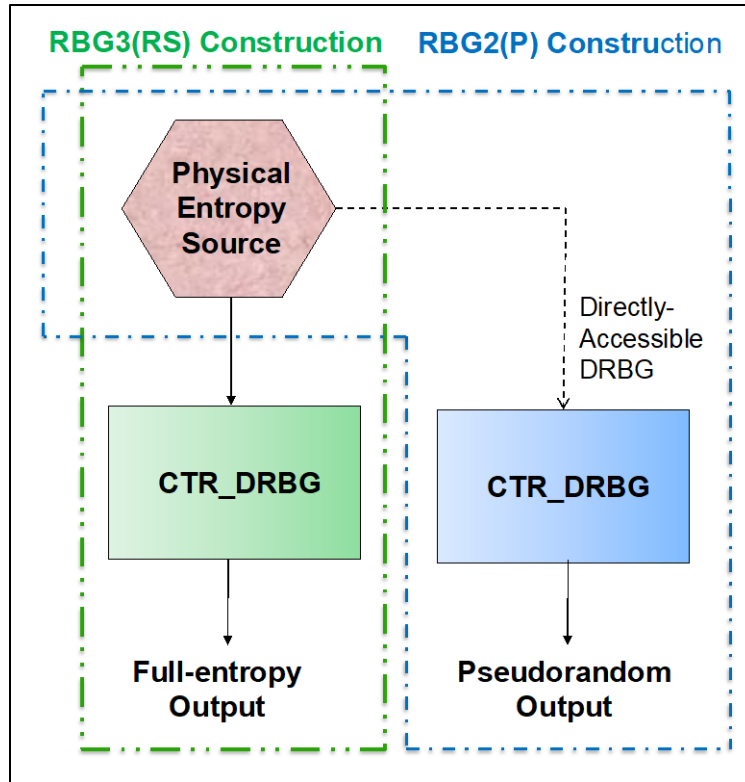


Fig. 24. RBG3(RS) Construction Example

The CTR_DRBG specified in [SP800-90A] will be used as the DRBG with AES-256 used as the underlying block cipher for the DRBG. The CTR_DRBG will be implemented using a derivation function (located inside the CTR_DRBG implementation). In this case, full-entropy output will not be required for the entropy source (see [SP800-90A]). However, an alternative example could use the CTR_DRBG without a derivation function. In that case, either the entropy source would need to provide full-entropy output, or a vetted conditioning function would be required to condition the entropy to provide full-entropy bits before providing it to the DRBG.

As specified in Section 6.2, a DRBG used as part of the RBG must be instantiated (and reseeded) at a security strength of 256 bits (which AES-256 can support).

For this example, the DRBG has a fixed security strength (256 bits), which is hard-coded into the implementation so will not be used as an input parameter.

Calls are made to the RBG as specified in Section 6.3.1. Calls made to the directly accessible DRBG (part of a RBG2(P) construction) use the RBG calls specified in Section 5.2. Since an entropy source is always available, the directly accessed DRBG can be reseeded and support prediction resistance.

If the entropy source produces output at a slow rate, a consuming application might call the RBG3(RS) construction only when full-entropy bits are required, obtaining all other output from the directly accessible DRBG.

This example provides the following capabilities:

- 2552 • Full-entropy output by the RBG3(RS) construction,
- 2553 • Fallback to the security strength of the RBG3(RS)'s DRBG instantiation (256 bits) if the
- 2554 entropy source has an undetected failure,
- 2555 • Direct access to an RBG2(P) construction with a security strength of 256 bits for faster
- 2556 output when full-entropy output is not required,
- 2557 • Access to an entropy source to instantiate and reseed the DRBG, and
- 2558 • Prediction resistance support for the directly accessed DRBG.

2559 **B.6.1. Instantiation of an RBG3(RS) Construction**

2560 Instantiation for this example consists of the instantiation of the CTR_DRBG used by the
2561 RBG3(RS) construction.

2562 The DRBG is initialized as follows:

2563 $(status, RBG3(RS)_{state_handle}) = \mathbf{RBG3(RS)_DRBG_Instantiate}(\text{"RBG3(RS) 2021"})$.

- 2564 • "RBG3(RS) 2021" is to be used as the personalization string for the DRBG instantiation
- 2565 used in the RBG3(RS) construction.
- 2566 • $RBG3(RS)_{state_handle}$ is returned as the state handle for the DRBG instantiation used
- 2567 by the RBG3(RS) construction.

2568 Appendices [B.6.2](#) and [B.6.3](#) will show the differences between the operation of the RBG3(RS)
2569 and RBG2(P) constructions.

2570 **B.6.2. Generation by an RBG3(RS) Construction**

2571 Assuming that the DRBG instantiation for the RBG3(RS) construction has been instantiated (see
2572 [Appendix B.6.1](#)), the RBG can be invoked by a consuming application to generate outputs with
2573 full entropy. The **RBG3(RS)_Generate** construction in [Section 6.3.1.2.1](#) is invoked using

2574 $(status, returned_bits) = \mathbf{RBG3(RS)_Generate}(RBG3(RS)_{state_handle}, n,$
2575 $additional_information)$.

- 2576 • The $RBG3(RS)_{state_handle}$ (obtained during instantiation; see [Appendix B.6.1](#)) is used
- 2577 to access the internal state information for the DRBG instantiation for the RBG3(RS)
- 2578 construction.
- 2579 • The consuming application requests n bits.
- 2580 • The input of $additional_information$ is optional.

2581 The process is specified in [Section 6.3.1.2.1](#). The state handle in the **Generate_function** is
2582 $RBG3(RS)_{state_handle}$, which was obtained during instantiation (see [Appendix B.6.1](#)).

2583 **B.6.3. Generation by the Directly Accessible DRBG**

2584 Assuming that the DRBG has been instantiated (see [Appendix B.6.1](#)), it can be accessed directly
2585 by a consuming application in the same manner as the RBG2(P) example in [Appendix B.4.2](#) using

2586 the *RBG3(RS)_state_handle* obtained during instantiation (see [Appendix B.6.1](#)) and using a
2587 **Generate_function** call:

2588 (*status, returned_bits*) = **Generate_function**(*RBG3(RS)_state_handle*, *n*,
2589 *prediction_resistance_request, additional_input*).

2590 Note that the security strength parameter (256) was omitted since its value has been hard coded.

2591 Requirement 2 in [Section 6.3.2](#) requires that the DRBG be reseeded whenever a request for
2592 generation by a directly accessible DRBG follows a request for generation by the RBG3(RS)
2593 construction. For this example, the internal state includes an indication about whether the last use
2594 of the DRBG was as part of the RBG3(RS) construction or was directly accessible. If the
2595 **Generate_function** (above) does not include a request for prediction resistance (e.g.,
2596 *prediction_resistance_request* was not set to TRUE), then the DRBG will be reseeded anyway
2597 using the entropy source before generating output if the previous use of the DRBG was part of the
2598 RBG3(RS) construction.

2599 **B.6.4. Reseeding a DRBG**

2600 When operating as part of the RBG3(RS) construction, the **Reseed_function** is invoked one or
2601 more times to produce full-entropy output when the **RBG3(RS)_Generate** function is invoked by
2602 a consuming application.

2603 When operating as part of the RBG2(P) construction (the directly accessible DRBG), the DRBG
2604 is reseeded 1) if explicitly requested by the consuming application, 2) automatically whenever a
2605 generation with prediction resistance is requested during a direct access of the DRBG (see
2606 [Appendix B.6.3](#)), 3) whenever the previous use of the DRBG was by the **RBG3(RS)_Generate**
2607 function (see [Appendix B.6.2](#)), or 4) automatically during a **Generate_function** call at the end of
2608 the seedlife of the RBG2(P) construction (see the **Generate_function** specification in [[SP800-90A](#)]).
2609

2610 The **Reseed_function** call is:

2611 *status* = **Reseed_function**(*RBG3(RS)_state_handle, additional_input*).

- 2612 • The *state_handle* is *RBG3(RS)_state_handle*, and
- 2613 • *additional_input* is optional.³⁸

2614 The DRBG is reseeded with a security strength of 256 bits as follows:

2615 (*status, entropy_bitstring*) = **Get_ES_Bitstring**(256).

2616 If *status* = SUCCESS is returned by **Get_ES_Bitstring**, *entropy_bitstring* consists of at least 256
2617 bits containing at least 256 bits of entropy. *Status* = SUCCESS is returned to the calling application
2618 by the **Reseed_function**.

³⁸ Note that when the **RBG3(RS)_Generate** function uses a Hash_DRBG, HMAC_DRBG, or CTR_DRBG with no derivation function and Method A, whereby 64 bits of additional entropy are required to produce *output_len* bits with full entropy (see Section 7.3.1, 2.1, step 3.1), the additional 64 bits of entropy obtained in step 3.1.1 is provided to the **Generate_function** (in step 3.1.3) with prediction requested. In Section 9.3 of SP 800-90A, the **Generate_function** reseeds the DRBG when prediction resistance is requested using entropy from the entropy source and any additional input that is provided – the additional 64 bits of entropy, in this case.

2619 If *status* ≠ SUCCESS (e.g., the entropy source has failed), *entropy_bitstring* is an empty (e.g., null)
2620 bitstring, the DRBG is not reseeded, and a FAILURE *status* is returned from **Reseed_function** to
2621 the calling application (e.g., the **Generate_function**).
2622

2623 **Appendix C. Addendum to SP 800-90A: Instantiating and Reseeding a CTR_DRBG**

2624 **C.1. Background and Scope**

2625 The CTR_DRBG, specified in [SP800-90A], uses the block cipher AES and has two versions that
2626 may be implemented: with or without a derivation function.

2627 When a derivation function is not used, SP 800-90A requires the use of bitstrings with full entropy
2628 for instantiating and reseeding a CTR_DRBG. This addendum permits the use of an RBG
2629 compliant with SP 800-90C to provide the required seed material for the CTR_DRBG when
2630 implemented as specified in SP 800-90C (see [Appendix C.2](#)).

2631 When a derivation function is used in a CTR_DRBG implementation, SP 800-90A specifies the
2632 use of the block cipher derivation function. This addendum modifies the requirements in SP 800-
2633 90A for the CTR_DRBG by specifying two additional derivation functions that may be used
2634 instead of the block cipher derivation function (see [Appendix C.3](#)).

2635 **C.2. CTR_DRBG without a Derivation Function**

2636 When a derivation function is not used, SP 800-90A requires that *seedlen* full-entropy bits be
2637 provided as the randomness input (e.g., from an entropy source that provides full-entropy output),
2638 where *seedlen* is the length of the key to be used by the CTR_DRBG plus the length of the output
2639 block.³⁹ SP 800-90C includes an approved method for externally conditioning the output of an
2640 entropy source to provide a bitstring with full entropy when using an entropy source that does not
2641 provide full-entropy output.

2642 SP 800-90C also permits the use of seed material from an RBG when the DRBG to be instantiated
2643 and reseeded is implemented and used as specified in SP 800-90C.

2644 **C.3. CTR_DRBG using a Derivation Function**

2645 When a derivation function is used within a CTR_DRBG, SP 800-90A specifies the use of the
2646 **Block_cipher_df** included in that document during instantiation and reseeding to adjust the length
2647 of the seed material to *seedlen* bits, where

2648
$$seedlen = \text{the security strength} + \text{the block length.}$$

2649 For AES, *seedlen* = 256, 320 or 384 bits (see [SP800-90A], Rev. 1). During generation, the length
2650 of any additional input provided during the generation request is adjusted to *seedlen* bits as well
2651 (see SP 800-90A).

³⁹ 128 bits for AES.

2652 Two alternative derivation functions are specified in Appendices [C.3.2](#) and [C.3.3](#). Appendix [C.3.1](#)
2653 discusses the keys and constants for use with the alternative derivation functions specified in
2654 Appendices [C.3.2](#) and [C.3.3](#).

2655 **C.3.1. Derivation Keys and Constants**

2656 Both of the derivation methods specified in Appendices [C.3.2](#) and [C.3.3](#) an AES derivation key
2657 (*df_Key*) whose length shall meet or exceed the instantiated security strength of the DRBG
2658 instantiation.

2659 The *df_Key* **may** be set to any value and **may** be the current value of a key used by the DRBG.

2660 These alternative methods use three 128-bit constants C_1 , C_2 and C_3 , which are defined as:

2661 $C_1 = 000000...00$

2662 $C_2 = 101010...10$

2663 $C_3 = 010101...01$

2664 The value of B used in Appendices [C.3.2](#) and [C.3.3](#) depends on the length of the AES derivation
2665 key (*df_Key*). When the length of *df_Key* = 128 bits, then $B = 2$. Otherwise, $B = 3$.

2666 **C.3.2. Derivation Function Using CMAC**

2667 CMAC is a block-cipher mode of operation specified in [[SP800-38B](#)]. The CMAC_*df* derivation
2668 function is specified as follows:

2669 **CMAC_*df*:**

2670 **Input:** bitstring *input_string*, integer *number_of_bits_to_return*.

2671 **Output:** bitstring Z .

2672 **Process:**

2673 1. Let C_1 , C_2 , C_3 be 128-bit blocks defined as 000000...0, 101010...10, 010101...01,
2674 respectively.

2675 2. Get *df_Key*. Comment: See [Appendix C.3.1](#).

2676 3. Z = the Null string.

2677 4. For $i = 1$ to B :

2678 $Z = Z \parallel \text{CMAC}(\text{df_Key}, C_i \parallel \text{input_string})$.

2679 5. $Z = \text{leftmost}(Z, \text{number_of_bits_to_return})$.

2680 6. Return(Z).

2681 **C.3.3. Derivation Function Using CBC-MAC**

2682 This CBC-MAC derivation function **shall** only be used when the *input_string* has the following
2683 properties:

- 2684 • The length of the *input_string* is always a fixed length.
- 2685 • The length of the *input_string* is an integer multiple of 128 bits. Let m be the number of
- 2686 128-bit blocks in the *input_string*.
- 2687 This derivation function is specified as follows:
- 2688 **CBC-MAC_df:**
- 2689 **Input:** bitstring *input_string*, integer *number_of_bits_to_return*.
- 2690 **Output:** bitstring Z .
- 2691 **Process:**
- 2692 1. Let C_1 , C_2 , C_3 be 128-bit blocks defined as 000000...0, 101010...10, 010101...01,
 - 2693 respectively.
 - 2694 2. Get *df_Key*. Comment: See [Appendix C.3.1](#).
 - 2695 3. Z = the *Null* string.
 - 2696 4. Let *input_string* = $S_1 \parallel S_2 \parallel \dots \parallel S_m$, where the S_i are contiguous 128-bit blocks.
 - 2697 5. For $j = 1$ to B :
 - 2698 5.1 $S_0 = C_j$.
 - 2699 5.2 V = 128-bit block of all zeroes.
 - 2700 5.3 For $i = 0$ to m :

2701 $V = \text{Encrypt}(\text{df_Key}, V \oplus S_i)$. Comment: Perform the cipher operation

2702 specified in [[FIPS197](#)].
 - 2703 5.4 $Z = Z \parallel V$.
 - 2704 6. $Z = \text{leftmost}(Z, \text{number_of_bit_to_return})$.
 - 2705 7. Return(Z).

2706 **Appendix D. List of Symbols, Abbreviations, and Acronyms**

2707 **AES**

2708 Advanced Encryption Standard⁴⁰

2709 **API**

2710 Application Programming Interface

2711 **CAVP**

2712 Cryptographic Algorithm Validation Program

2713 **CDF**

2714 Cumulative Distribution Function

2715 **CMVP**

2716 Cryptographic Module Validation Program

2717 **DRBG**

2718 Deterministic Random Bit Generator⁴¹

2719 **FIPS**

2720 Federal Information Processing Standard

2721 **ITL**

2722 Information Technology Laboratory

2723 **MAC**

2724 Message Authentication Code

2725 **NIST**

2726 National Institute of Standards and Technology

2727 **RAM**

2728 Random Access Memory

2729 **RBG**

2730 Random Bit Generator

2731 **SP**

2732 (NIST) Special Publication

2733 **Sub-DRBG**

2734 Subordinate DRBG

2735 **TDEA**

2736 Triple Data Encryption Algorithm⁴²

2737 **XOR**

2738 Exclusive-Or (operation)

2739 **0^x**

2740 A string of x zeroes

2741 **[x]**

⁴⁰ As specified in [FIPS 197].

⁴¹ Mechanism specified in [SP800-90A].

⁴² As specified in [SP 800-67], Recommendation for the Triple Data Encryption Algorithm (TDEA) Block Cipher.

2742	The ceiling of x ; the least integer number that is not less than the real number x . For example, $\lceil 3 \rceil = 3$, and $\lceil 5.5 \rceil = 6$.
2743	ϵ
2744	A positive constant that is assumed to be smaller than 2^{-32}
2745	$E(X)$
2746	The expected value of the random variable X
2747	$\text{len}(x)$
2748	The length of x in bits
2749	$\min(a, b)$
2750	The minimum of a and b
2751	output_len
2752	The bit length of the output block of a cryptographic primitive
2753	s
2754	The security strength
2755	$X \oplus Y$
2756	Boolean bitwise exclusive-or (also bitwise addition modulo 2) of two bitstrings X and Y of the same length
2757	$+$
2758	Addition over real numbers
2759	\times
2760	Multiplication over real numbers

2761 **Appendix E. Glossary**

2762 **adversary**

2763 A malicious entity whose goal is to determine, to guess, or to influence the output of an RBG.

2764 **approved**

2765 An algorithm or technique for a specific cryptographic use that is specified in a FIPS or NIST Recommendation,
2766 adopted in a FIPS or NIST Recommendation, or specified in a list of NIST-approved security functions.

2767 **backtracking resistance**

2768 A property of a DRBG that provides assurance that compromising the current internal state of the DRBG does not
2769 weaken previously generated outputs. See [SP 800-90A](#) for a more complete discussion. (Contrast with *prediction*
2770 *resistance*.)

2771 **biased**

2772 A random variable is said to be biased if values of the finite sample space are selected with unequal probability.
2773 Contrast with unbiased.

2774 **big-endian format**

2775 A format in which the most significant bytes (the bytes containing the high-order or leftmost bits) are stored in the
2776 lowest address with the following bytes in sequentially higher addresses.

2777 **bitstring**

2778 An ordered sequence (string) of 0s and 1s. The leftmost bit is the most significant bit.

2779 **block cipher**

2780 A parameterized family of permutations on bitstrings of a fixed length; the parameter that determines the permutation
2781 is a bitstring called the key.

2782 **conditioning function (external)**

2783 As used in SP 800-90C, a deterministic function that is used to produce a bitstring with full entropy.

2784 **consuming application**

2785 An application that uses random outputs from an RBG.

2786 **cryptographic boundary**

2787 An explicitly defined physical or conceptual perimeter that establishes the physical and/or logical bounds of a
2788 cryptographic module and contains all of the hardware, software, and/or firmware components of a cryptographic
2789 module.

2790 **cryptographic module**

2791 The set of hardware, software, and/or firmware that implements cryptographic functions (including cryptographic
2792 algorithms and key generation) and is contained within the cryptographic boundary.

2793 **deterministic random bit generator (DRBG)**

2794 An RBG that produces random bitstrings by applying a deterministic algorithm to initial seed material.

2795 *Note:* A DRBG at least has access to a randomness source initially.

2796 *Note:* A portion of the seed material is secret.

2797 **digitization**

2798 The process of generating raw discrete digital values from non-deterministic events (e.g., analog noise sources) within
2799 a noise source.

2800 **entropy**

2801 A measure of disorder, randomness, or variability in a closed system.

2802 *Note:* The entropy of a random variable X is a mathematical measure of the amount of information gained by an
2803 observation of X .

2804 *Note:* The most common concepts are Shannon entropy and min-entropy. Min-entropy is the measure used in SP 800-
2805 90.

2806 **entropy rate**

2807 The validated rate at which an entropy source provides entropy in terms of bits per entropy-source output (e.g., five
2808 bits of entropy per eight-bit output sample).

2809 **entropy source**

2810 The combination of a noise source, health tests, and optional conditioning component that produce bitstrings
2811 containing entropy. A distinction is made between entropy sources having physical noise sources and those having
2812 non-physical noise sources.

2813 *Note:* Health tests are comprised of continuous tests and startup tests.

2814 **fresh entropy**

2815 A bitstring that is output from a non-deterministic randomness source that has not been previously used to generate
2816 output or has otherwise been made externally available.

2817 *Note:* The randomness source should be an entropy source or RBG3 construction.

2818 **full-entropy bitstring**

2819 A bitstring with ideal randomness (i.e., the amount of entropy per bit is equal to 1). This Recommendation assumes
2820 that a bitstring has *full entropy* if the entropy rate is at least $1 - \varepsilon$, where ε is at most 2^{-32} .

2821 **hash function**

2822 A (mathematical) function that maps values from a large (possibly very large) domain into a smaller range. The
2823 function satisfies the following properties:

- 2824 1. (One-way) It is computationally infeasible to find any input that maps to any pre-specified output.
2825 2. (Collision-free) It is computationally infeasible to find any two distinct inputs that map to the same output.

2826 **health testing**

2827 Testing within an implementation immediately prior to or during normal operations to obtain assurance that the
2828 implementation continues to perform as implemented and validated.

2829 **ideal randomness source**

2830 The source of an ideal random sequence of bits. Each bit of an ideal random sequence is unpredictable and unbiased,
2831 with a value that is independent of the values of the other bits in the sequence. Prior to an observation of the sequence,
2832 the value of each bit is equally likely to be 0 or 1, and the probability that a particular bit will have a particular value
2833 is unaffected by knowledge of the values of any or all of the other bits. An ideal random sequence of n bits contains n
2834 bits of entropy.

2835 **independent entropy sources**

2836 Two entropy sources are *independent* if knowledge of the output of one entropy source provides no information about
2837 the output of the other entropy source.

2838 **instantiate**

2839 The process of initializing a DRBG with sufficient randomness to generate pseudorandom bits at the desired security
2840 strength.

2841 **internal state (of a DRBG)**

2842 The collection of all secret and non-secret information about an RBG or entropy source that is stored in memory at a
2843 given point in time.

2844 **known-answer test**

2845 A test that uses a fixed input/output pair to detect whether a deterministic component was implemented correctly or
2846 to detect whether it continues to operate correctly.

2847 **min-entropy**

2848 A lower bound on the entropy of a random variable. The precise formulation for min-entropy is $(-\log_2 \max p_i)$ for a
2849 discrete distribution having probabilities p_1, \dots, p_k . Min-entropy is often used as a measure of the unpredictability of a
2850 random variable.

2851 **must**

2852 Used in SP 800-90C to indicate a requirement that may not be testable by a CMVP testing lab. Note that **must** may
2853 be coupled with **not** to become **must not**.

2854 **noise source**

2855 A source of unpredictable data that outputs raw discrete digital values. The digitization mechanism is considered part
2856 of the noise source. A distinction is made between physical noise sources and non-physical noise sources.

2857 **non-physical entropy source**

2858 An entropy source whose primary noise source is non-physical.

2859 **non-physical noise source**

2860 A noise source that typically exploits system data and/or user interaction to produce digitized random data.

2861 **non-validated entropy source**

2862 An entropy source that has not been validated by the CMVP as conforming to [SP 800-90B](#).

2863 **null string**

2864 An empty bitstring.

2865 **personalization string**

2866 An optional input value to a DRBG during instantiation to make one DRBG instantiation behave differently from
2867 other instantiations.

2868 **physical entropy source**

2869 An entropy source whose primary noise source is physical.

2870 **physical noise source**

2871 A noise source that exploits physical phenomena (e.g., thermal noise, shot noise, jitter, metastability, radioactive
2872 decay, etc.) from dedicated hardware designs (using diodes, ring oscillators, etc.) or physical experiments to produce
2873 digitized random data.

2874 **prediction resistance**

2875 A property of a DRBG that provides assurance that compromising the current internal state of the DRBG does not
2876 allow future DRBG outputs to be predicted past the point where the DRBG has been reseeded with sufficient entropy.
2877 See [SP 800-90A](#) for a more complete discussion. (Contrast with *backtracking resistance*.)

2878 **pseudocode**

2879 An informal, high-level description of a computer program, algorithm, or function that resembles a simplified
2880 programming language.

2881 **random bit generator (RBG)**

2882 A device or algorithm that outputs a random sequence that is effectively indistinguishable from statistically
2883 independent and unbiased bits.

2884 **randomness**

2885 As used in this Recommendation, the unpredictability of a bitstring. If the randomness is produced by a non-deterministic
2886 source (e.g., an entropy source or RBG3 construction), the unpredictability is dependent on the quality of the source. If

2887 the randomness is produced by a deterministic source (e.g., a DRBG), the unpredictability is based on the capability of
2888 an adversary to break the cryptographic algorithm for producing the pseudorandom bitstring.

2889 **randomness input**

2890 An input bitstring from a randomness source that provides an assessed minimum amount of randomness (e.g., entropy)
2891 for a DRBG. See *min-entropy*.

2892 **randomness source**

2893 A source of randomness for an RBG. The randomness source may be an entropy source or an RBG construction.

2894 **RBG1 construction**

2895 An RBG construction with the DRBG and the randomness source in separate cryptographic modules.

2896 **RBG2 construction**

2897 An RBG construction with one or more entropy sources and a DRBG within the same cryptographic module. This
2898 RBG construction does not provide full-entropy output.

2899 **RBG2(NP) construction**

2900 A non-physical RBG2 construction. An RBG2 construction that obtains entropy from one or more validated non-
2901 physical entropy sources and possibly from one or more validated physical entropy sources. This RBG construction
2902 does not provide full-entropy output.

2903 **RBG2(P) construction**

2904 A physical RBG2 construction. An RBG construction that includes a DRBG and one or more entropy sources in the
2905 same cryptographic module. Only the entropy from validated physical entropy sources is counted when fulfilling an
2906 entropy request within the RBG. This RBG construction does not provide full-entropy output.

2907 **RBG3 construction**

2908 An RBG construction that includes a DRBG and one or more entropy sources in the same cryptographic module.
2909 When working properly, bitstrings that have full entropy are produced. Sometimes called a *non-deterministic random*
2910 *bit generator* (NRBG) or true random number (or bit) *generator*.

2911 **reseed**

2912 To refresh the internal state of a DRBG with seed material. The seed material should contain sufficient entropy to
2913 allow recovery from a possible compromise.

2914 **sample space**

2915 The set of all possible outcomes of an experiment.

2916 **secure channel**

2917 A physically protected secure path for transferring data between two cryptographic modules that ensures
2918 confidentiality, integrity, and replay protection as well as mutual authentication between the modules.

2919 **security boundary**

2920 For an entropy source: A conceptual boundary that is used to assess the amount of entropy provided by the values
2921 output from the entropy source. The entropy assessment is performed under the assumption that any observer
2922 (including any adversary) is outside of that boundary during normal operation.

2923 For a DRBG: A conceptual boundary that contains all of the DRBG functions and internal states required for a DRBG.

2924 For an RBG: A conceptual boundary that is defined with respect to one or more threat models that includes an
2925 assessment of the applicability of an attack and the potential harm caused by the attack.

2926 **security strength**

2927 A number associated with the amount of work (i.e., the number of basic operations of some sort) that is required to
2928 “break” a cryptographic algorithm or system in some way. In this Recommendation, the security strength is specified
2929 in bits and is a specific value from the set {128, 192, 256}. If the security strength associated with an algorithm or
2930 system is s bits, then it is expected that (roughly) 2^s basic operations are required to break it.

- 2931 *Note:* This is a classical definition that does not consider quantum attacks. This definition will be revised to address
2932 quantum issues in the future.
- 2933 **seed**
2934 To initialize the internal state of a DRBG with seed material. The seed material should contain sufficient entropy to
2935 meet security requirements.
- 2936 **seed material**
2937 A bitstring that is used as input to a DRBG. The seed material determines a portion of the internal state of the DRBG.
- 2938 **seedlife**
2939 The period of time between instantiating or reseeding a DRBG with seed material and reseeding the DRBG with seed
2940 material containing fresh entropy or uninstantiation of the DRBG.
- 2941 **shall**
2942 The term used to indicate a requirement that is testable by a testing lab. **Shall** may be coupled with **not** to become
2943 **shall not**. See *Testable requirement*.
- 2944 **should**
2945 The term used to indicate an important recommendation. Ignoring the recommendation could result in undesirable
2946 results. Note that **should** may be coupled with **not** to become **should not**.
- 2947 **state handle**
2948 A pointer to the internal state information for a particular DRBG instantiation.
- 2949 **subordinate DRBG (sub-DRBG)**
2950 A DRBG that is instantiated by an RBG1 construction.
- 2951 **support a security strength (by a DRBG)**
2952 The DRBG has been instantiated at a security strength that is equal to or greater than the security strength requested
2953 for the generation of random bits.
- 2954 **targeted security strength**
2955 The security strength that is intended to be supported by one or more implementation-related choices (e.g., algorithms,
2956 cryptographic primitives, auxiliary functions, parameter sizes, and/or actual parameters).
- 2957 **testable requirement**
2958 A requirement that can be tested for compliance by a testing lab via operational testing, a code review, or a review of
2959 relevant documentation provided for validation. A testable requirement is indicated using a **shall** statement.
- 2960 **threat model**
2961 A description of a set of security aspects that need to be considered. A threat model can be defined by listing a set of
2962 possible attacks along with the probability of success and the potential harm from each attack.
- 2963 **unbiased**
2964 A random variable is said to be unbiased if all values of the finite sample space are chosen with the same
2965 probability. Contrast with biased.
- 2966 **uninstantiate**
2967 The termination of a DRBG instantiation.
- 2968 **validated entropy source**
2969 An entropy source that has been successfully validated by the CAVP and CMVP for conformance to [SP 800-90B](#).